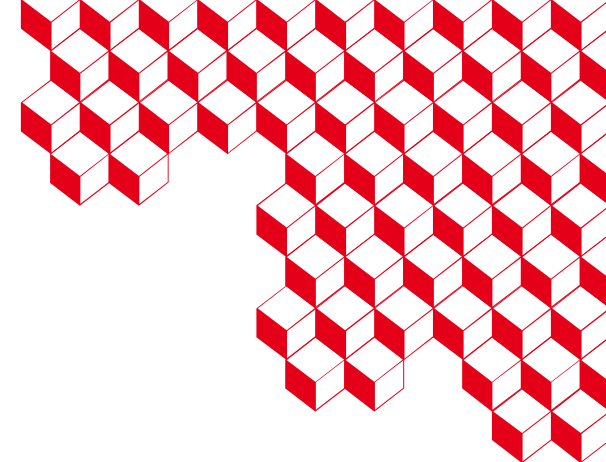




list



Performance evaluation and control improvements for solving optimization problems on Noisy Intermediate-Scale Quantum (NISQ) platforms

PhD Thesis defense

Valentin GILBERT

18th December 2024

Advisors: Stéphane LOUISE, Renaud SIRDEY

Reviewers: Frank PHILLIPSON, Caroline PRODHON

Examiners: Daniel ESTEVE, Jeannette LORENZ

Introduction and Context

Collection criteria

- Score of interest (beauty, color, shape, brand ...)



Introduction and Context

Collection criteria

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$$S_{\text{interest}} = 0.03$$

Introduction and Context

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$$S_{\text{interest}} = 0.8$$

Introduction and Context

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$$S_{\text{interest}} = 0.8$$

- Similarity score: color and shape similarity

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$$S_{\text{similarity}} = 0.95$$

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Collection representation

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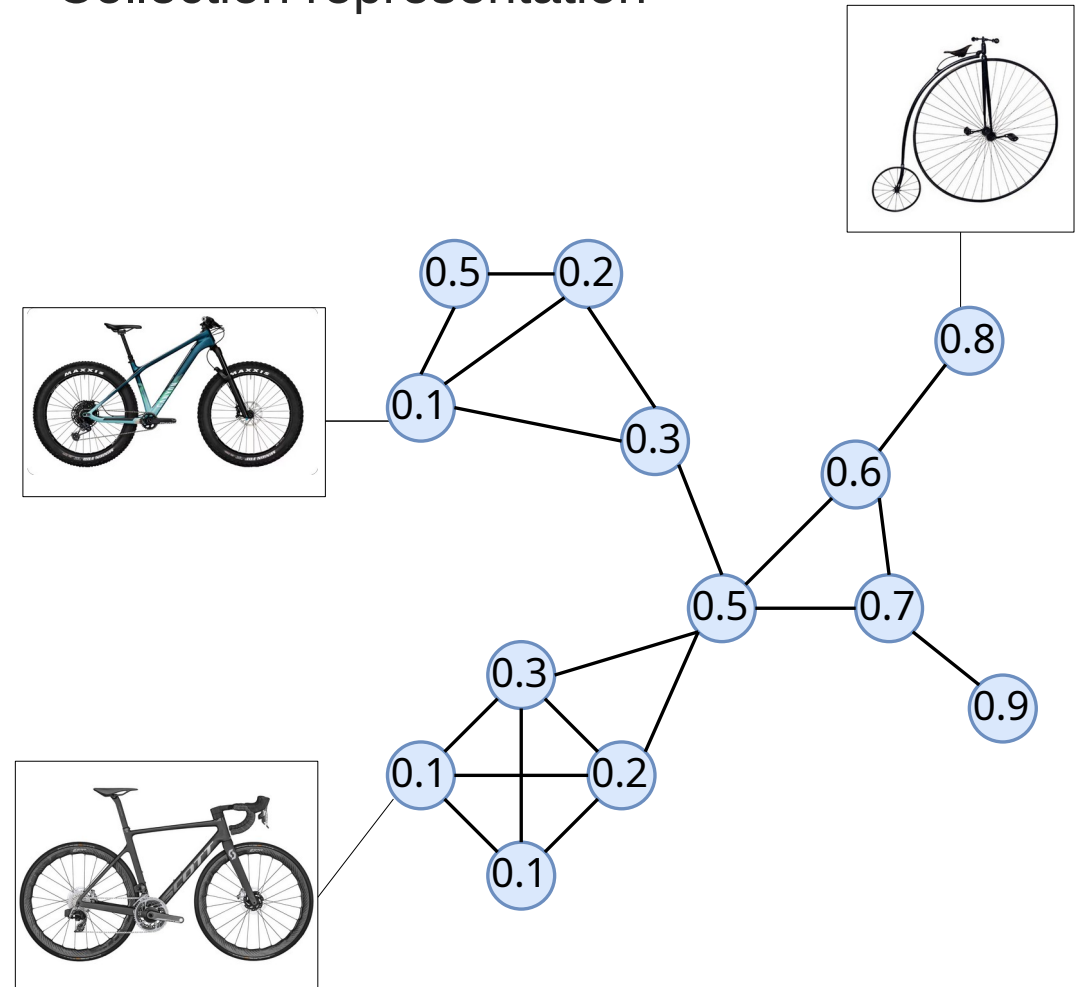
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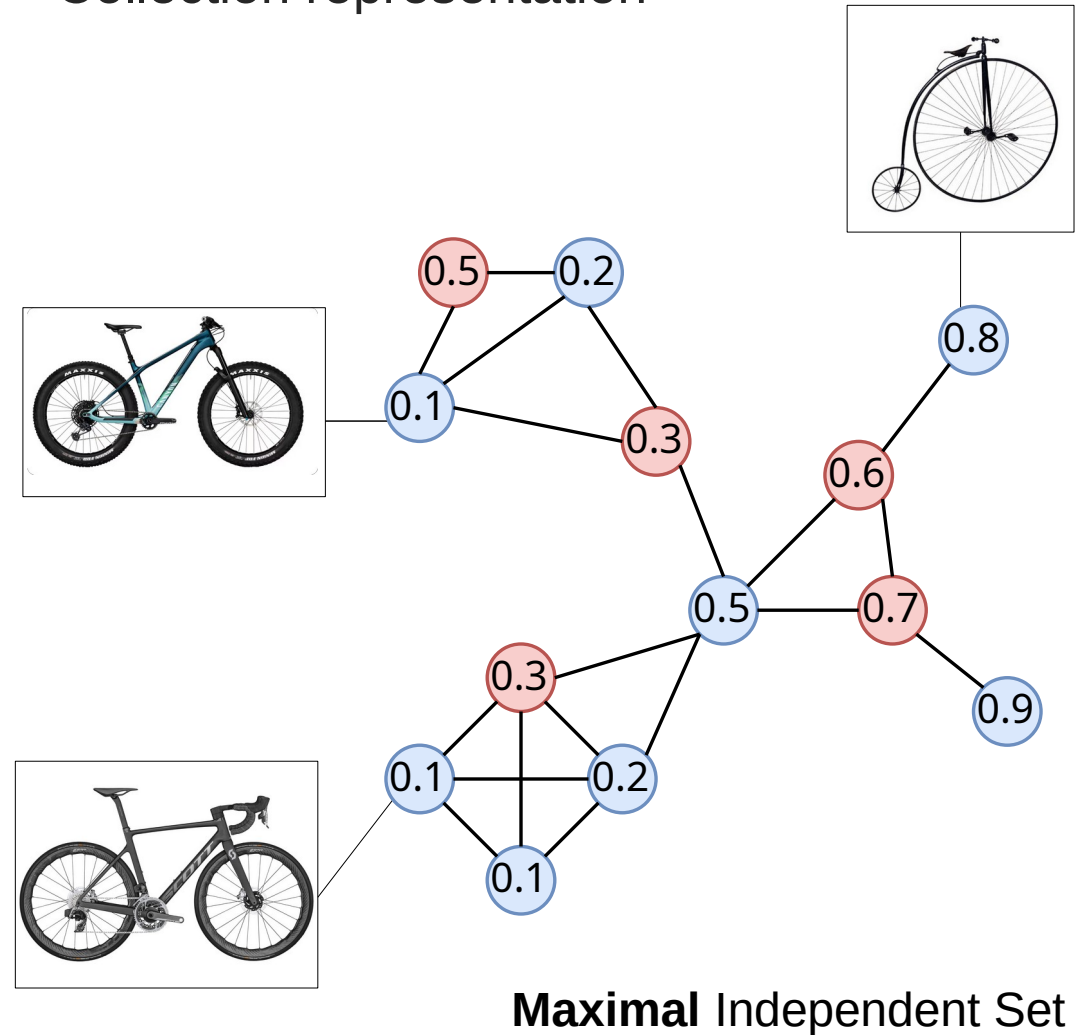
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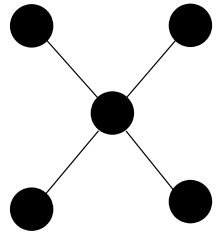
Introduction and Context

Introduction and Context

Maximum Independent Set (MIS) problem

$$G = (V_s, E_s)$$

$$x_v \in \{0, 1\}$$

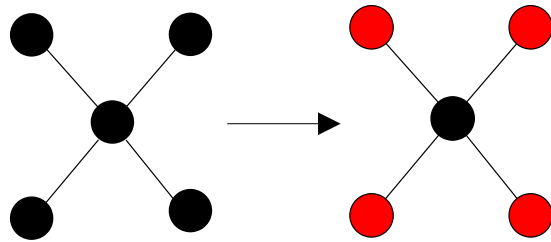


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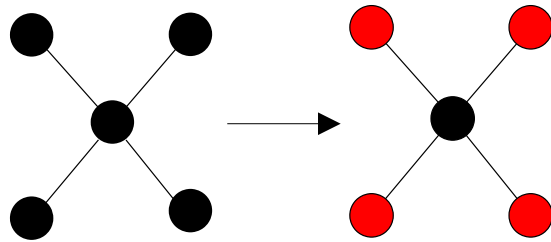


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■ Problem formulation

$$\max \sum_{v \in V_s} x_v$$

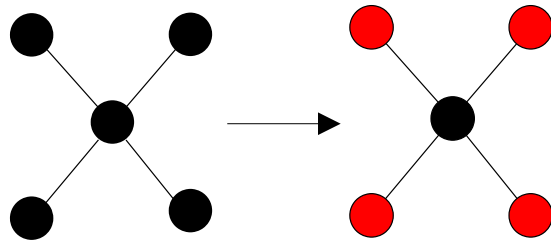
subject to $x_u \neq x_v$ if $(u, v) \in E_s$

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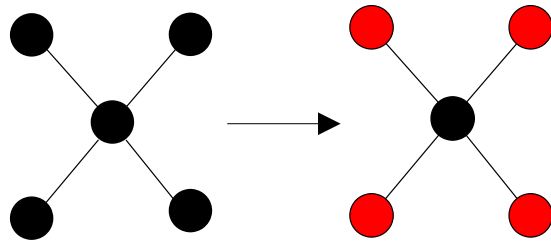
$$\max \sum_{v \in V_s} x_v - 2 \sum_{(u, v) \in E_s} x_u x_v$$

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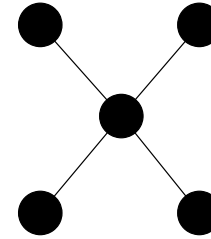
$$x_v \in \{0, 1\}$$



Max-cut problem

$$G = (V_s, E_s)$$

$$s_v \in \{-1, +1\}$$



■ Problem formulation

$$\max \sum_{v \in V_s} x_v$$

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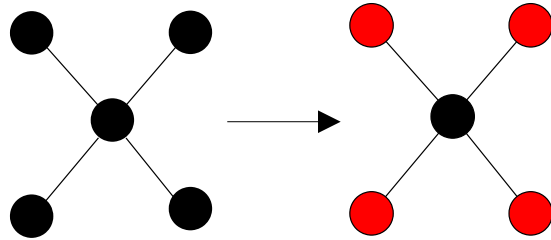
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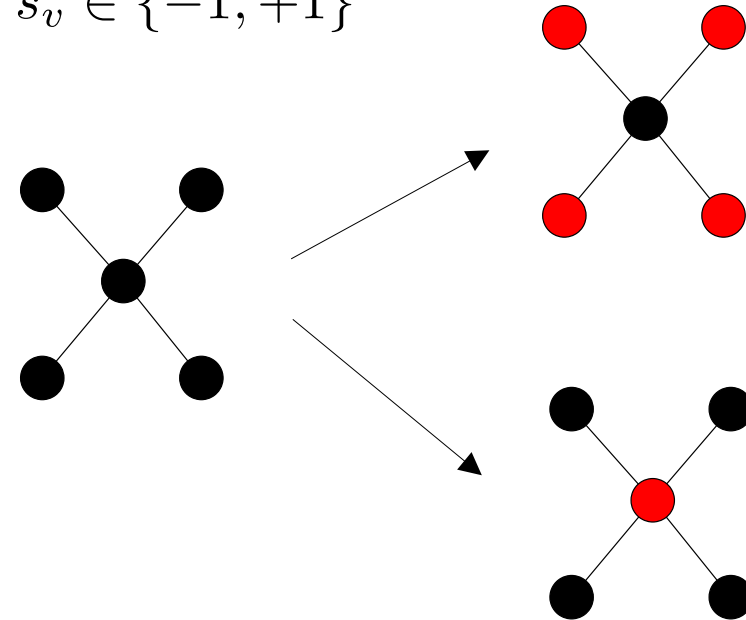
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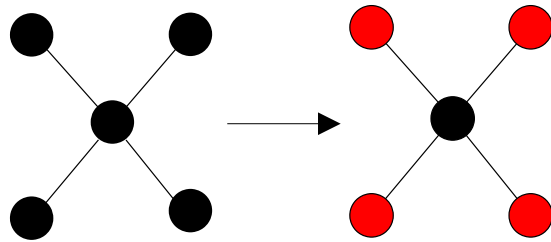


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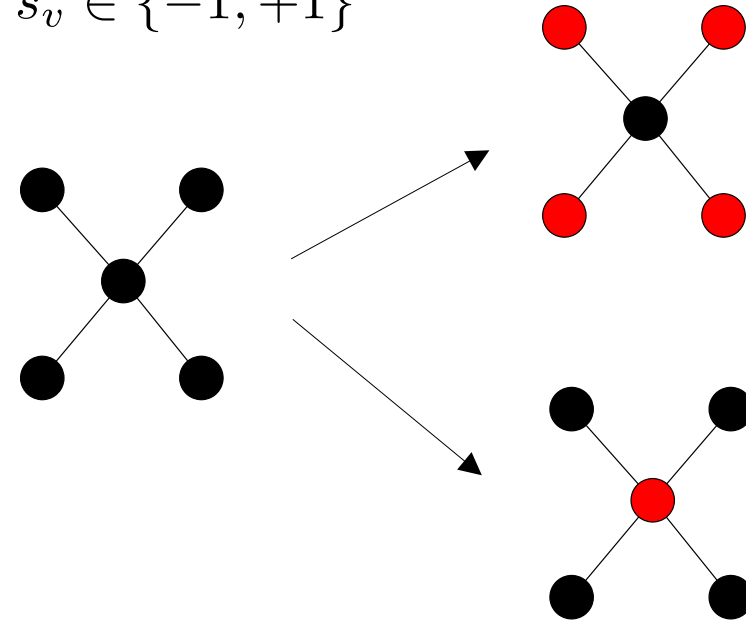
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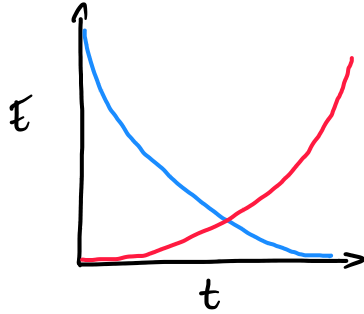
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Introduction - Quantum approaches for optimization problems



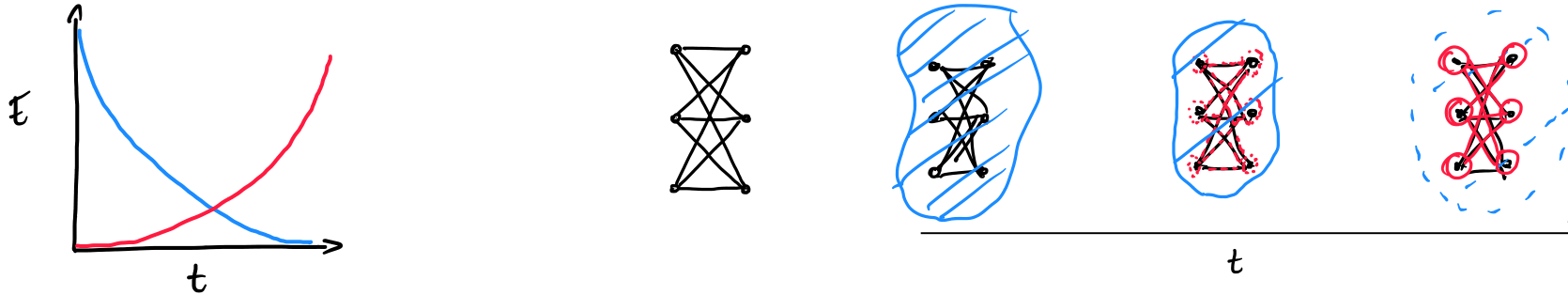
Introduction - Quantum approaches for optimization problems

- Quantum Annealers (analog-based) (NISQ) [KN98]



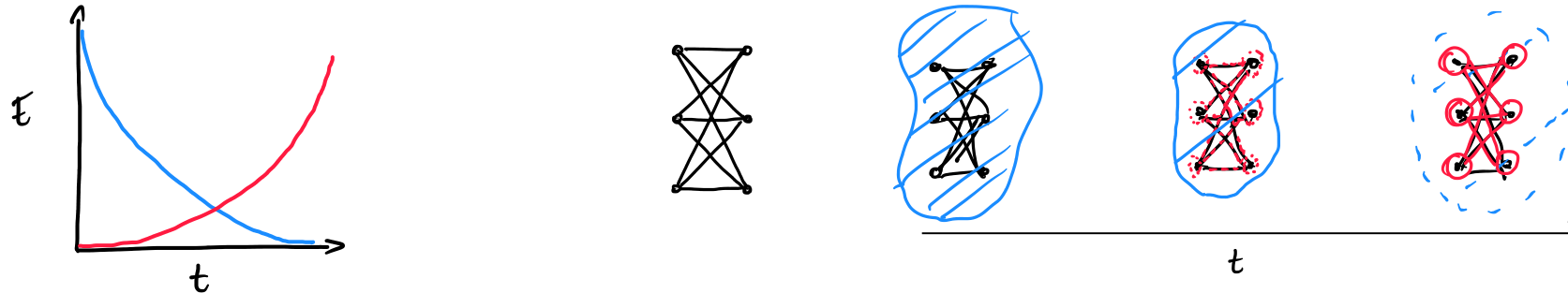
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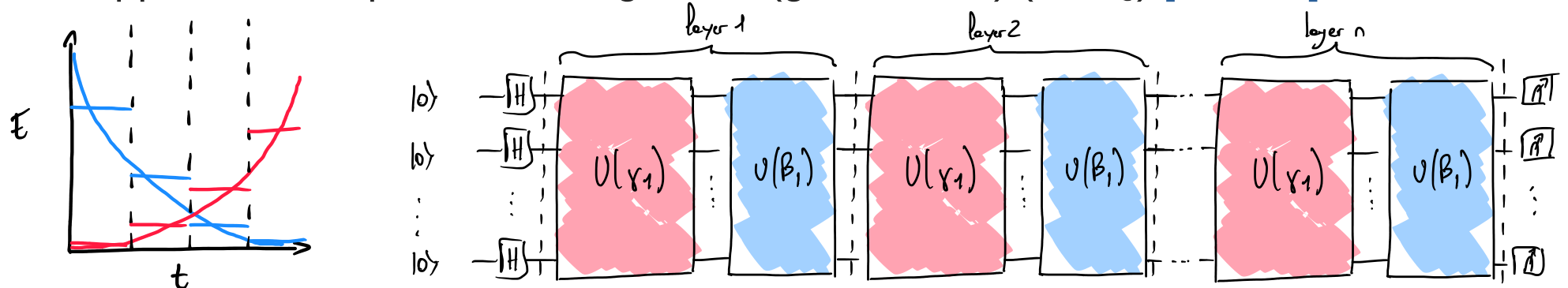


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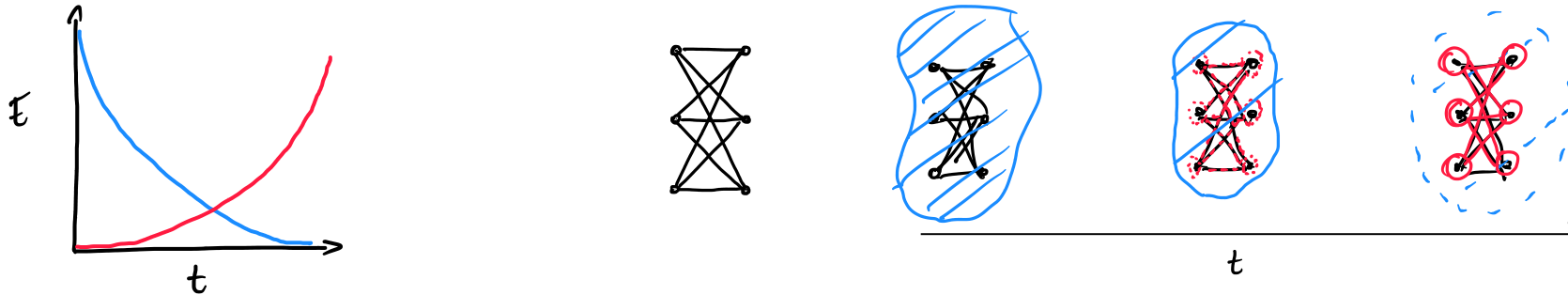


- The Quantum Approximate Optimization Algorithm (gate-based) (NISQ) [FGS14]

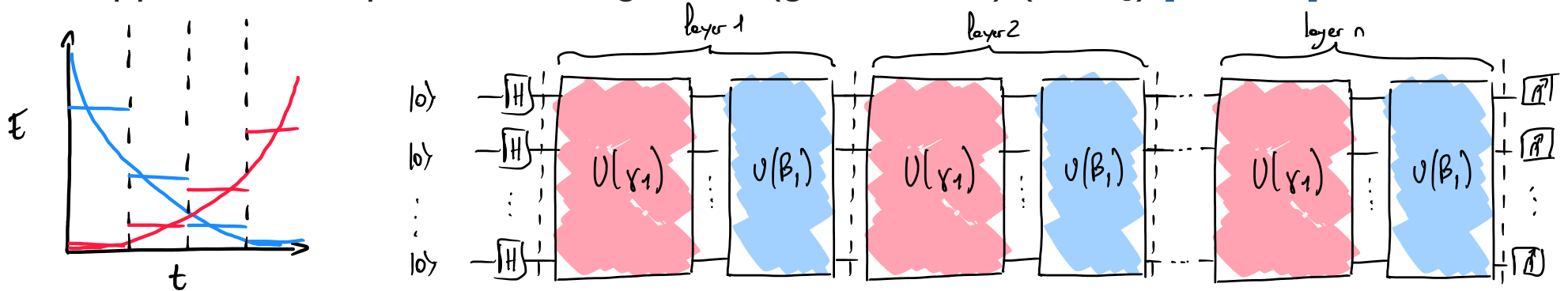


Introduction - Quantum approaches for optimization problems

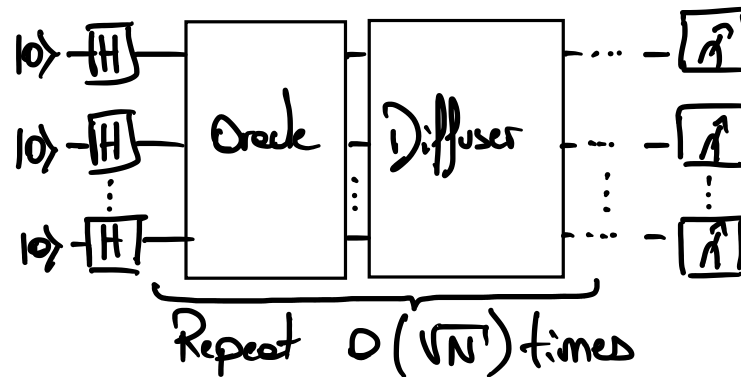
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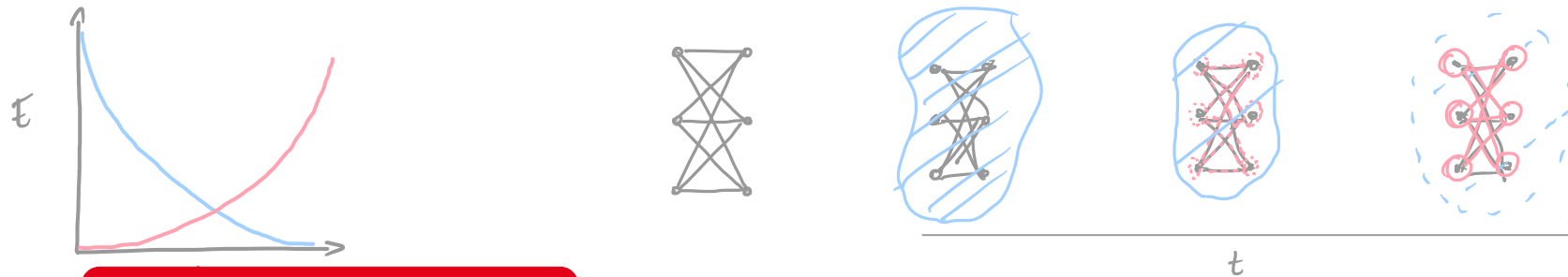


- Exact resolution using quantum circuits (gate-based) (not NISQ) [Gro96]



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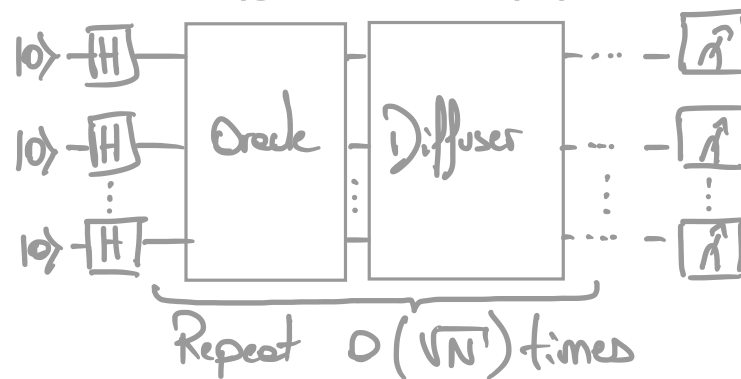


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Research question

- What makes a Quantum Annealer perform well?
- How to improve their performance?

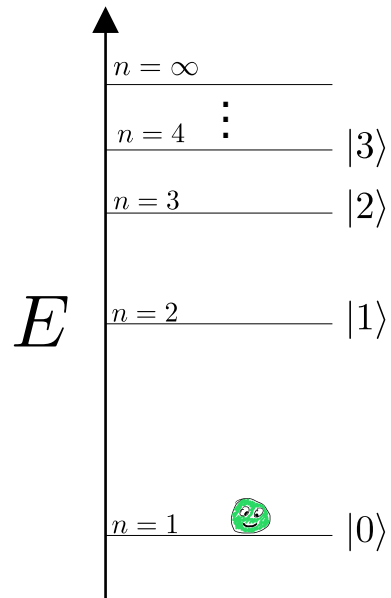
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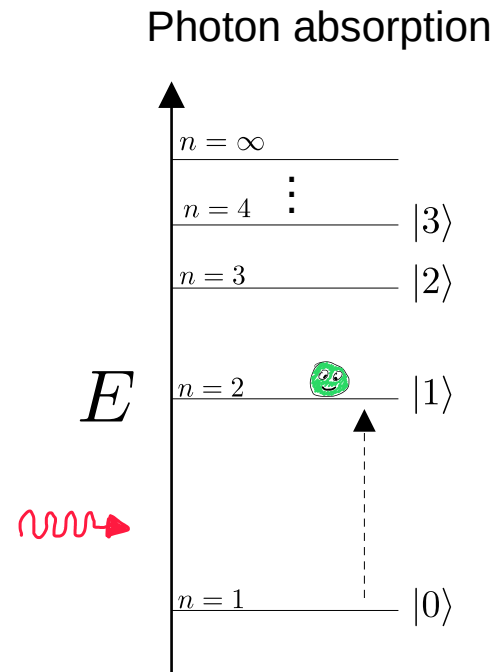


1 ■ Introduction to Quantum Annealing

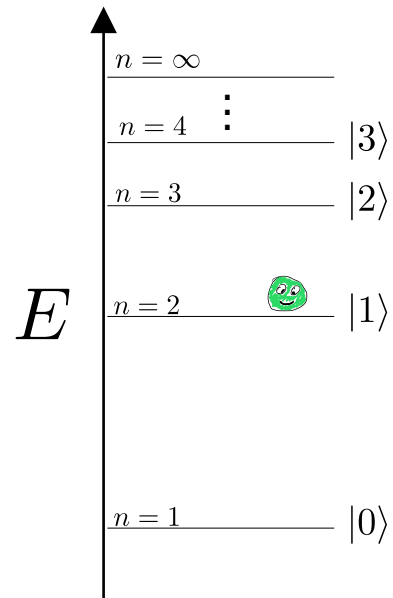
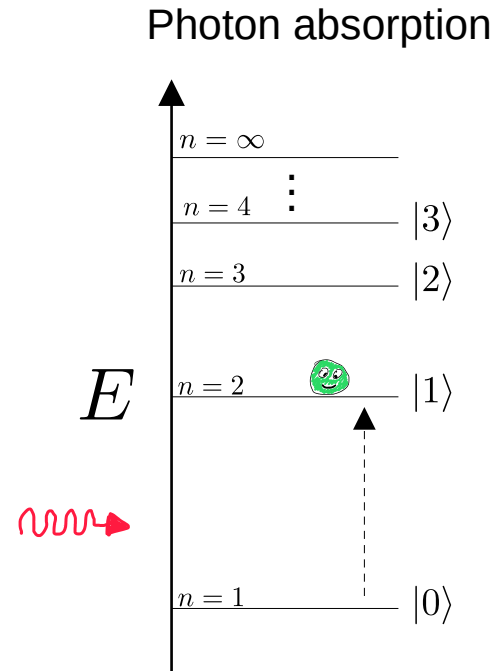
1- Ground state



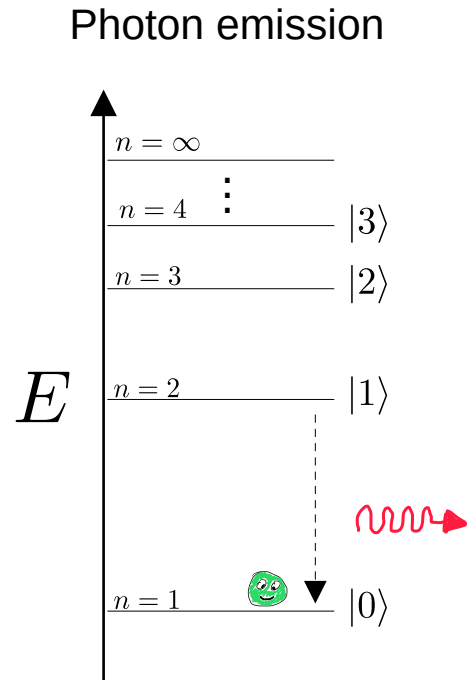
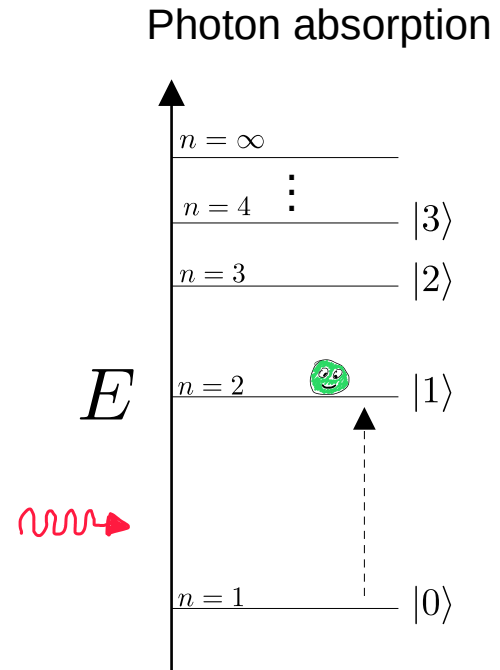
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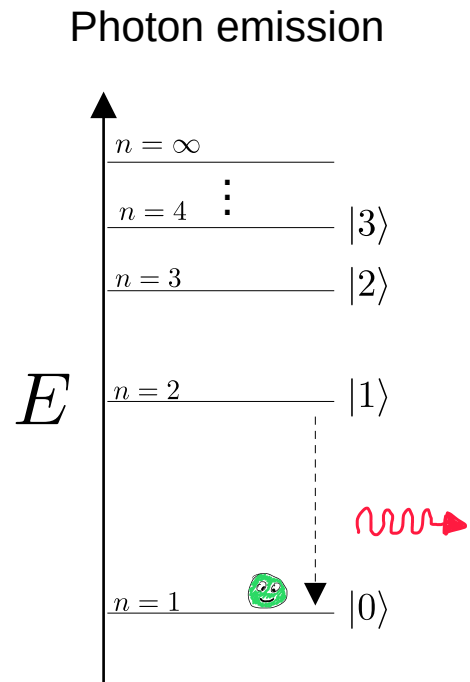
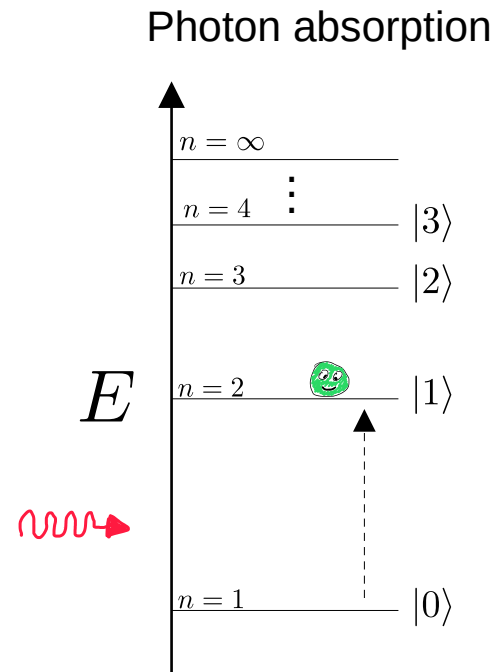
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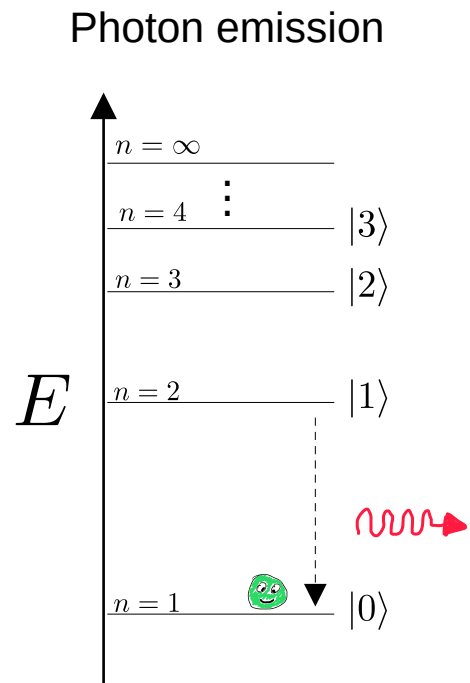
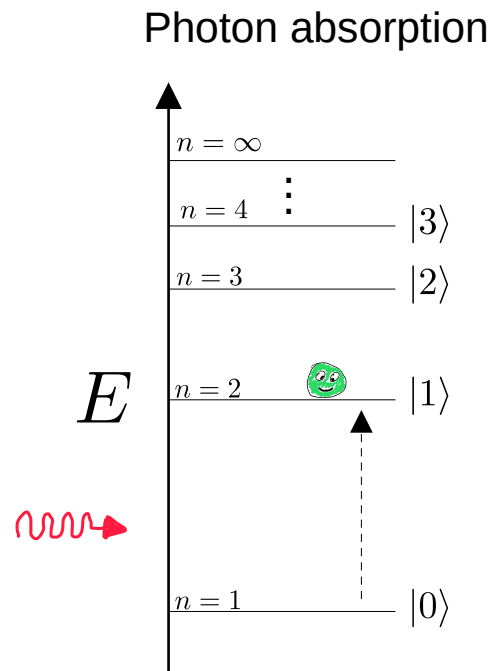


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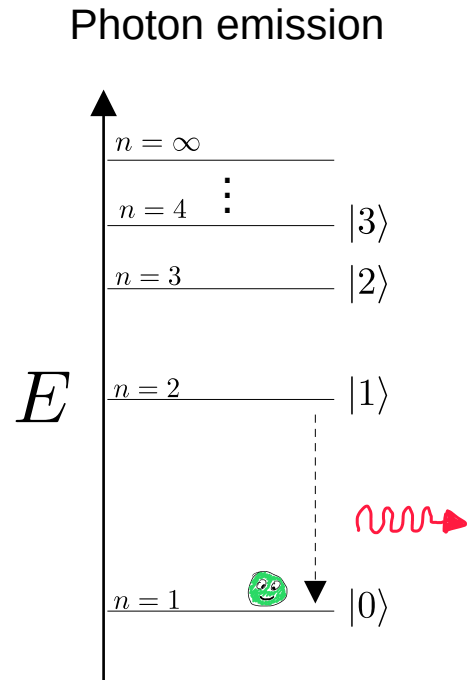
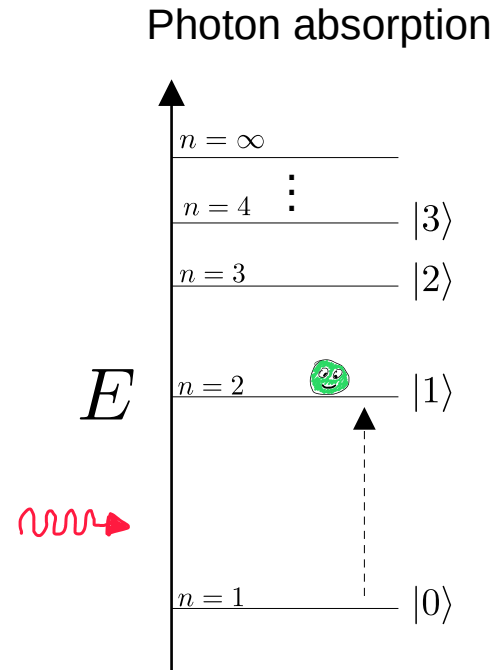


1- Ground state

- Hamiltonian
Energy of distinguishable states



1- Ground state

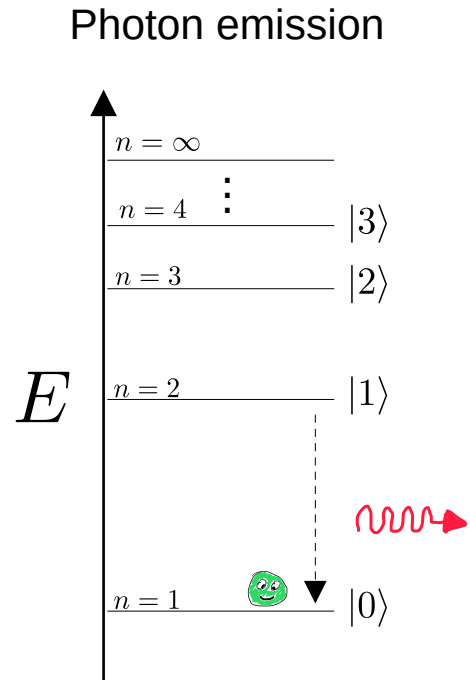
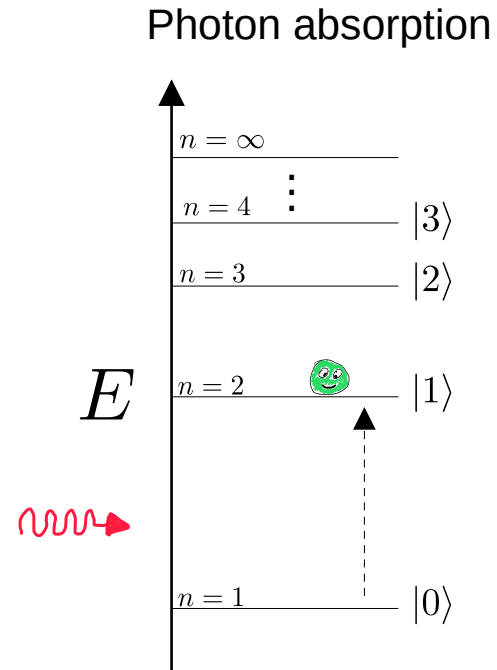


■ Hamiltonian

Energy of distinguishable states

$$H |0\rangle = E_0 |0\rangle \quad |0\rangle = \begin{pmatrix} 1 \\ 0 \end{pmatrix}$$

1- Ground state



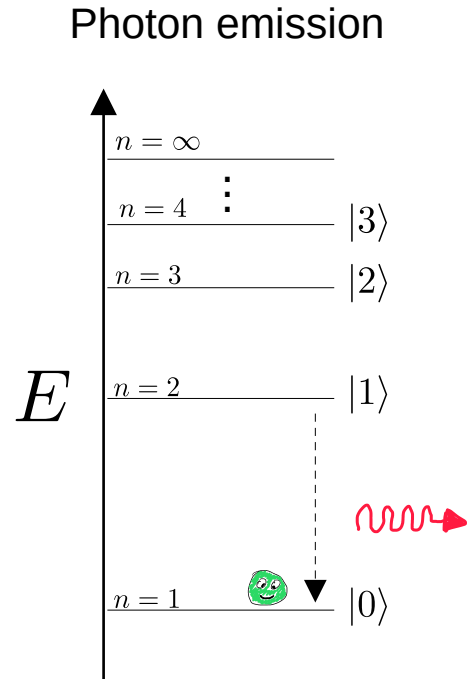
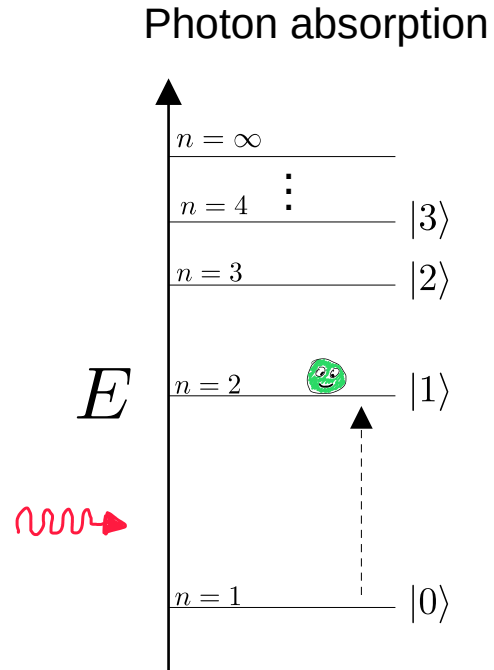
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1- Ground state



■ Hamiltonian

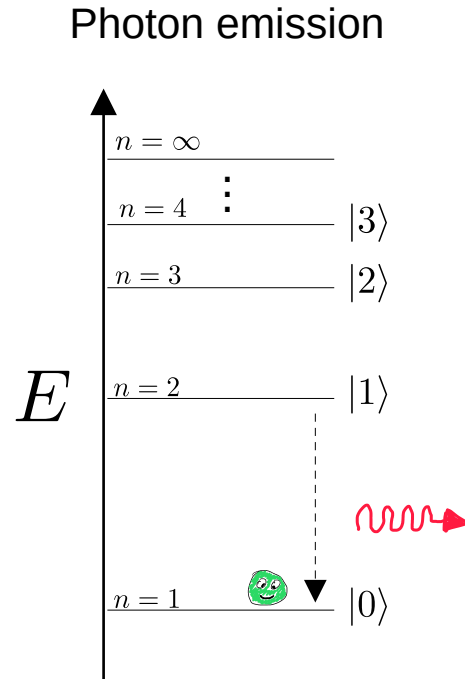
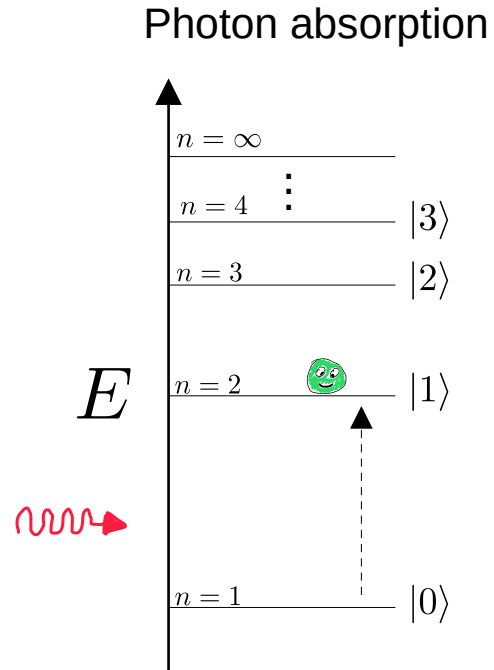
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■ n -qubit Hamiltonian

1- Ground state



■ Hamiltonian

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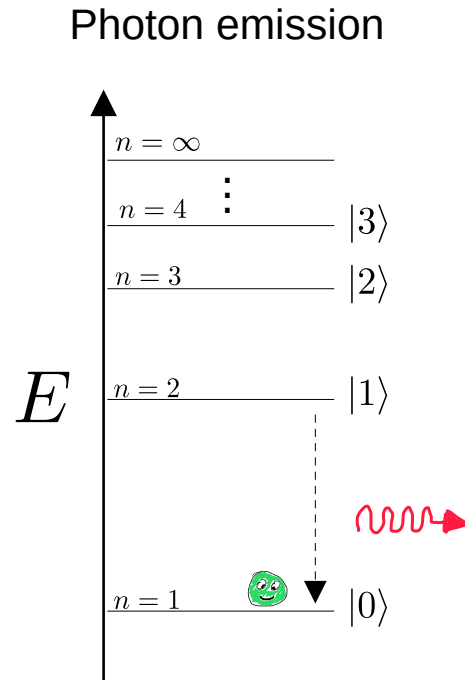
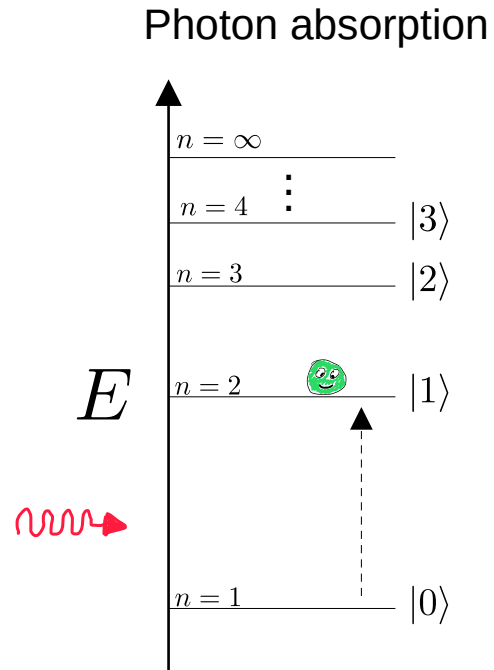
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■ n -qubit Hamiltonian

$$H |\psi_p\rangle = E_p |\psi_p\rangle$$

1- Ground state



■ Hamiltonian

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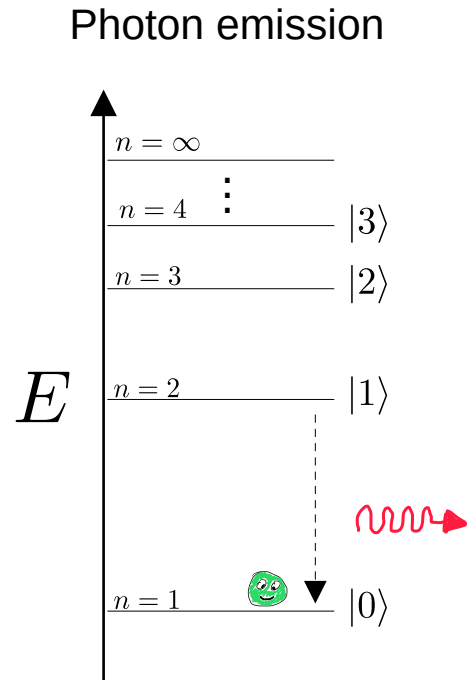
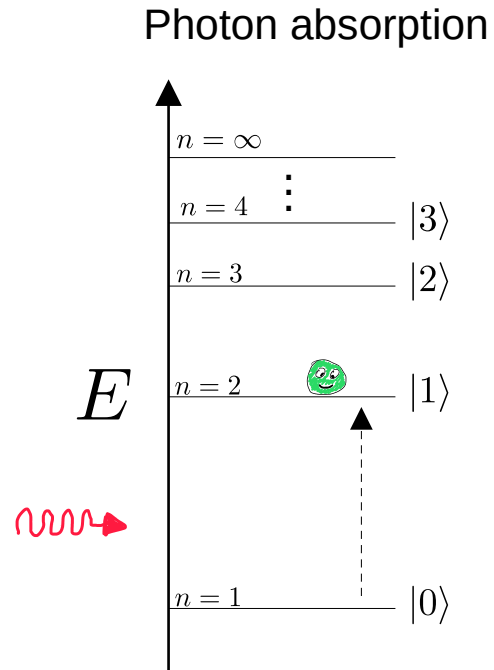
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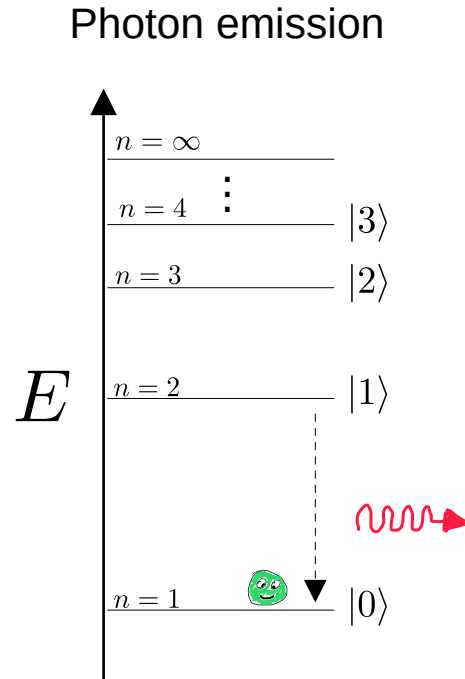
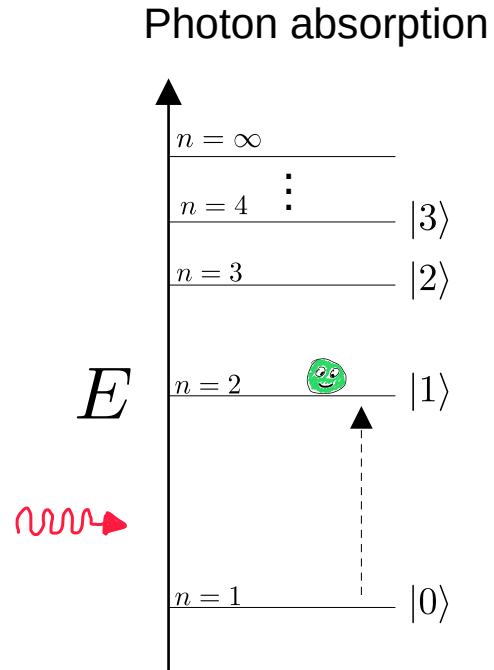
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■ Time-dependent Hamiltonian

1- Ground state



■ Hamiltonian

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■ Time-dependent Hamiltonian

$$H(t) |\psi_p(t)\rangle = E_p(t) |\psi_p(t)\rangle$$

1- Adiabatic theorem

Quantum Adiabatic Theorem

“A quantum system described by a time-dependent Hamiltonian $H(t)$ initially prepared in an eigenstate $|\psi_p(0)\rangle$ of $H(0)$ (e.g., the ground state), will approximately remain in the instantaneous eigenstate $|\psi_p(t)\rangle$, given that t varies sufficiently slowly”

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- Linear interpolation of Hamiltonians:

$$H(t) = -A(t)H_{init} + B(t)H_{final}$$

$$t \in [0, T]$$

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- Assumption

- Ideal closed quantum system

1- Adiabatic theorem

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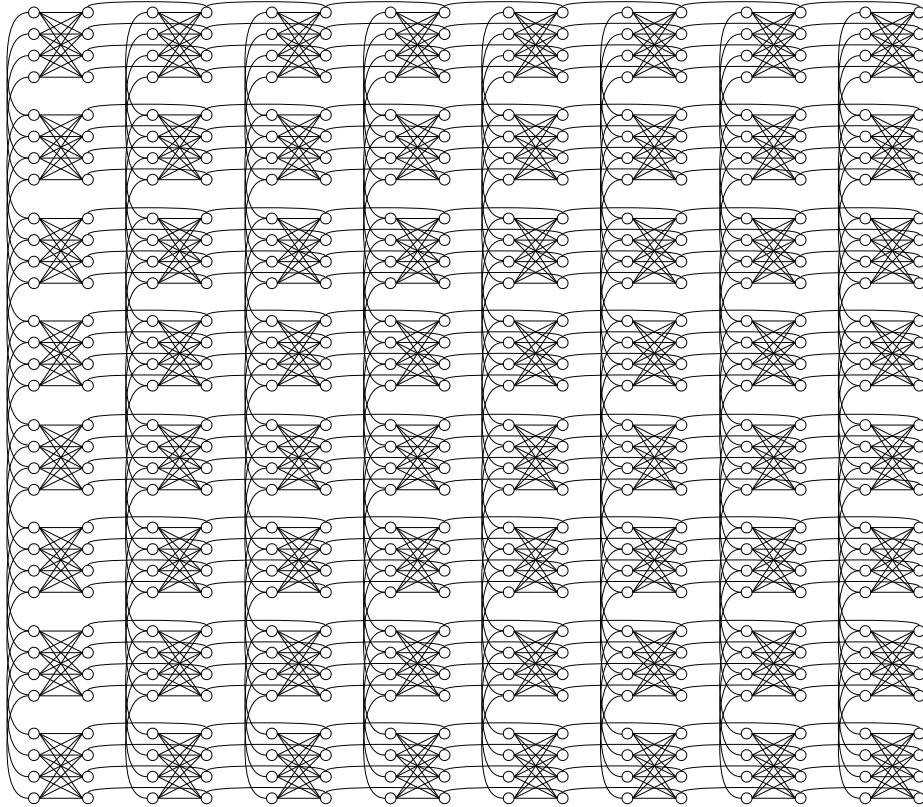
- Assumption

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- Theoretical advantage:

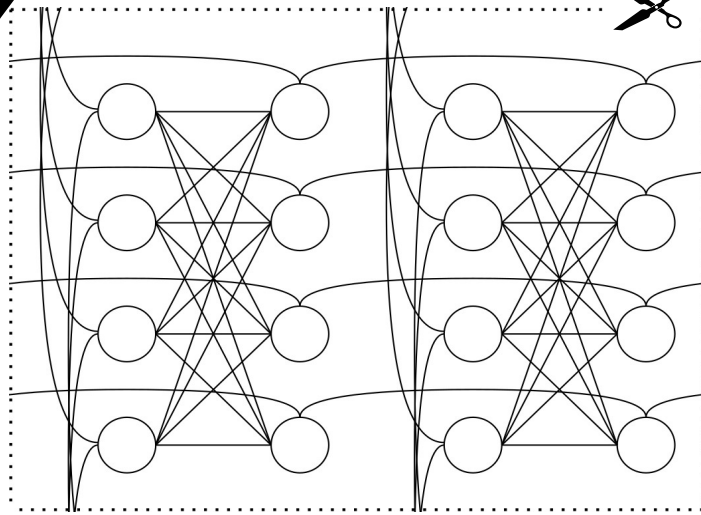
- Universal (as quantum circuit)
- Performance depends on the « sufficiently slowly »

1- Quantum Annealers (D-Wave systems)



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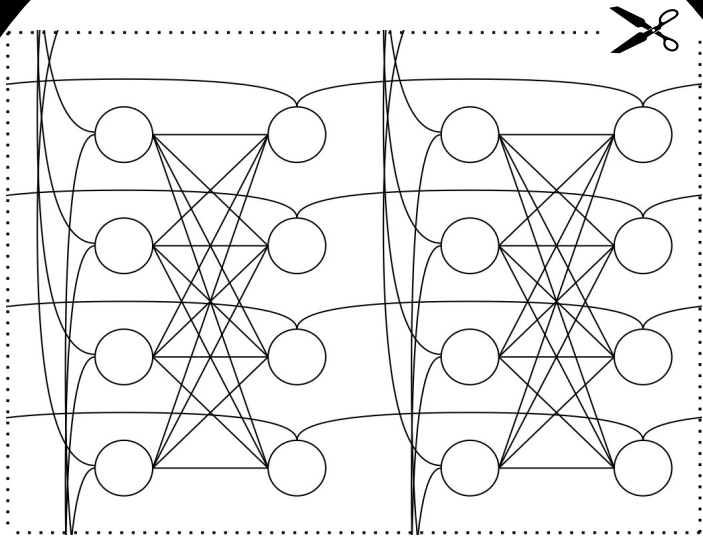
Qubit interaction graph



$$G_t = (V_t, E_t)$$

1- Quantum Annealers (D-Wave systems)

Qubit interaction graph



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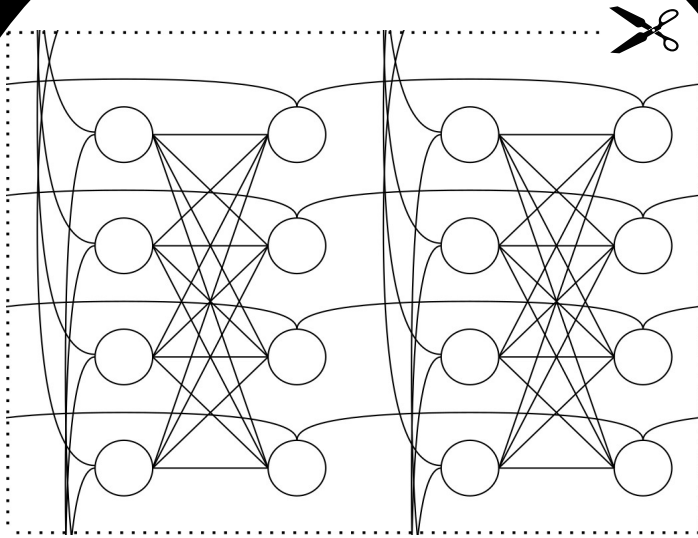
- System Hamiltonian (Transverse Field Ising model)

$$H(s) = -A(s)H_{init} + B(s)H_{final}$$

$$s = \frac{t}{T}, s \in [0, 1]$$

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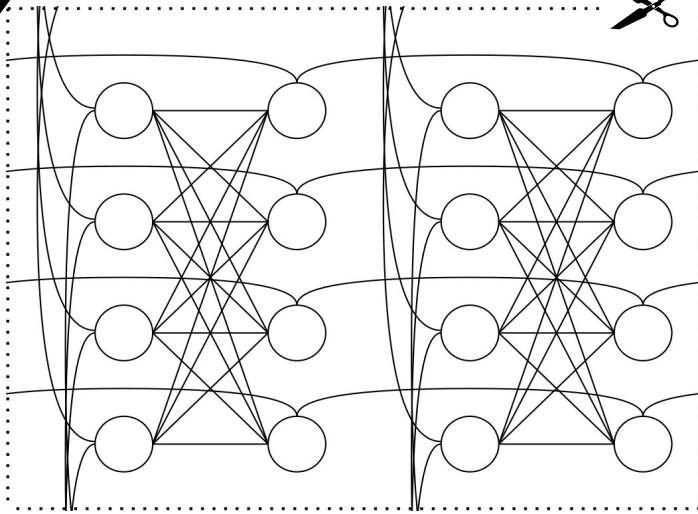
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- Initial Hamiltonian and ground state:

$$H_{init} = \sum_{v \in V_t} \sigma_v^x \quad \sigma^x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

1- Quantum Annealers (D-Wave systems)

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- What eigenstate minimizes the expression $-A(s) H_{init}$?

$$|\psi\rangle = \frac{1}{\sqrt{2^n}} \sum_{x=0}^{2^n-1} |x\rangle$$

1- Quantum Annealers

■ Problem Hamiltonian

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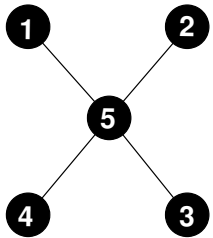
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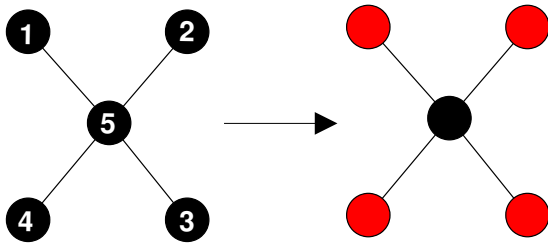
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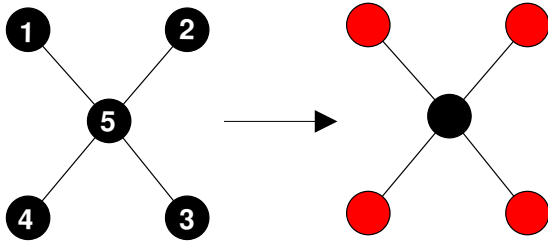
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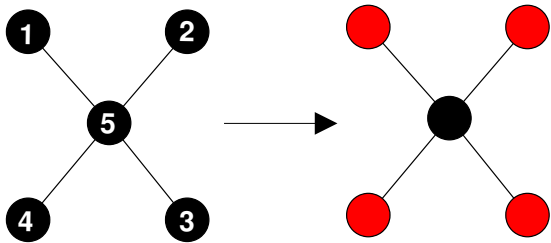
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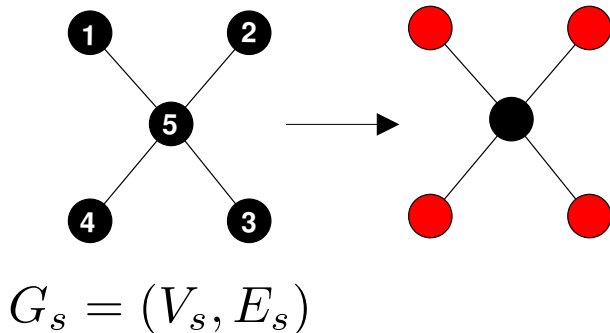
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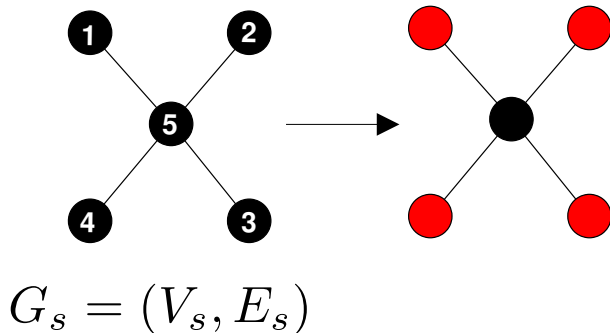
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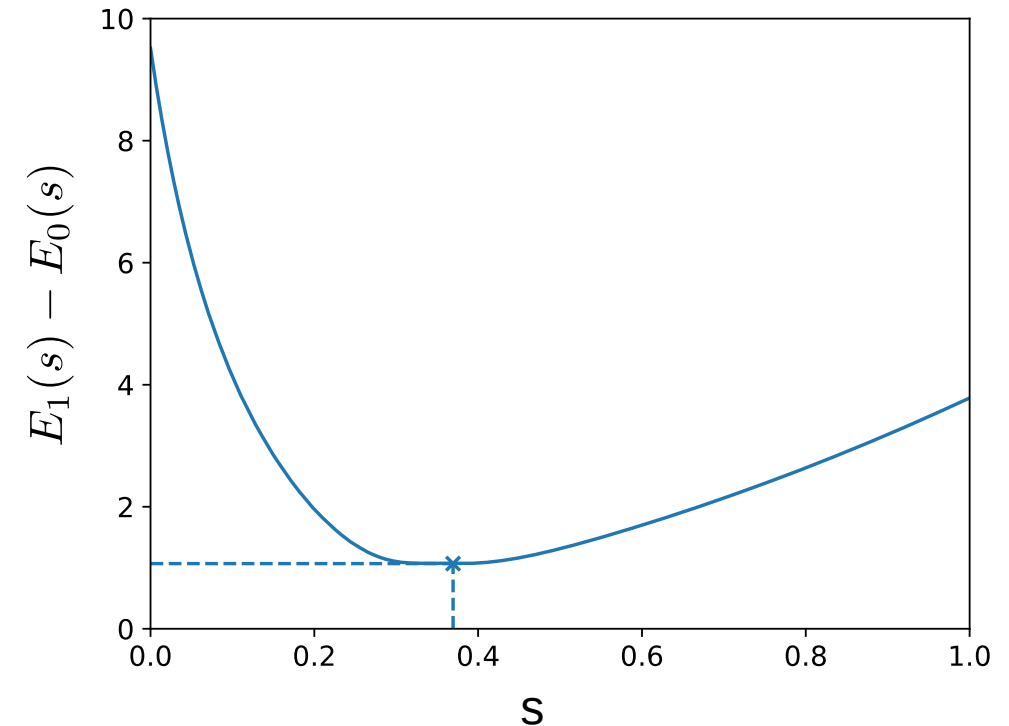
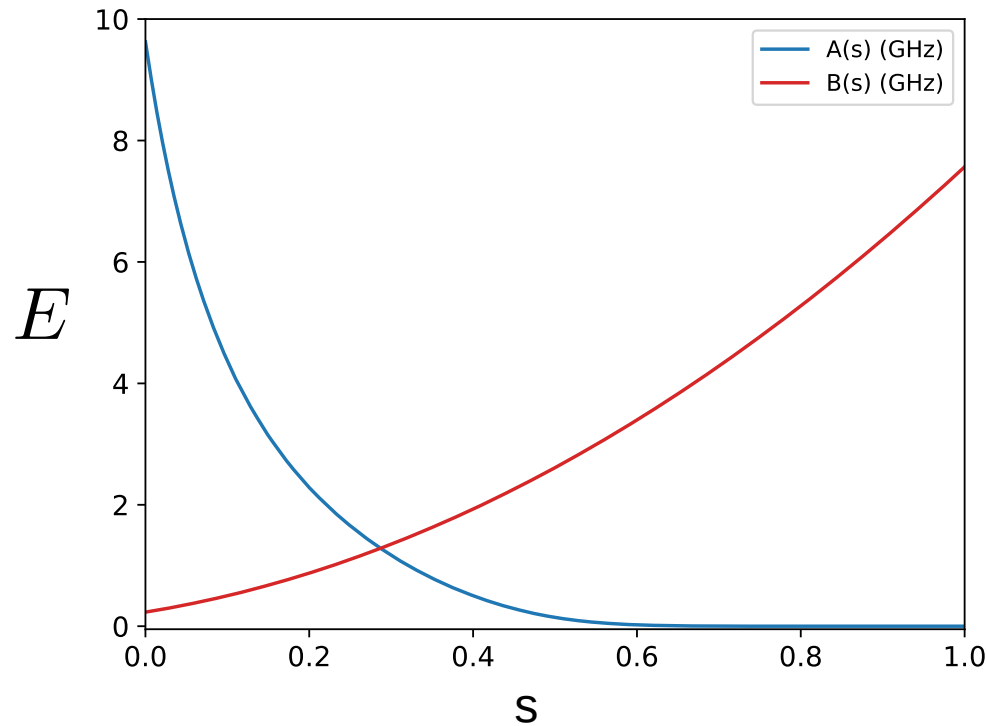
$$H(s) |\psi_i(s)\rangle = E_i(s) |\psi_i(s)\rangle, \text{ with } E_0(s) \leq E_1(s) \leq E_2(s) \leq \dots \leq E_p(s)$$

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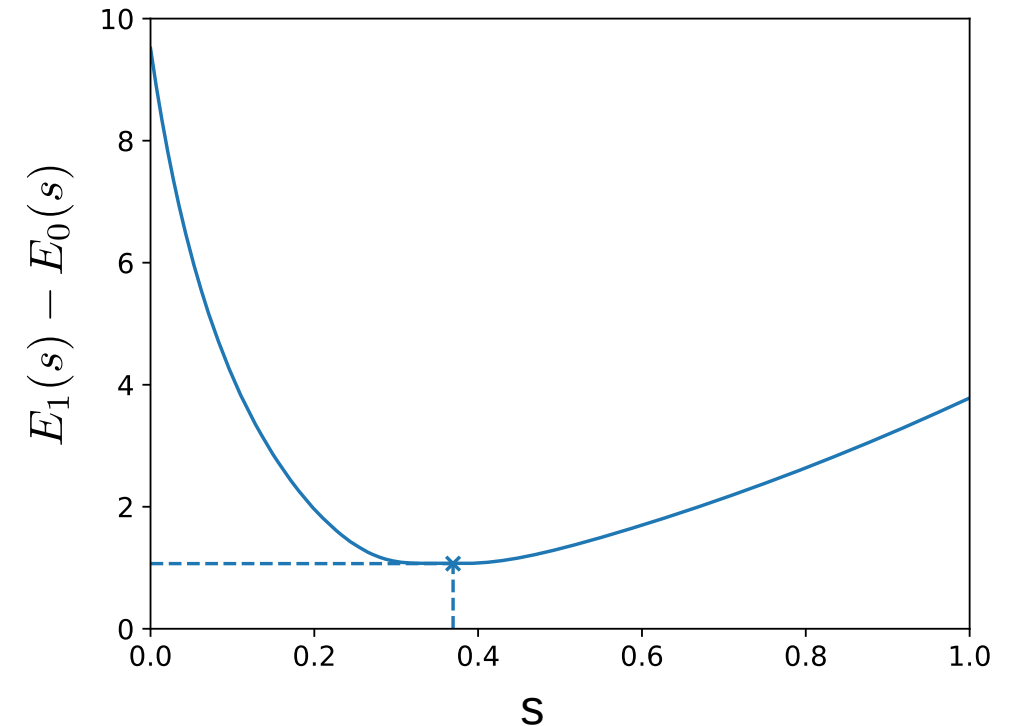
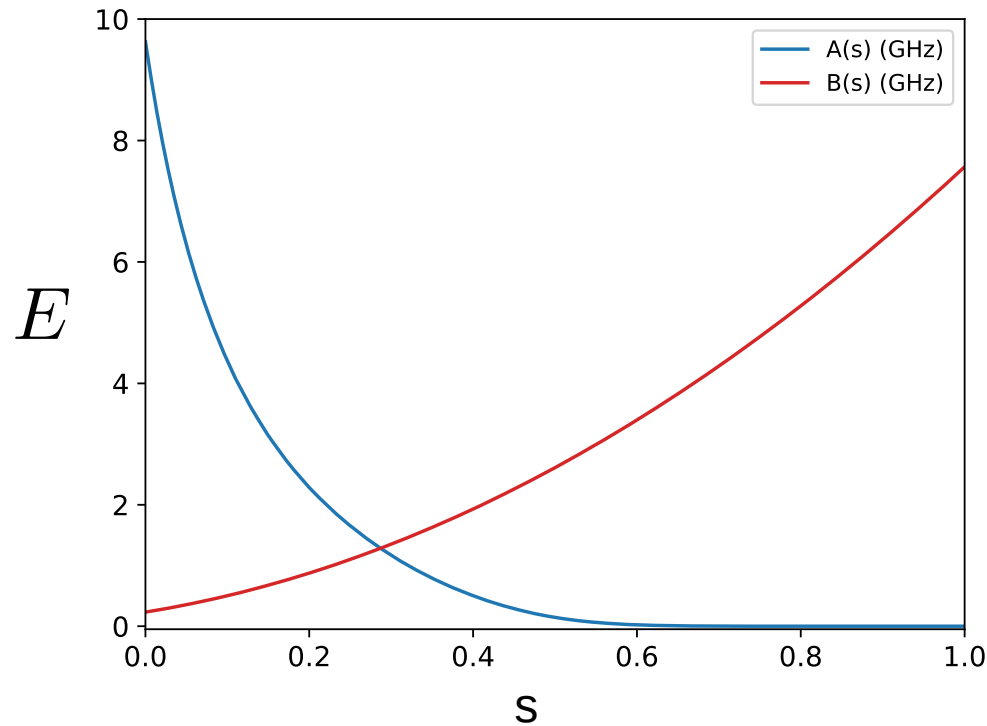
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$$g_{\min} = \min_{0 \leq s \leq 1} E_1(s) - E_0(s)$$

$$T \propto O(1/g_{\min}^2)$$



Contribution #1

Performance evaluation of Quantum Annealers

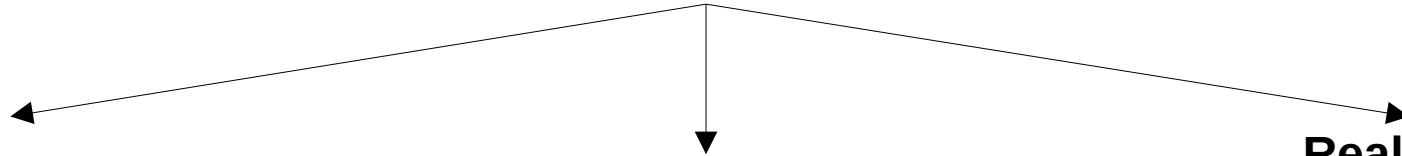


Generated with openai

2- Performance evaluation of Quantum Annealers



Application benchmark instances



Random instances :

- + Generate large sets of instances
- Relevance of the set can be debated
- Phase transition issue for some instances

Crafted instances :

- + Easy to find instances hard for classical solvers
- + Create instances with planted solutions
- Some classical methods may take advantage of the structure of the instance

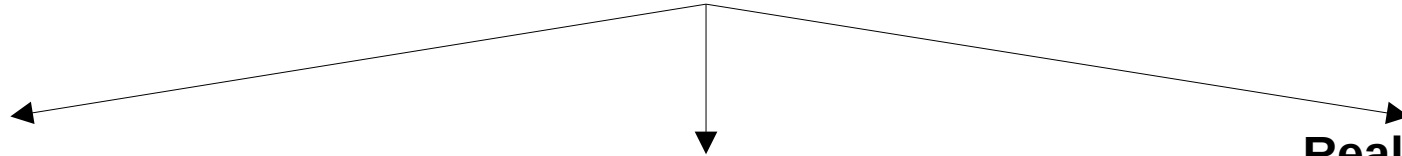
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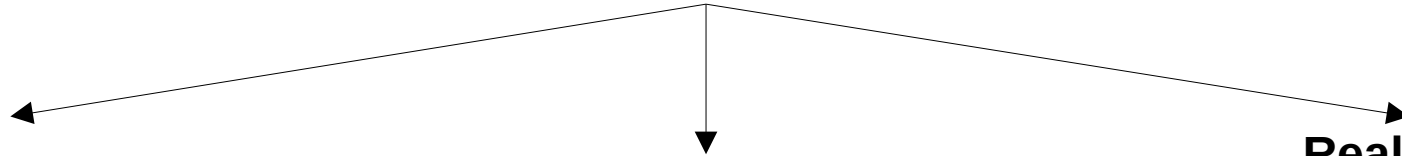
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2023 [LCG ⁺ 24]	3-regular graph	4-320	Adv4.1	QPU chip-subgraph
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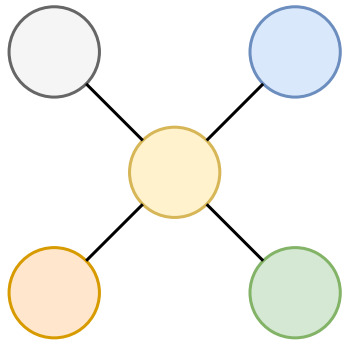
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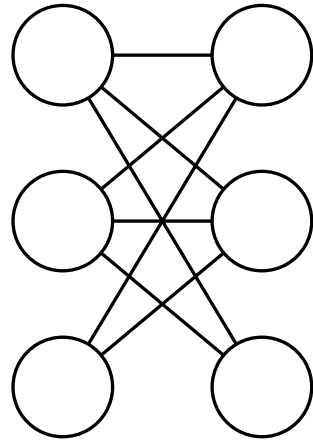
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2- Generation of instances

- Minor-embedding (Graph Minor Theory [\[RS95\]](#))



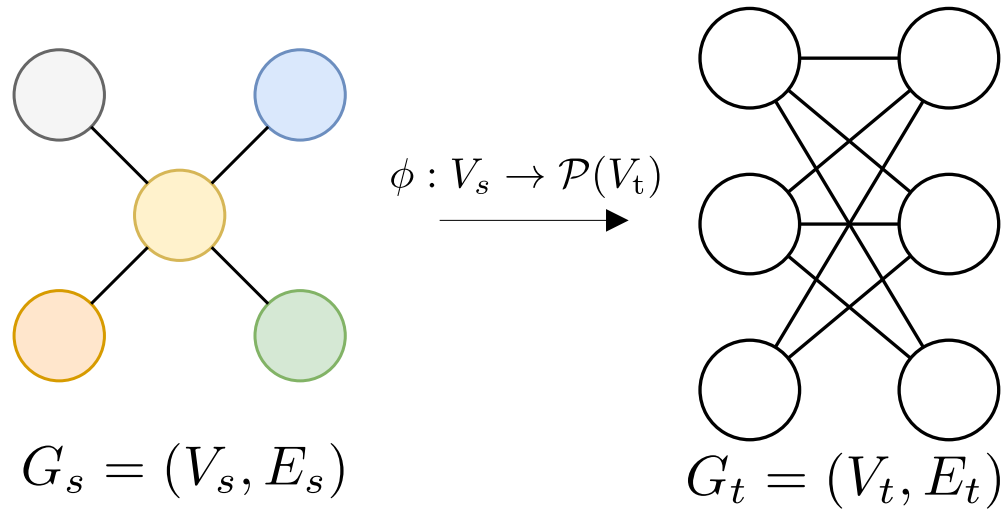
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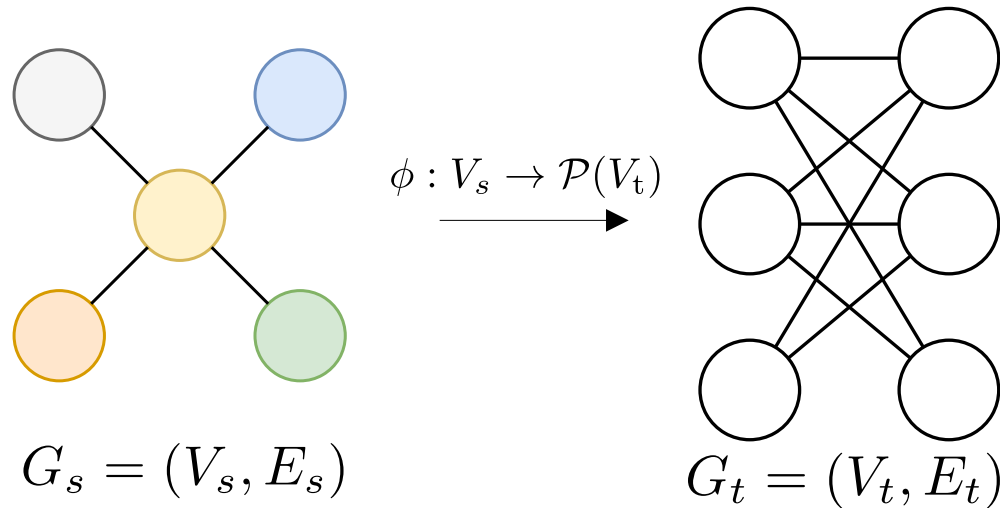
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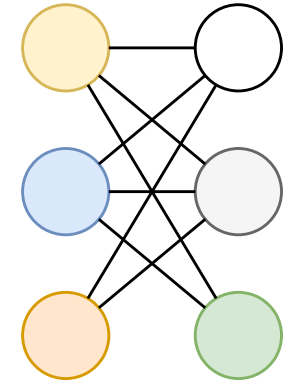


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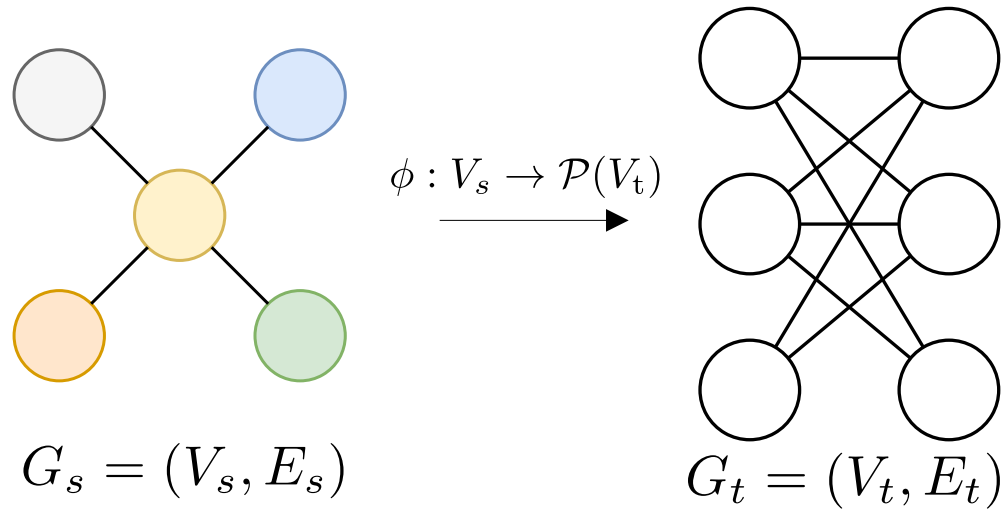


Rule 1:
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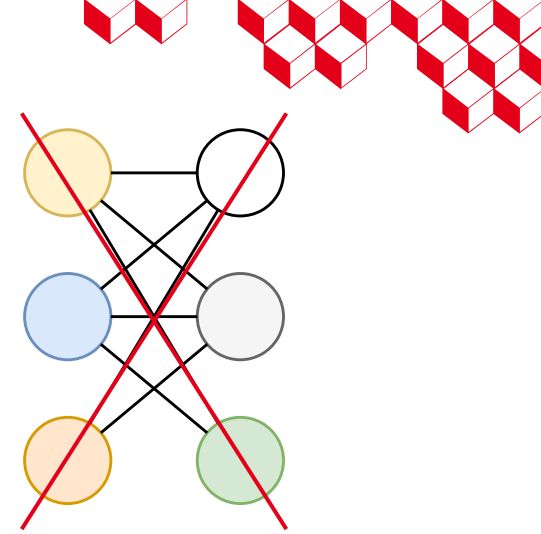


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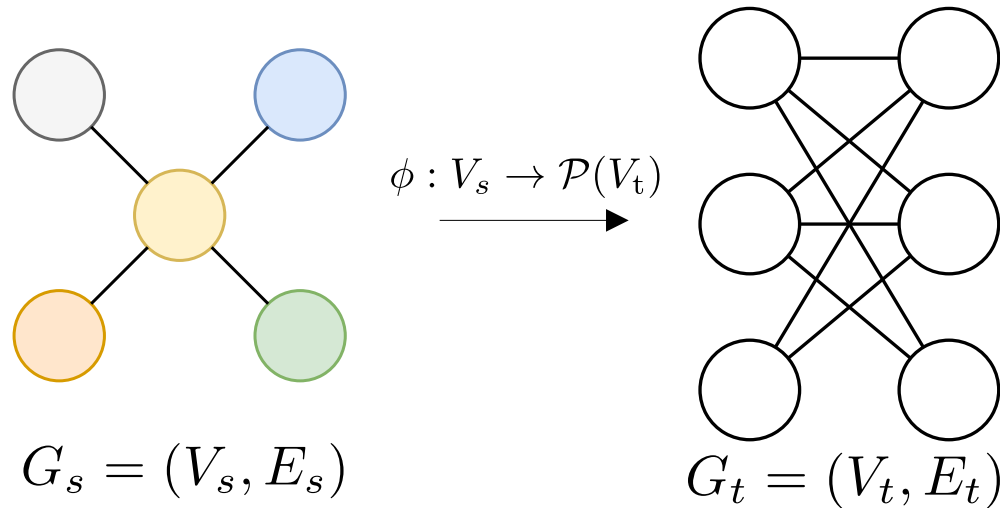


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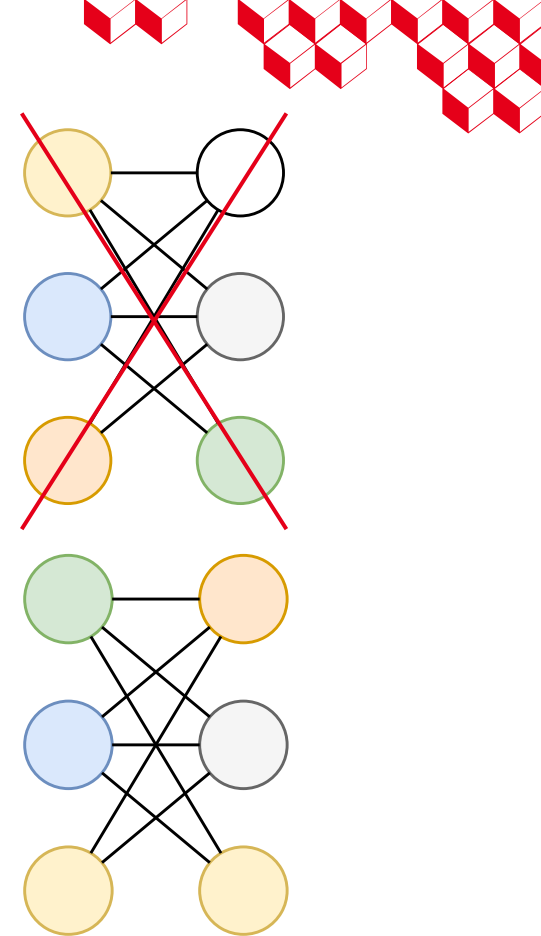
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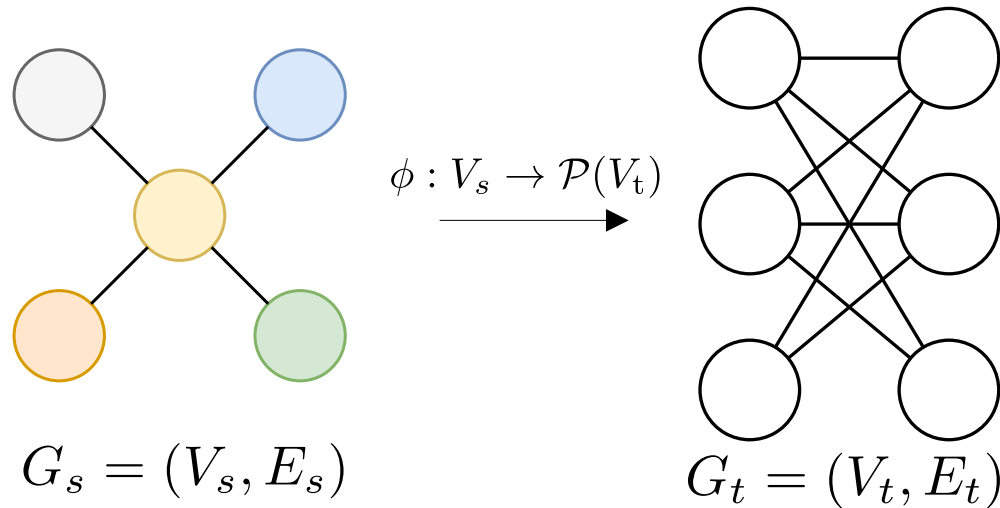
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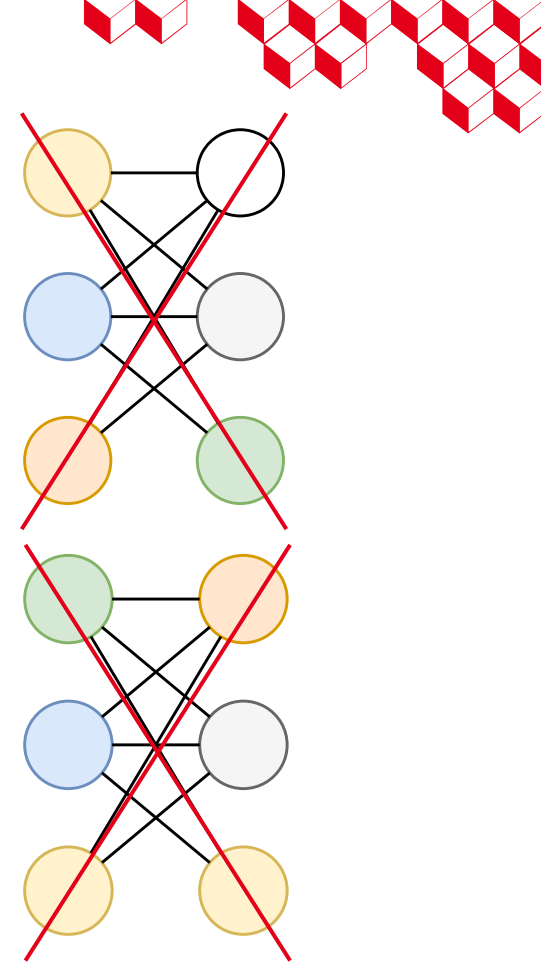
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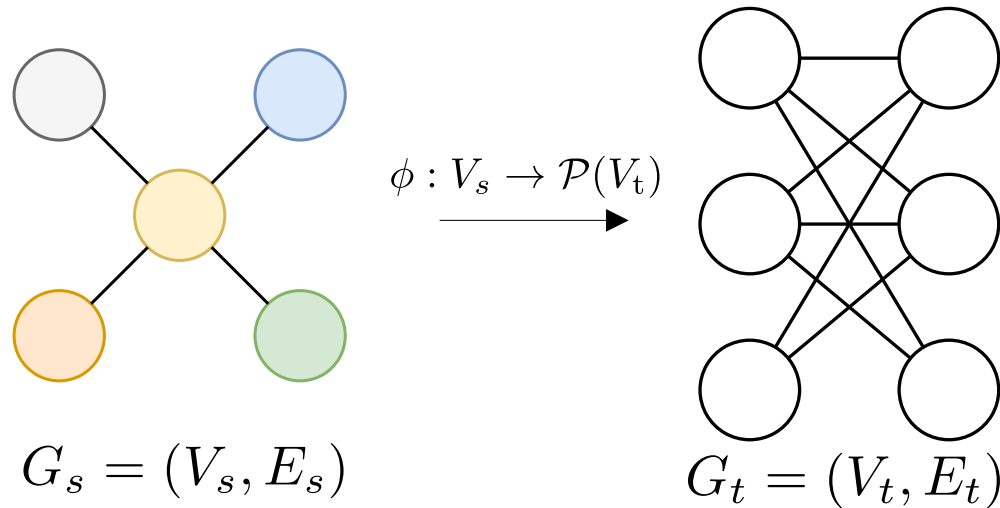
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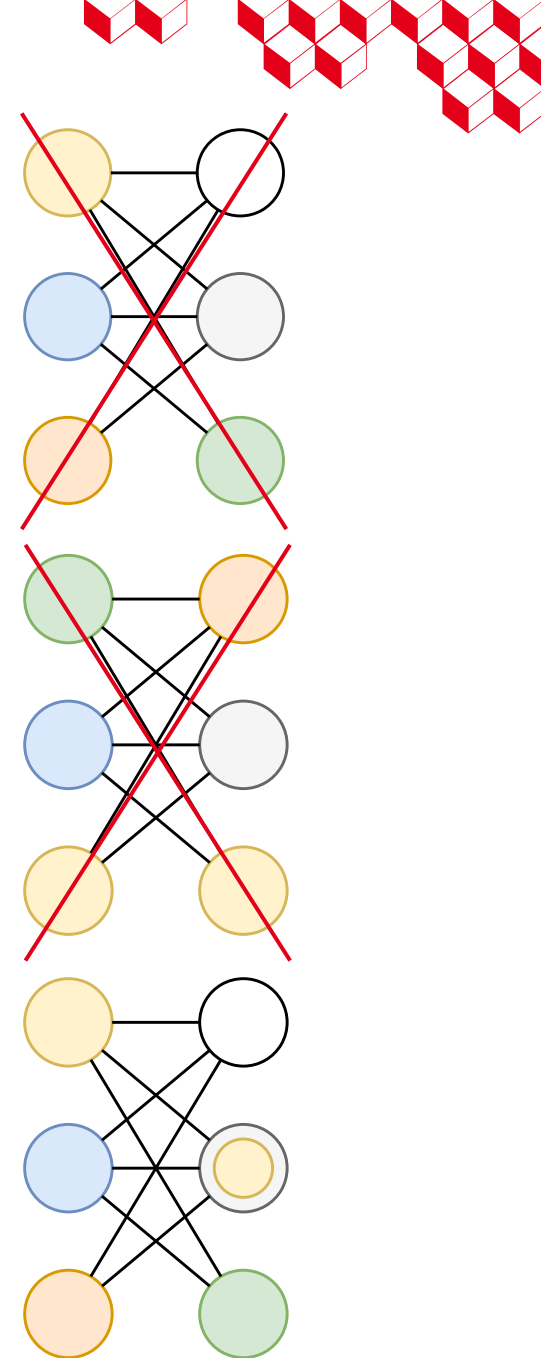
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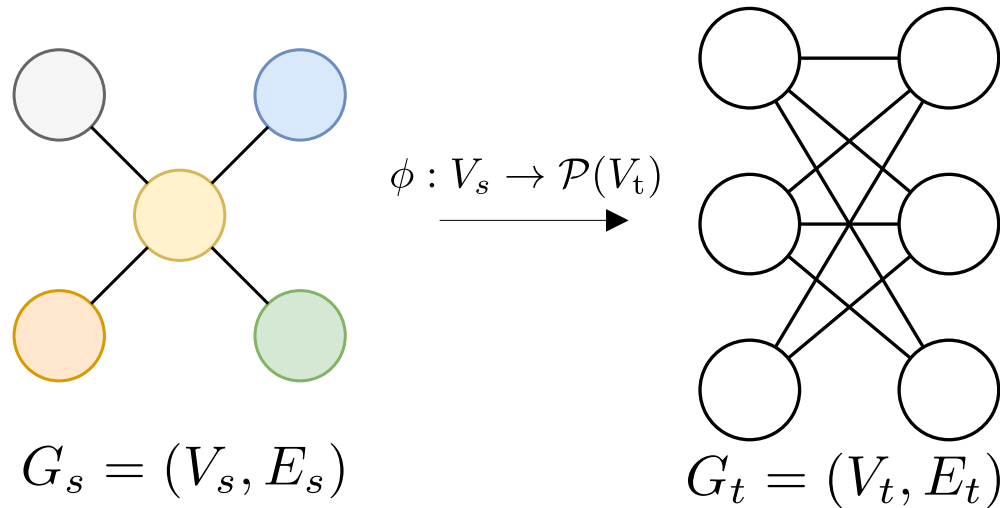
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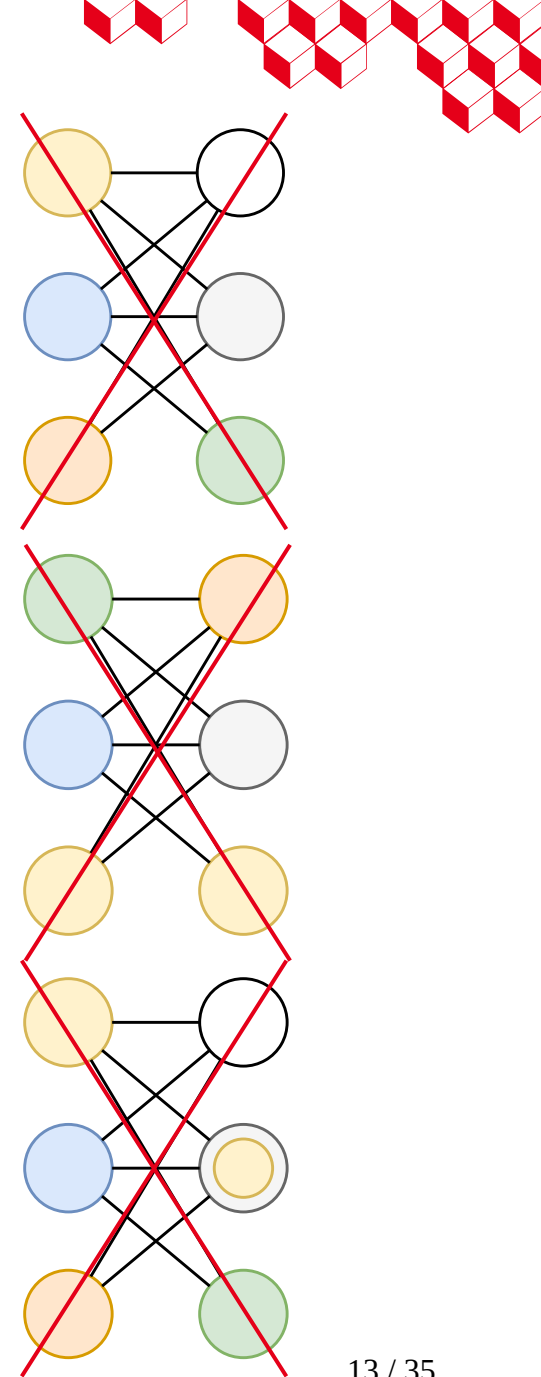
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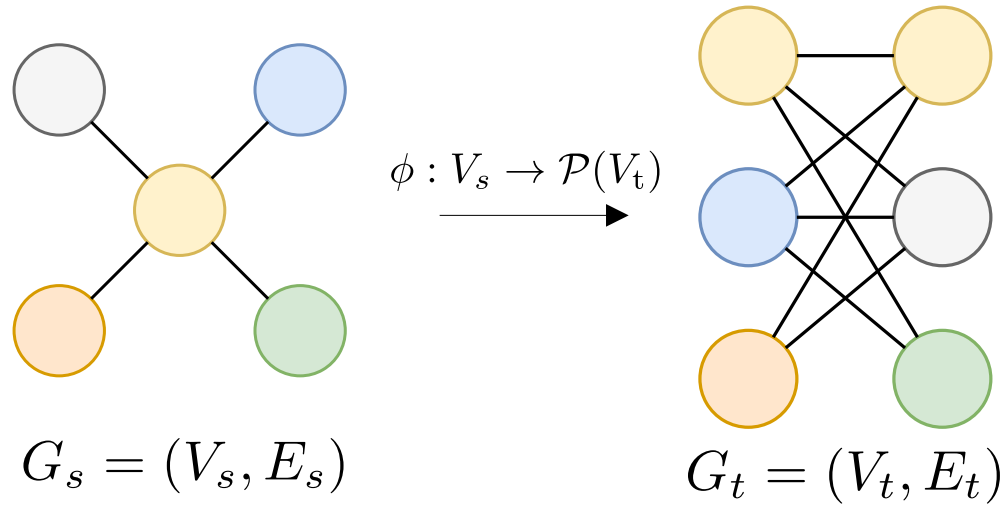
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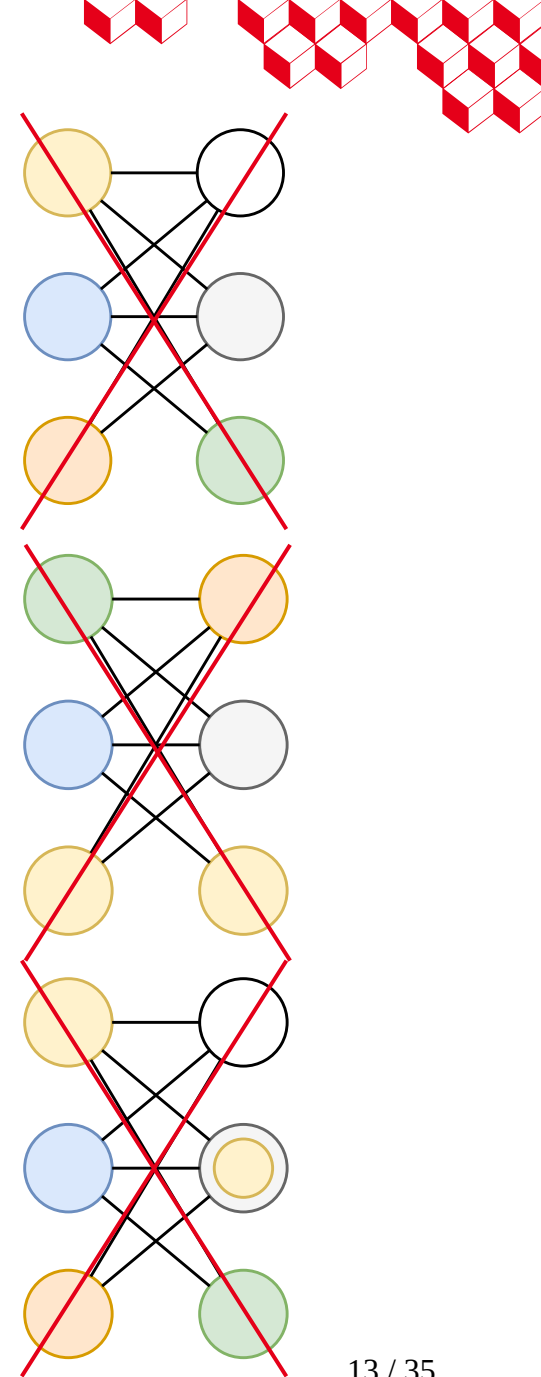
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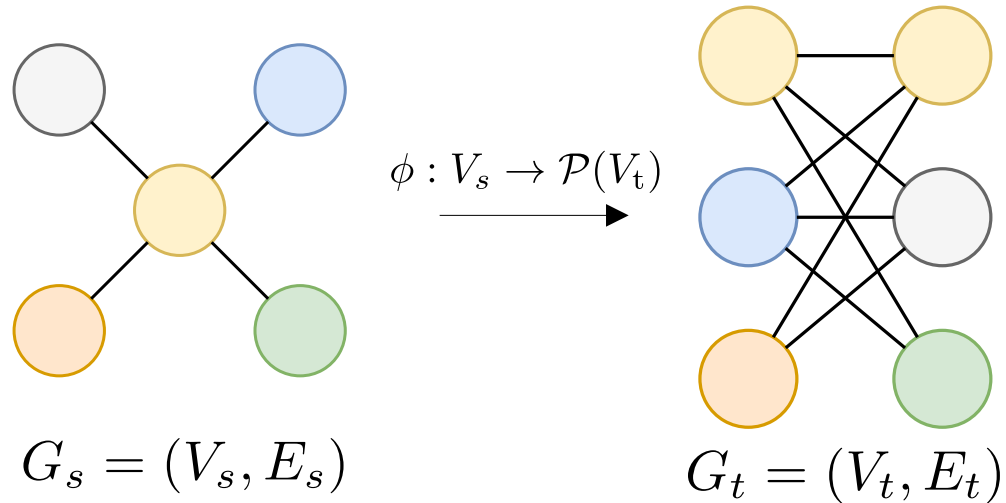
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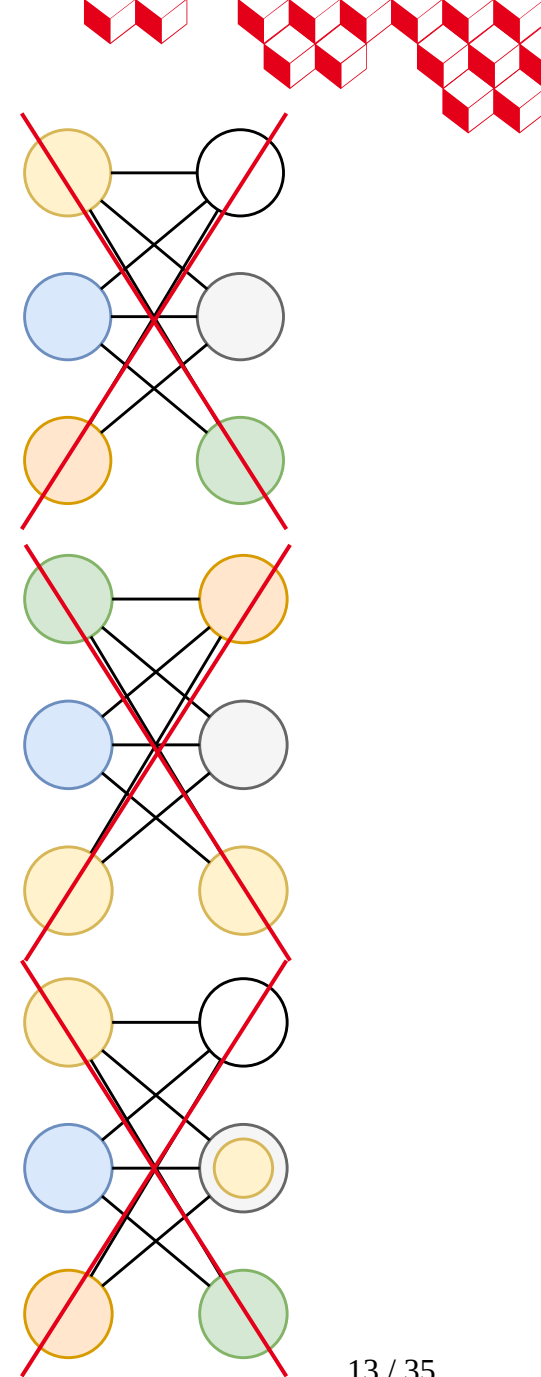


NP-Hard problem for arbitrary graphs

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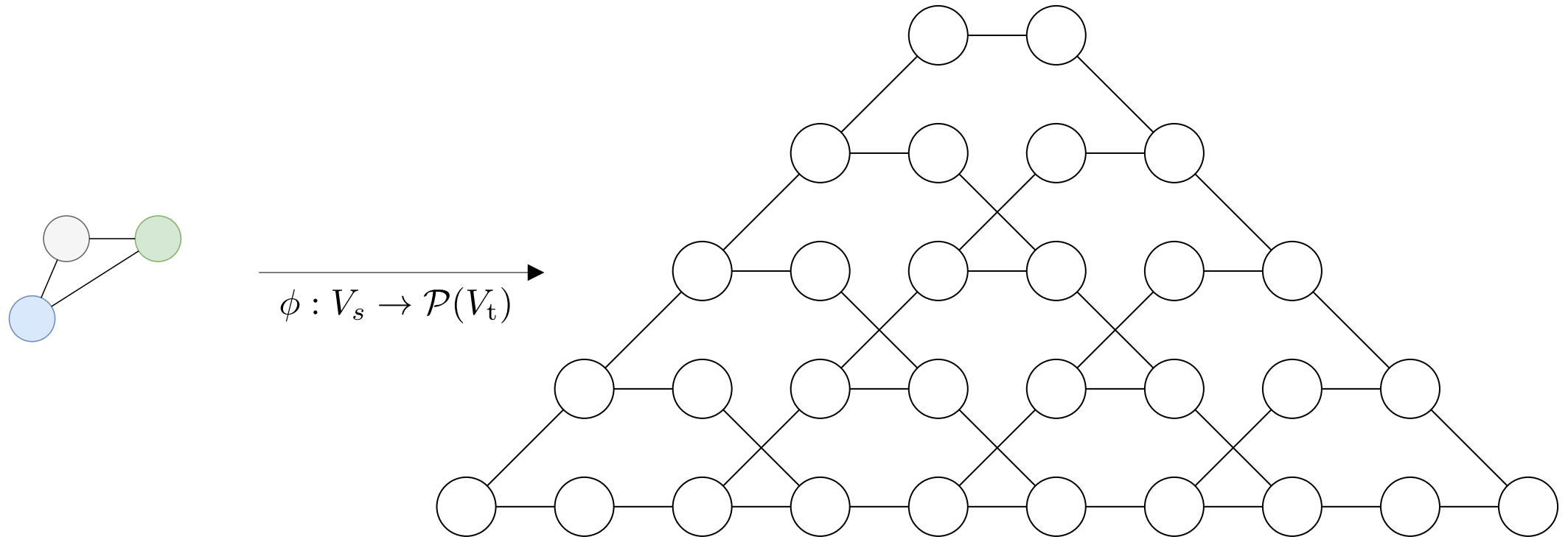
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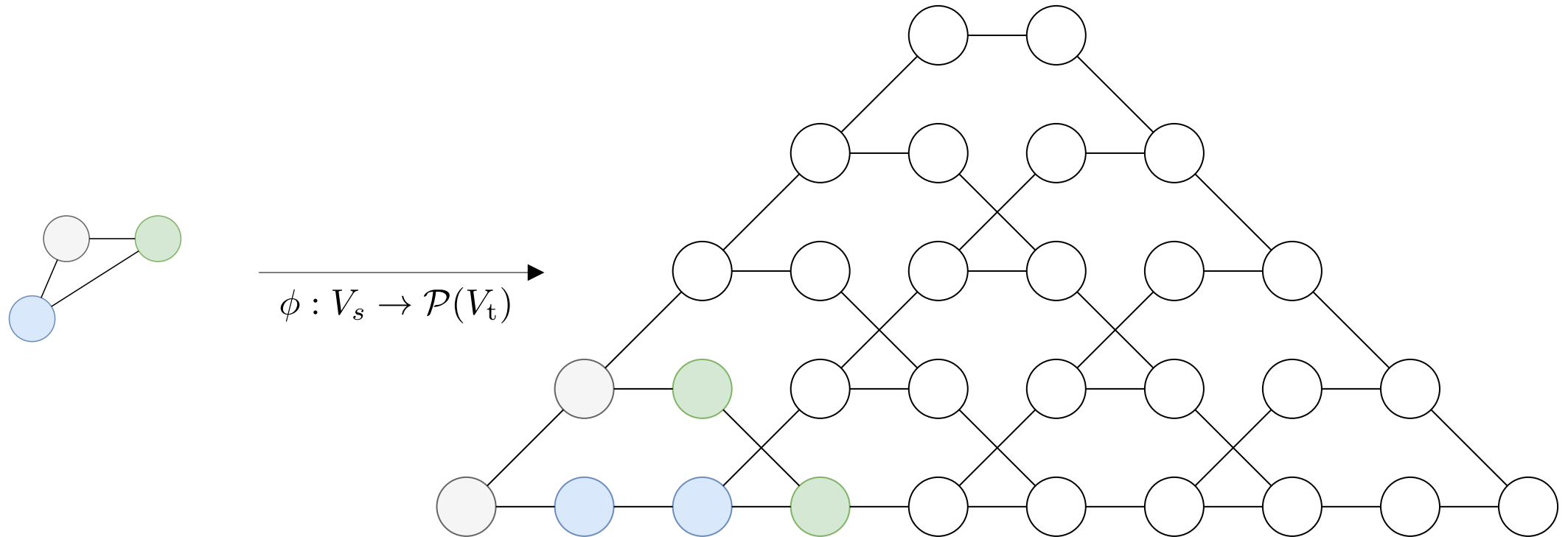
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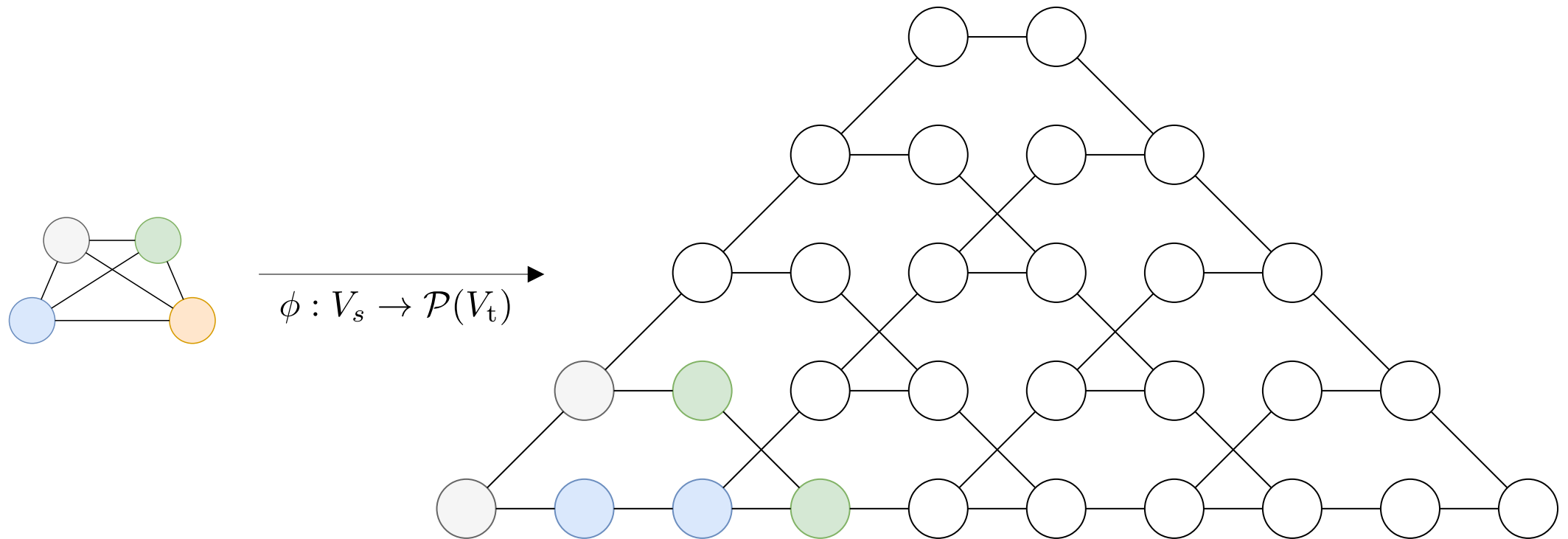
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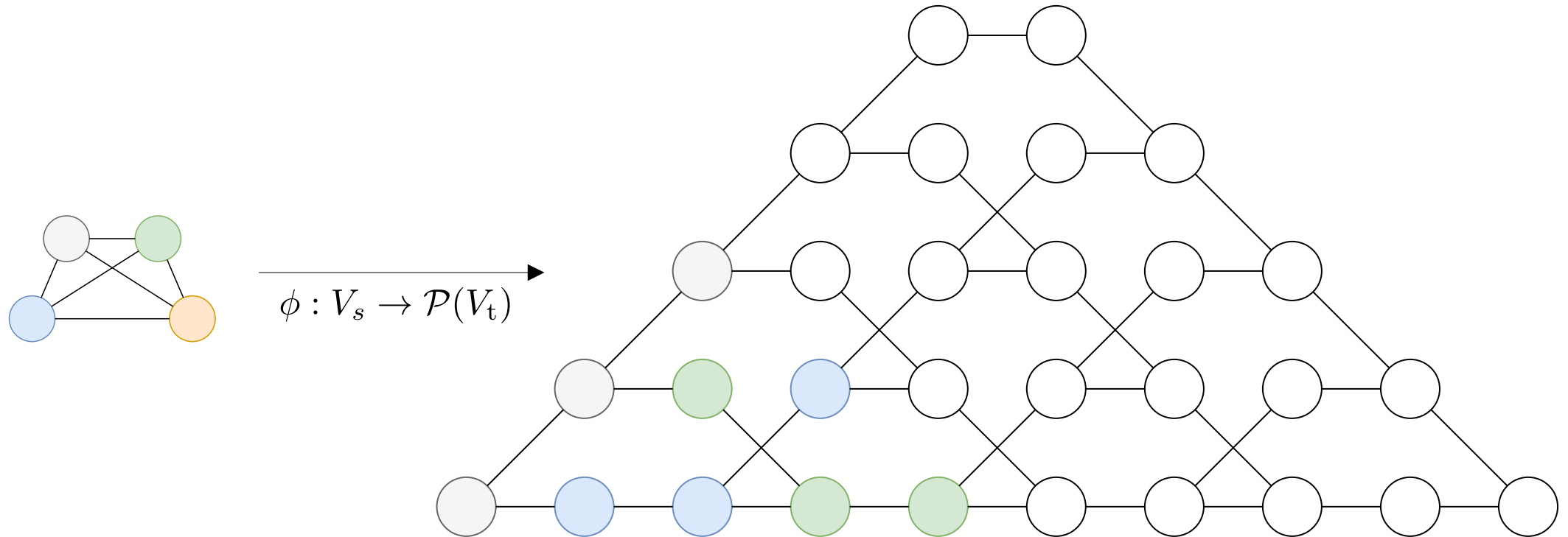
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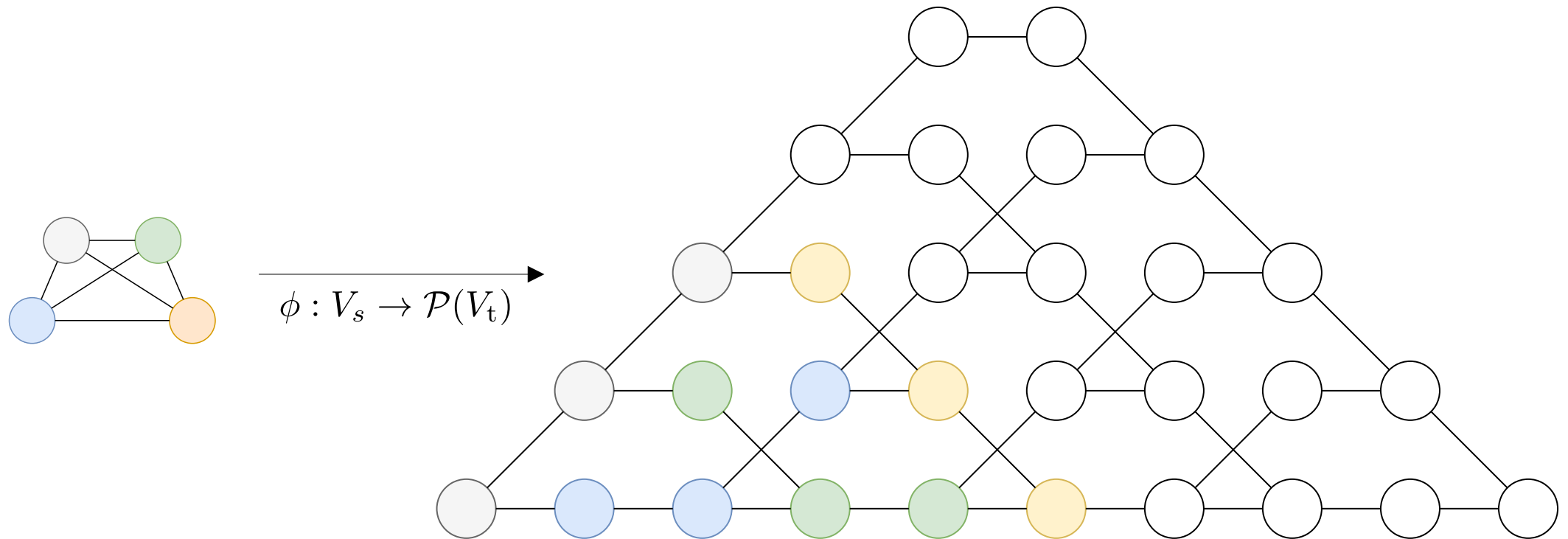
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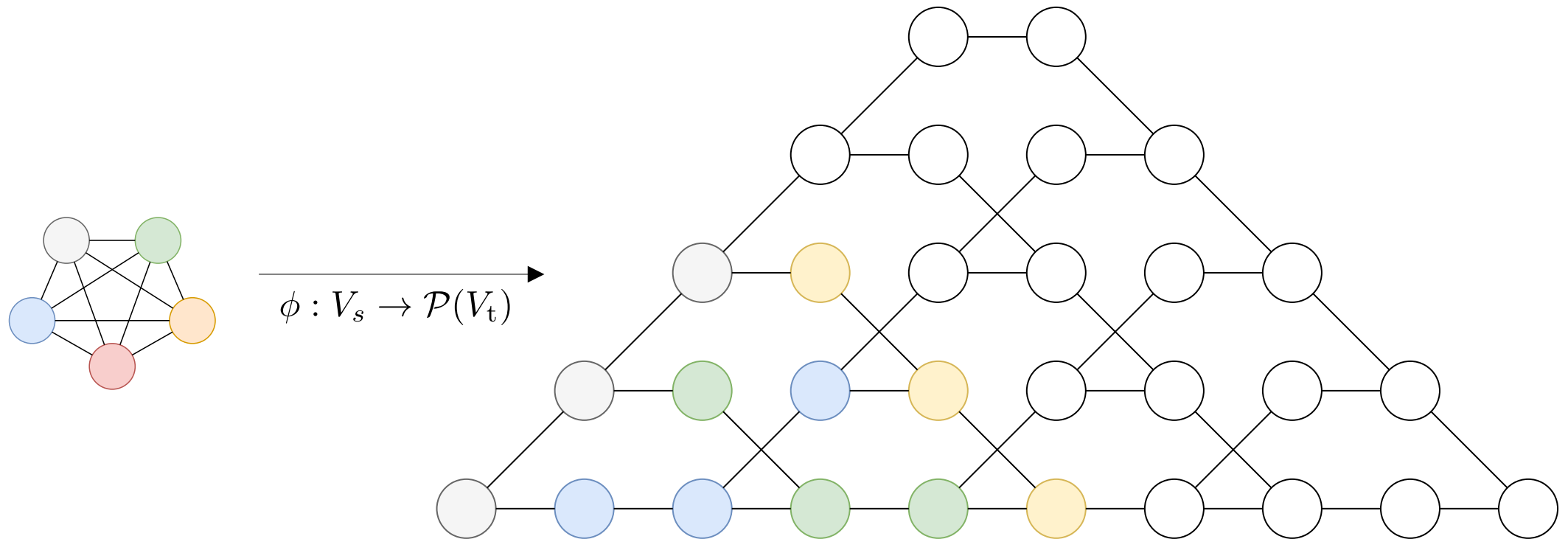
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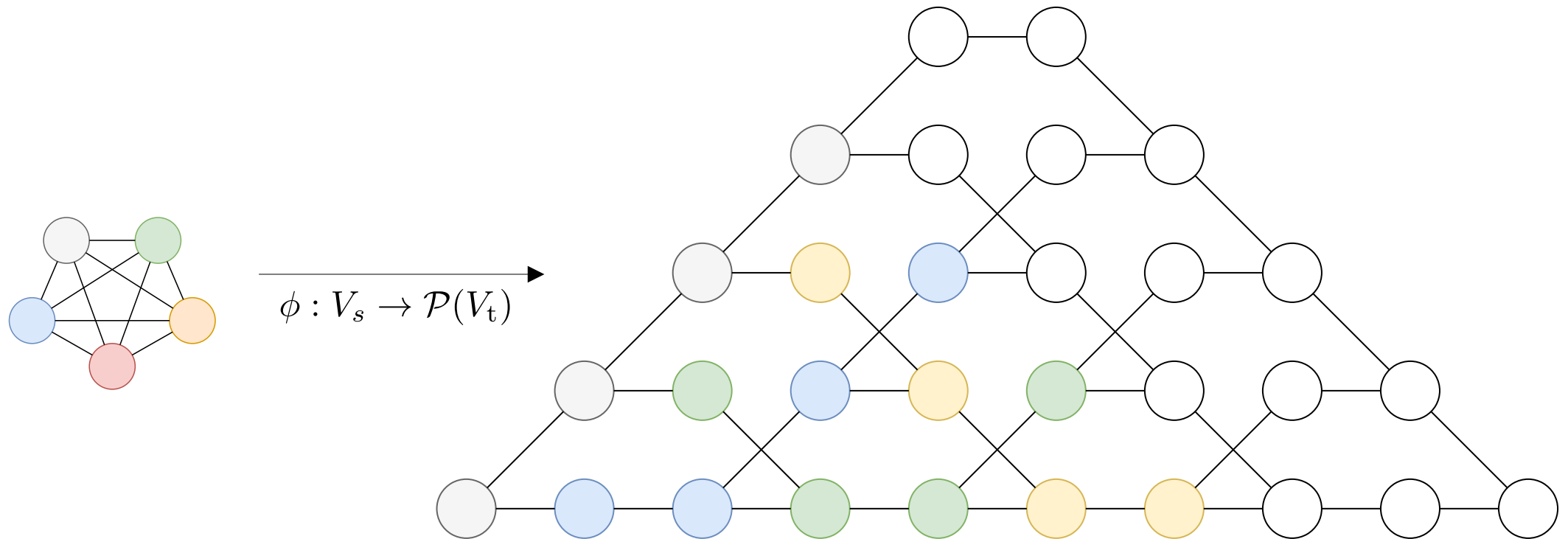
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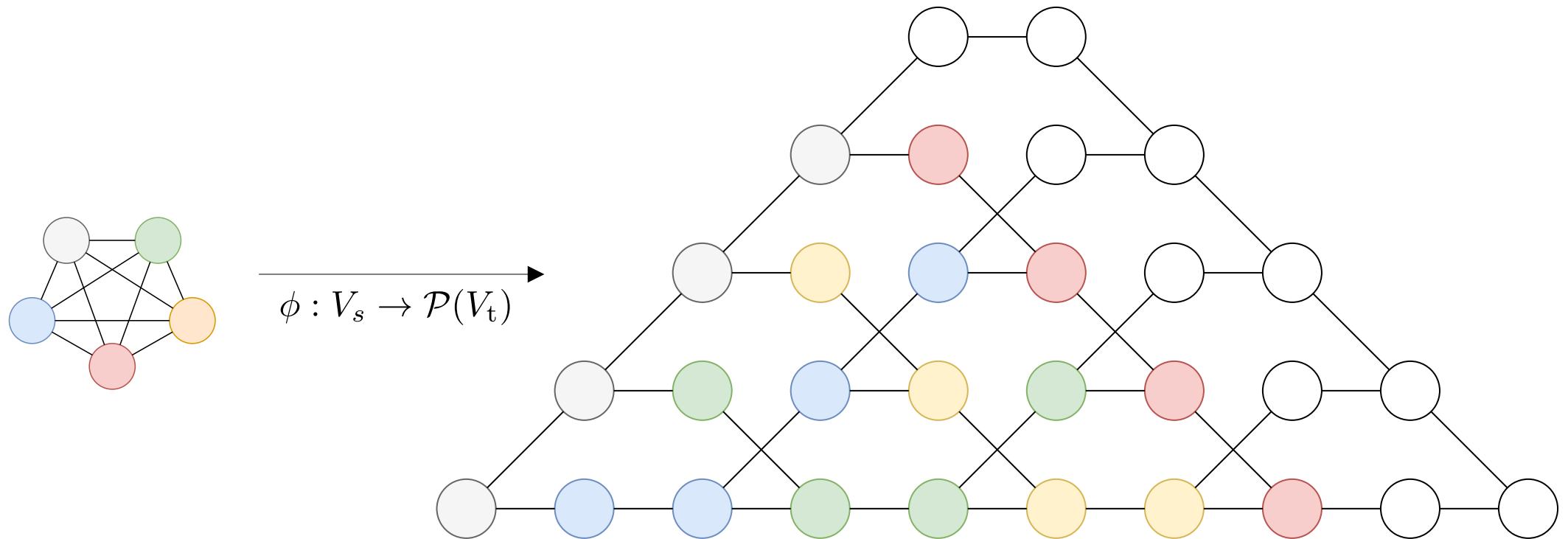
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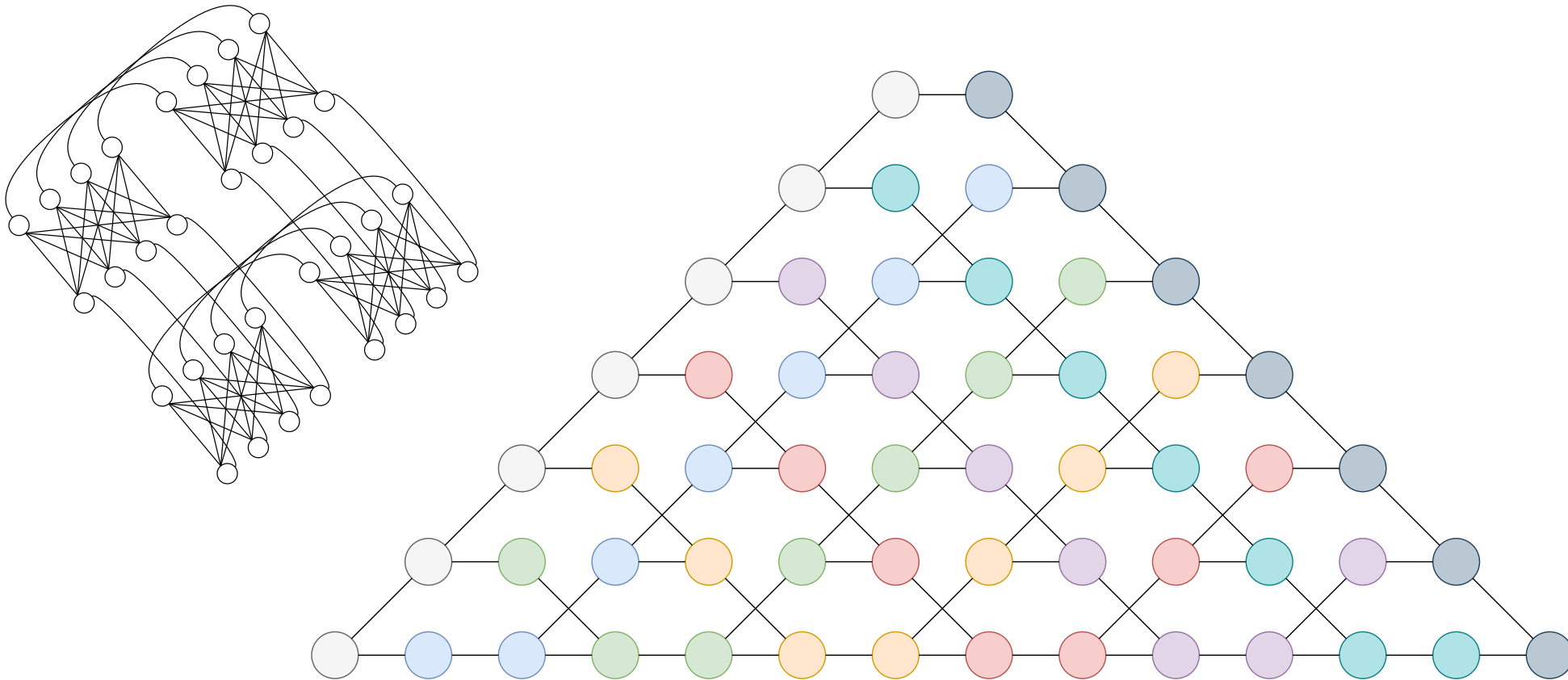
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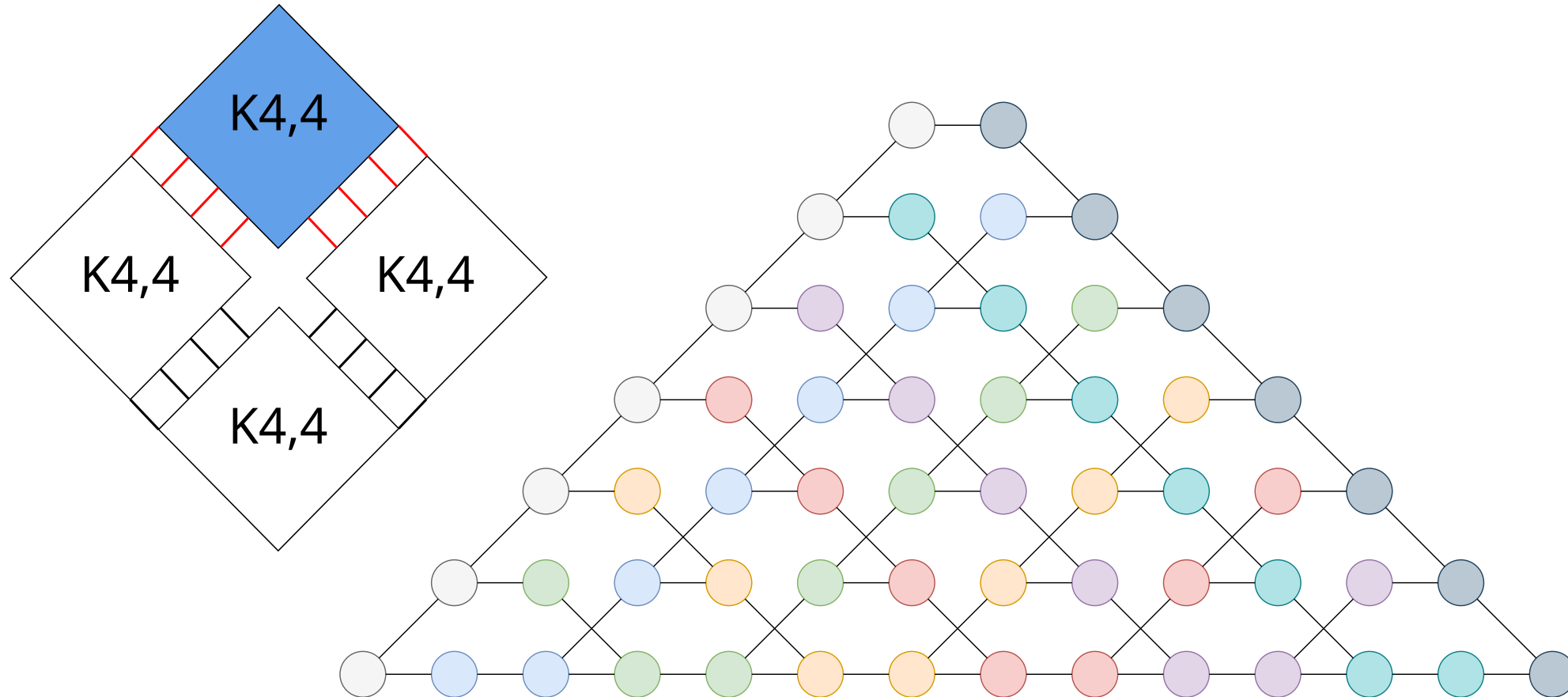
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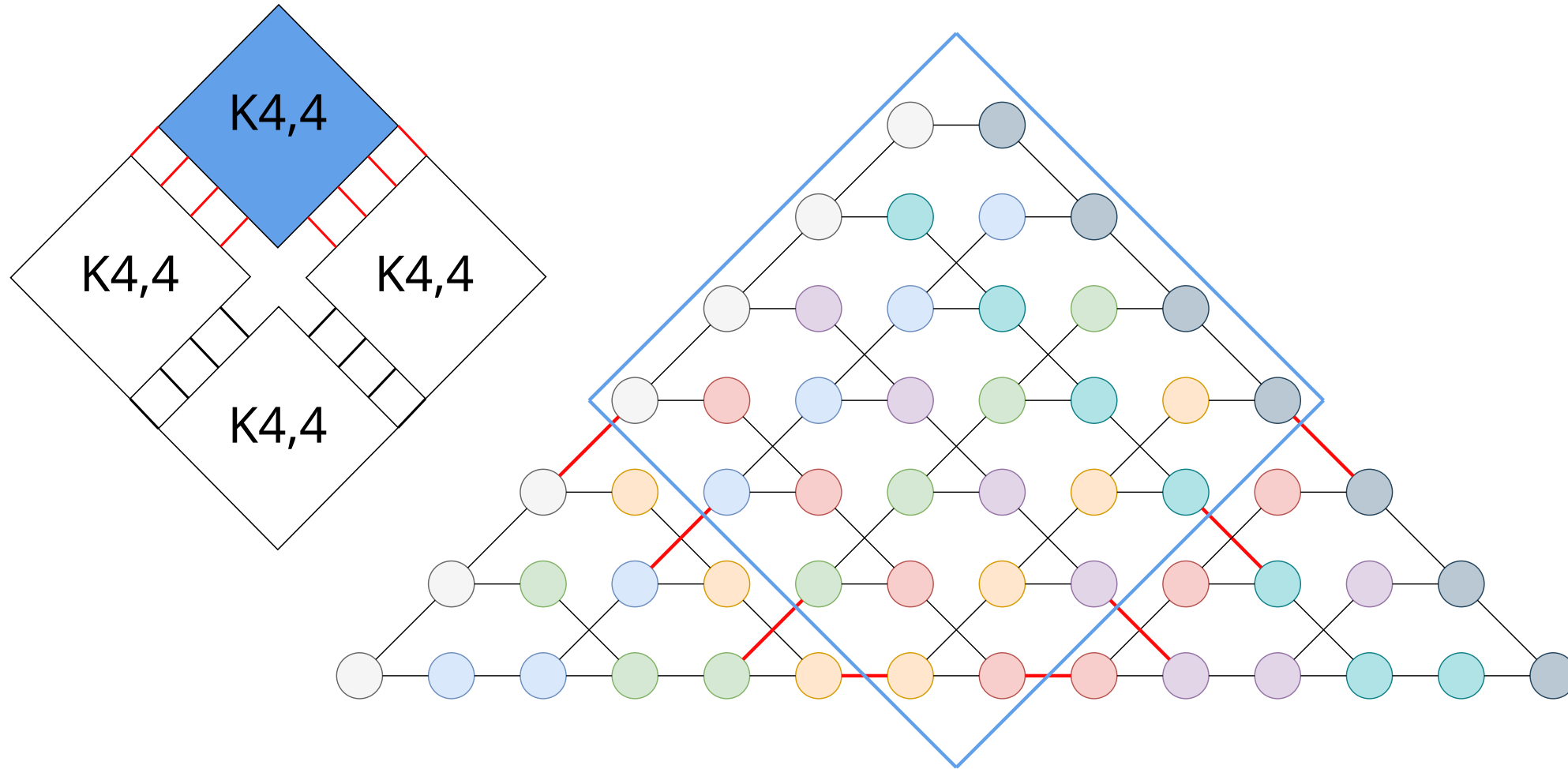
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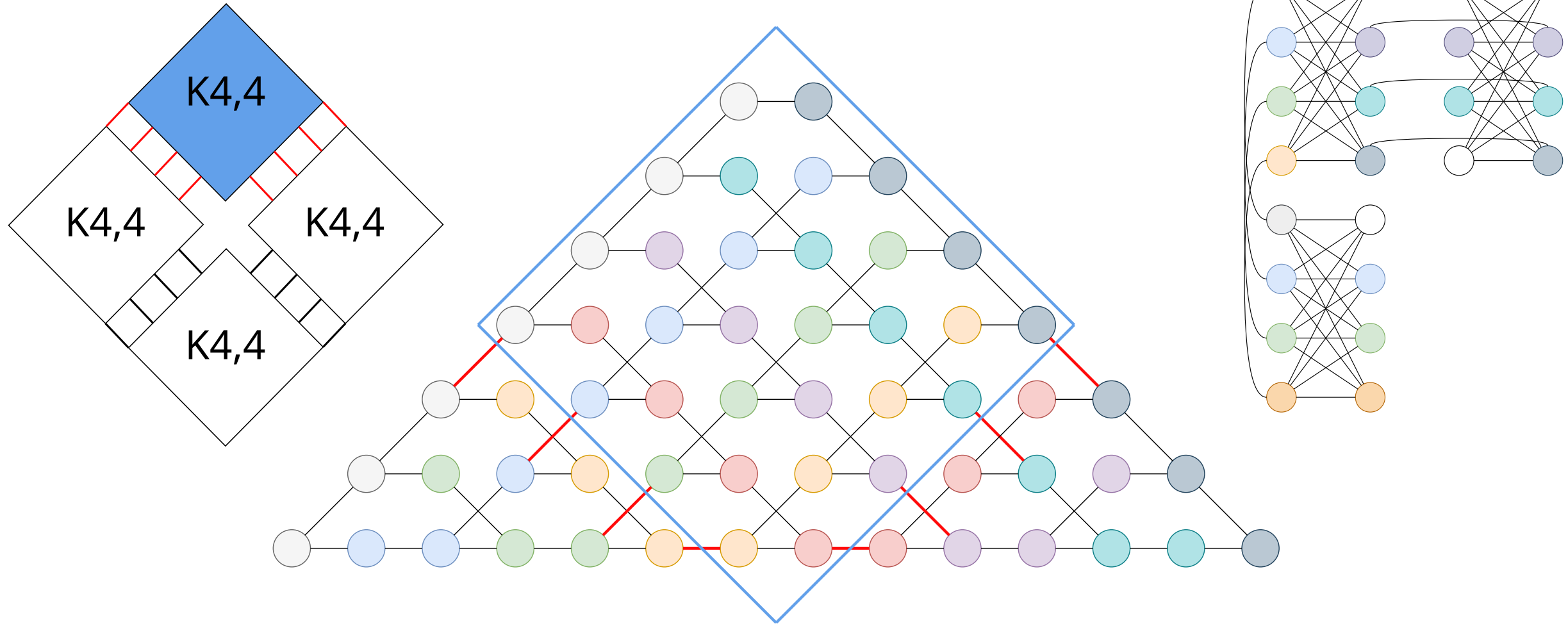
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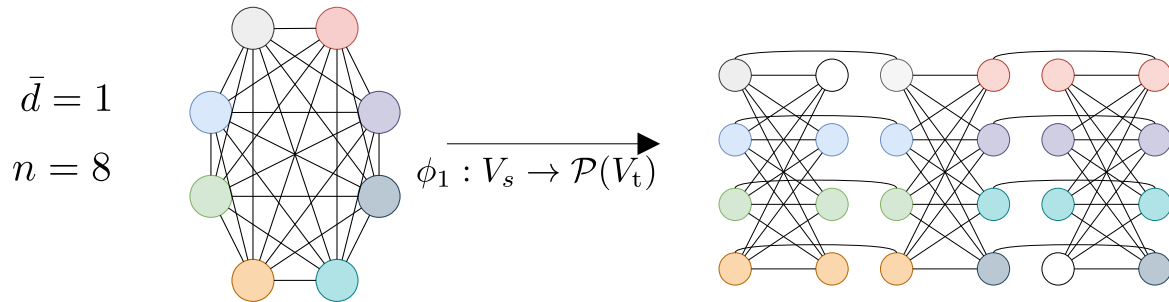
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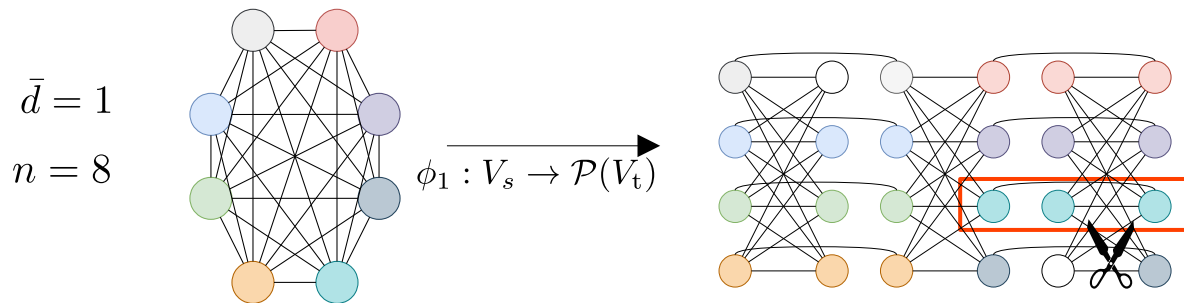
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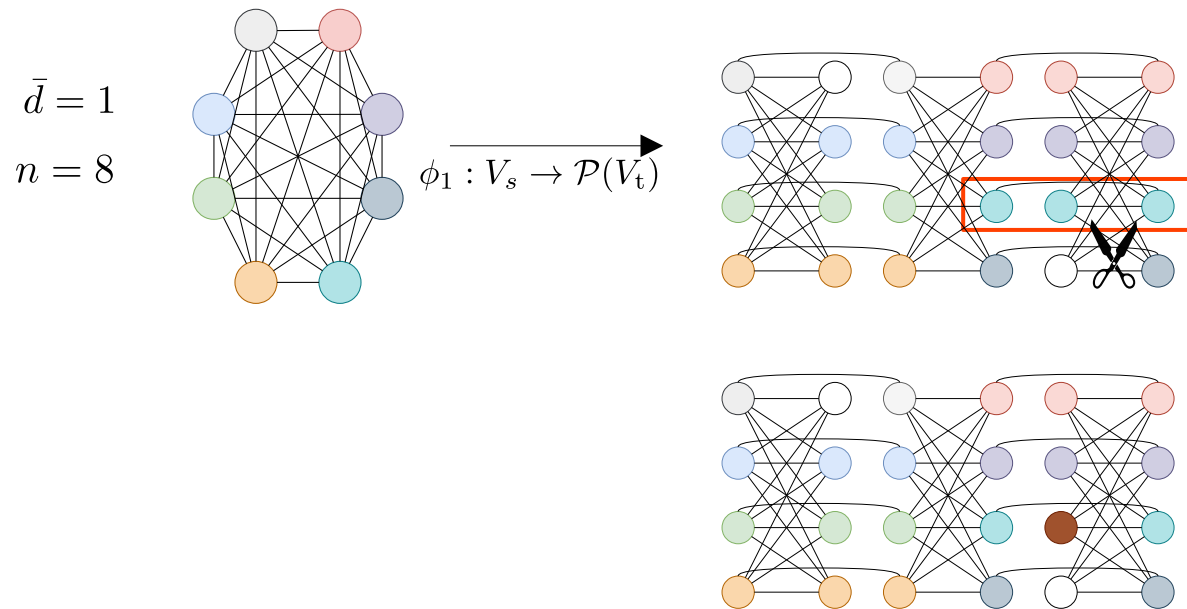
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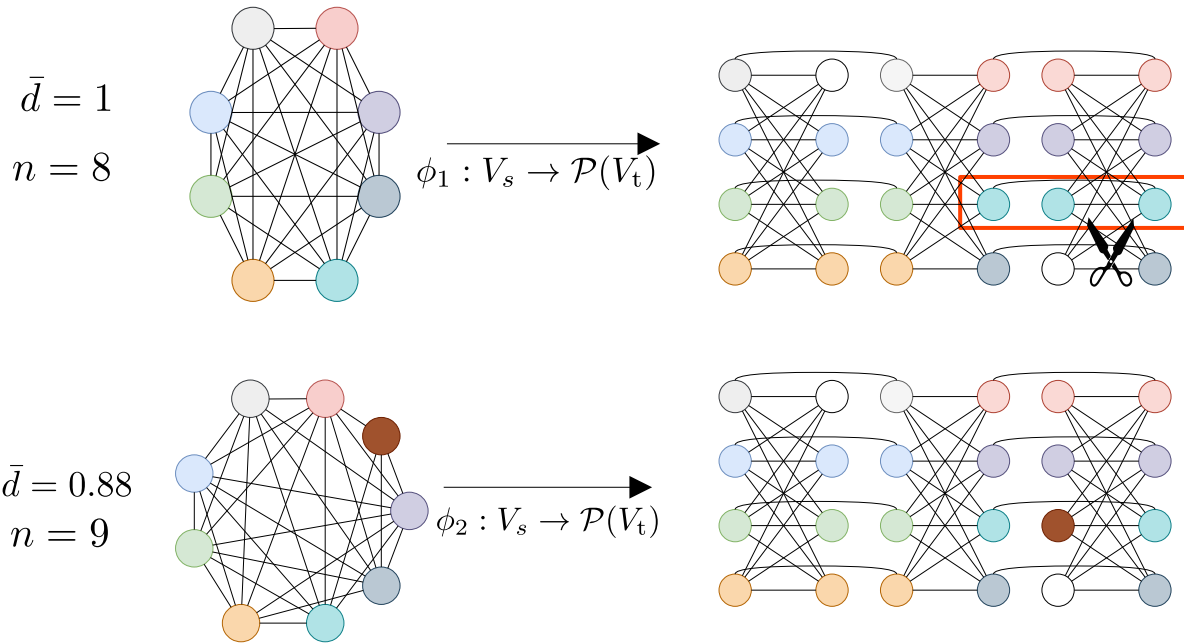
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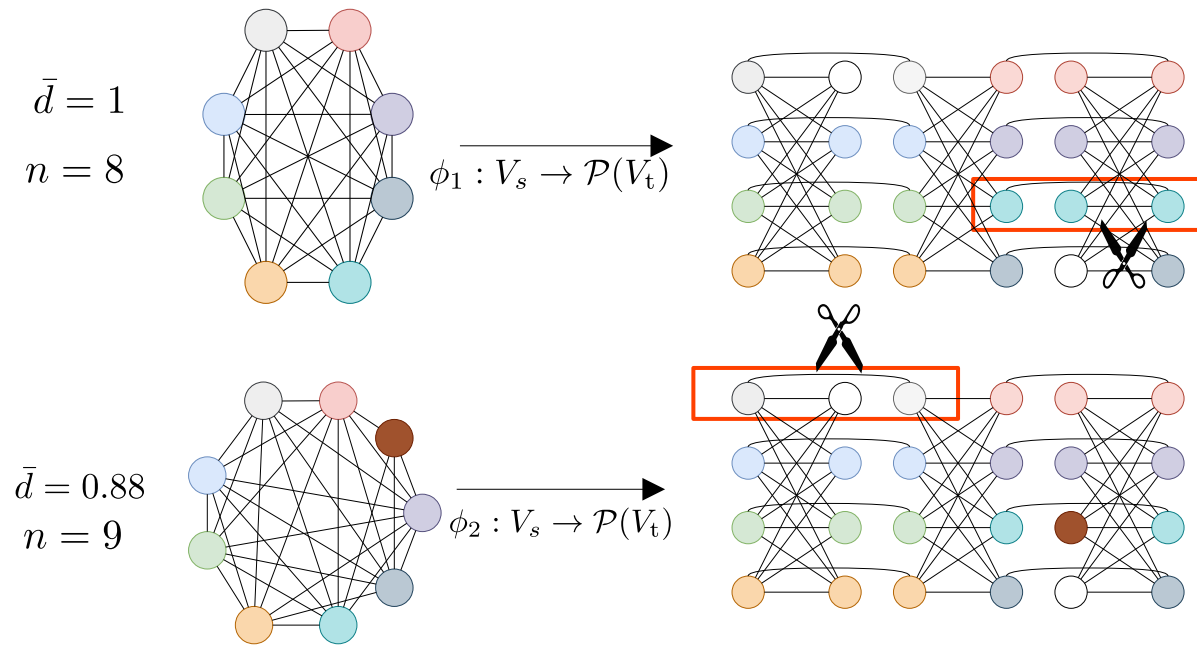
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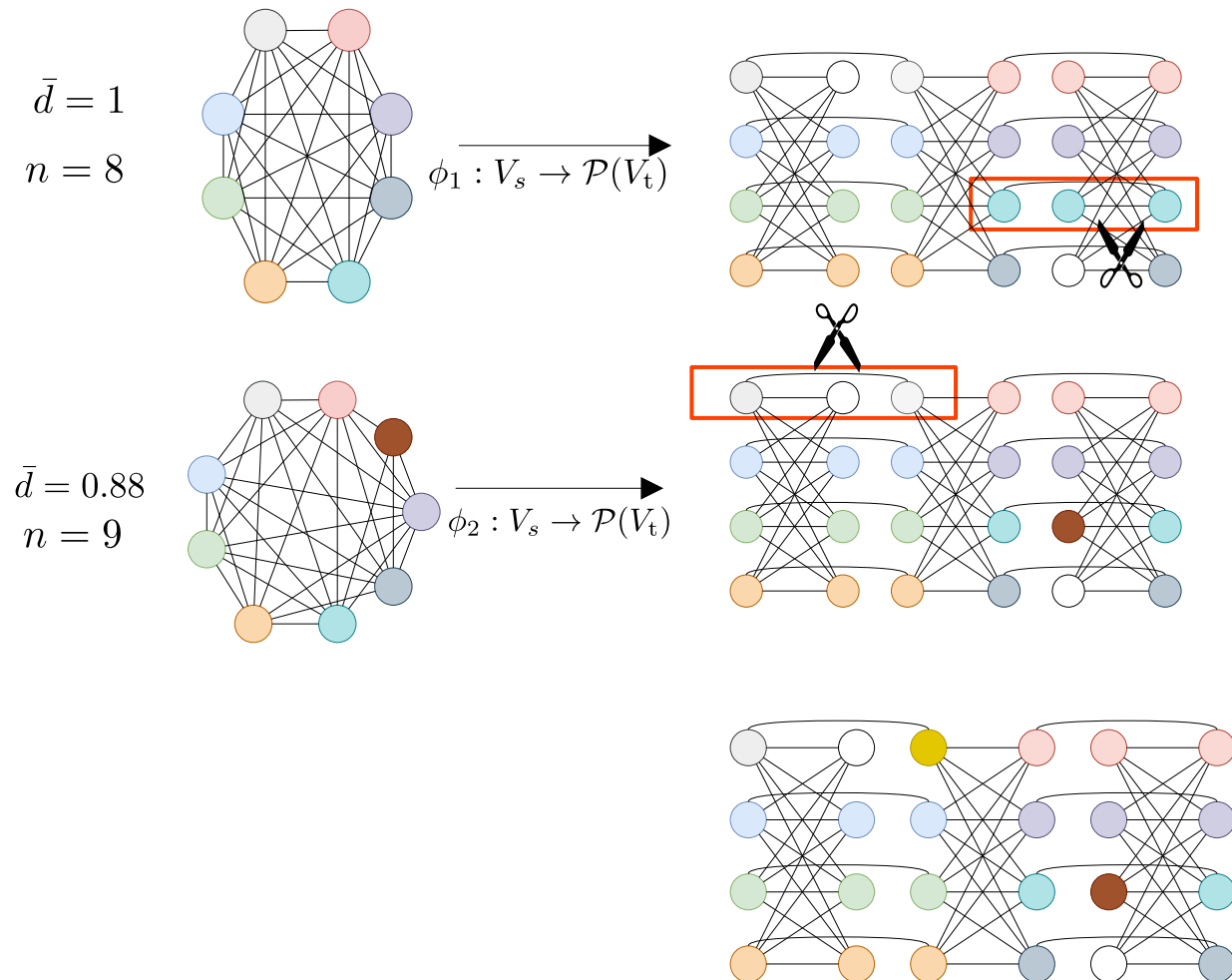
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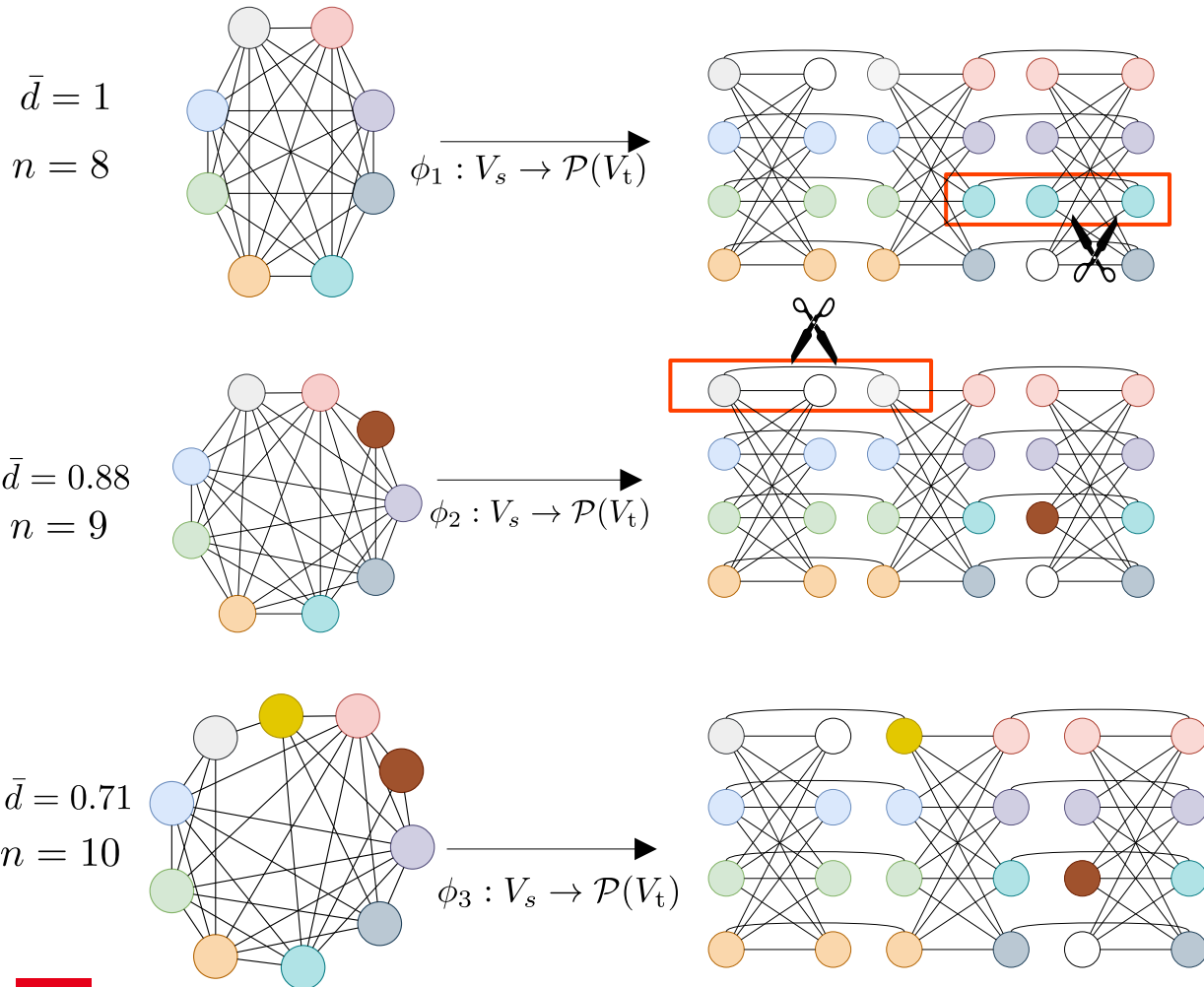
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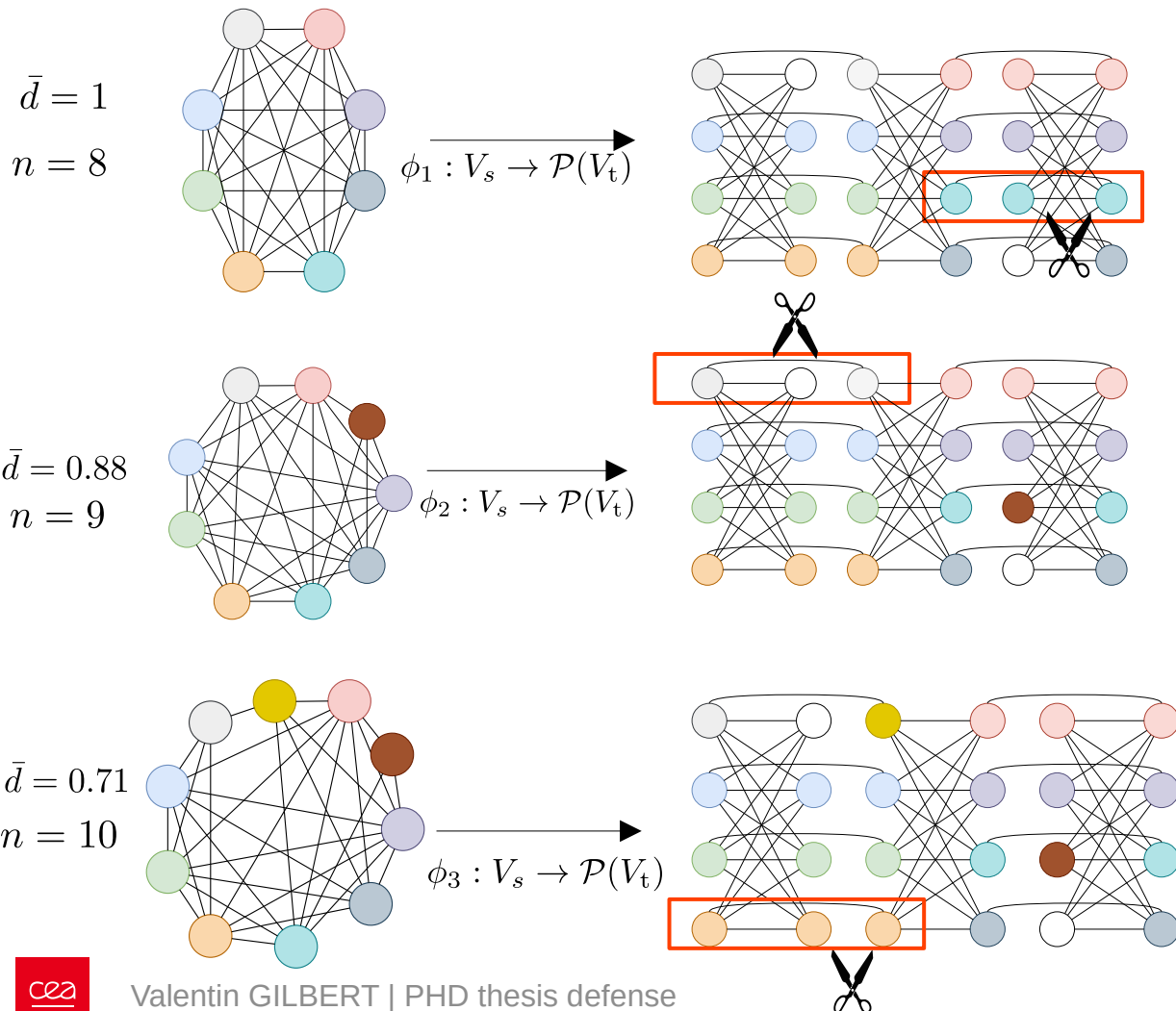
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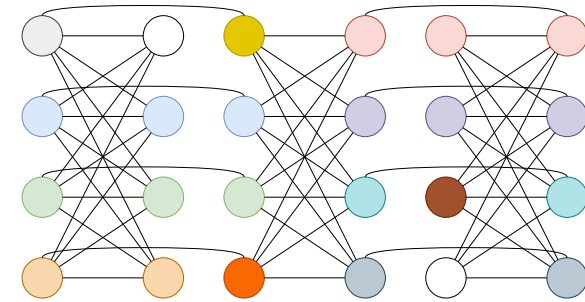
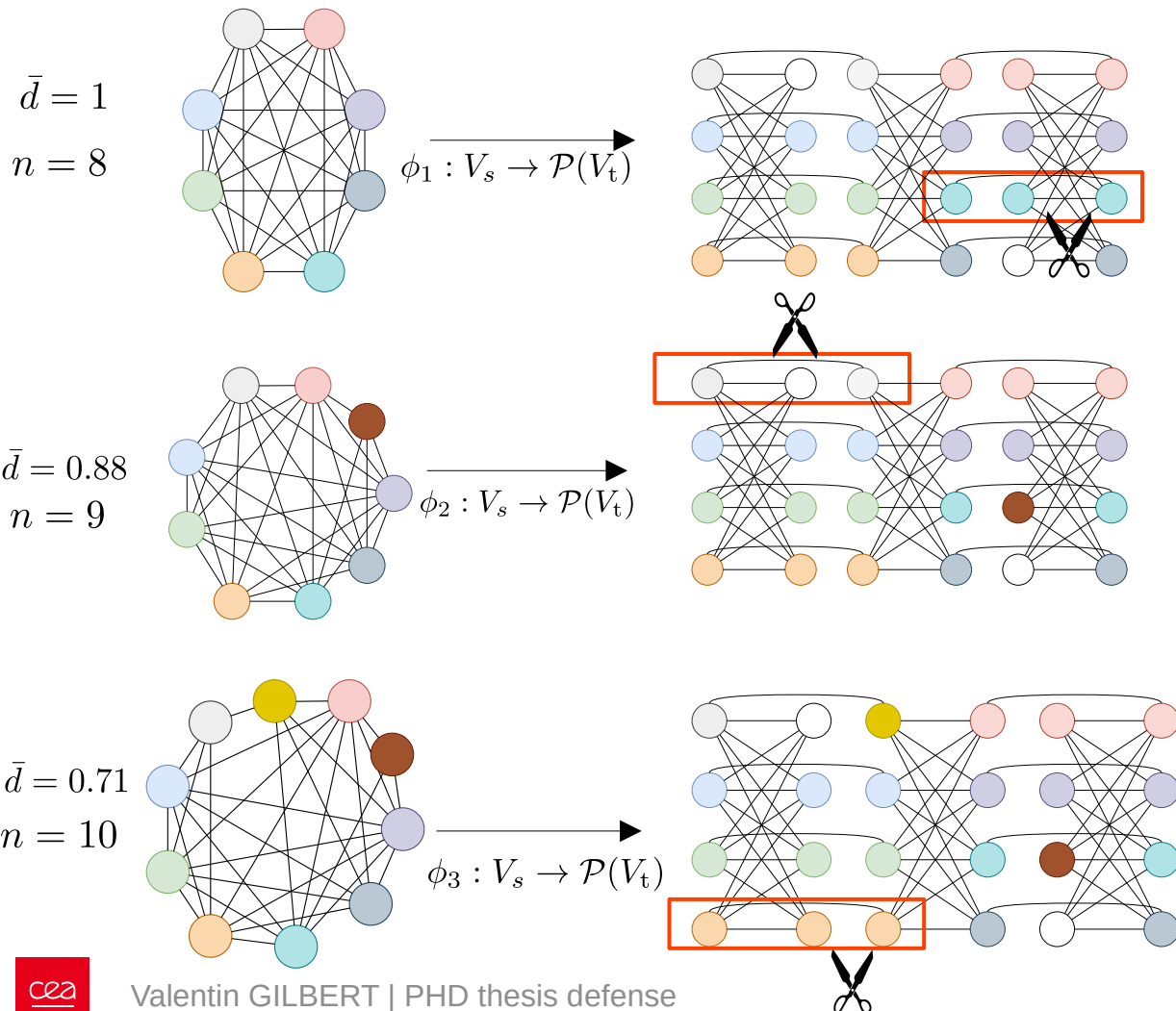
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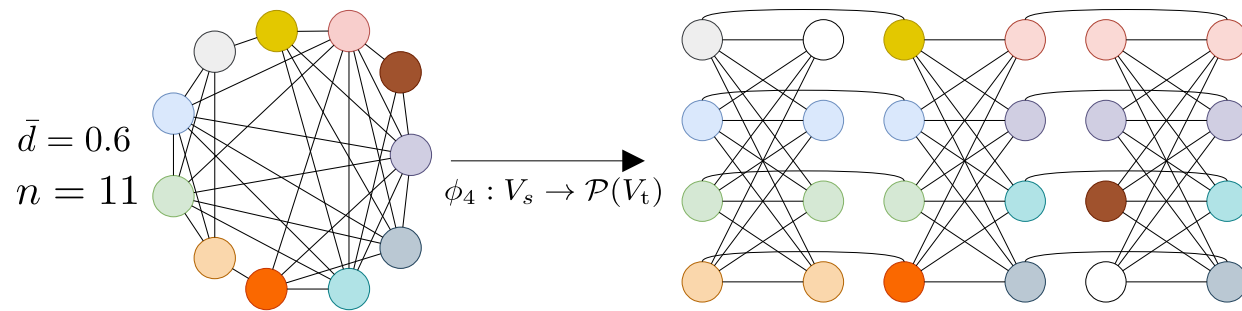
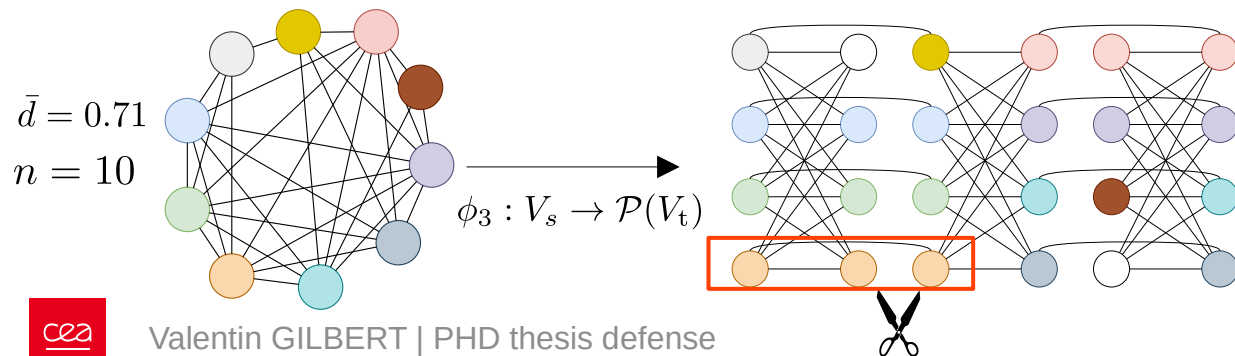
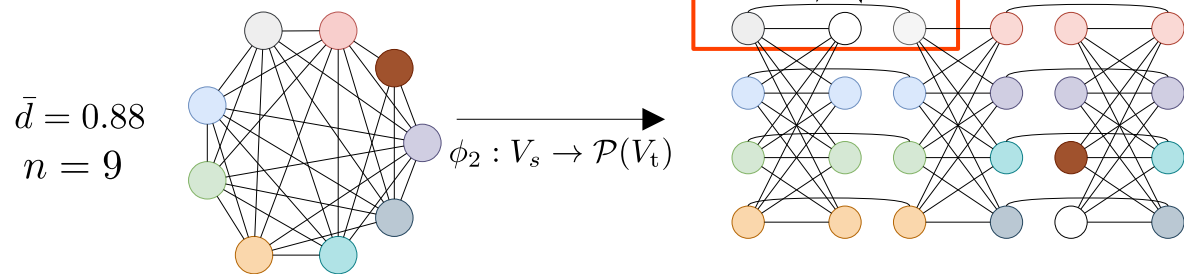
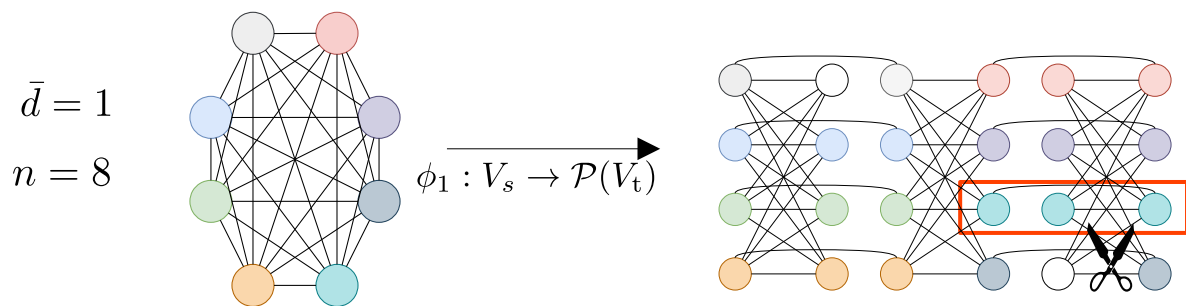
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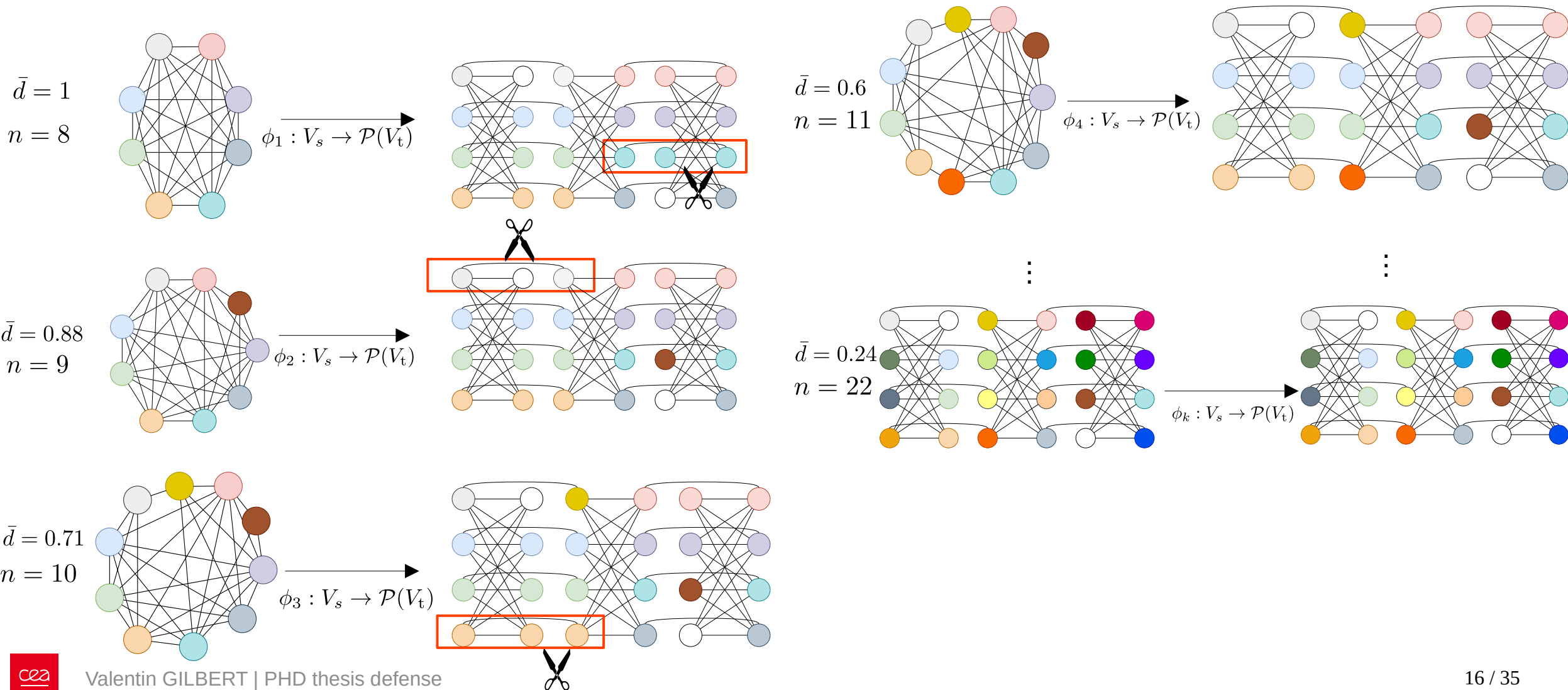
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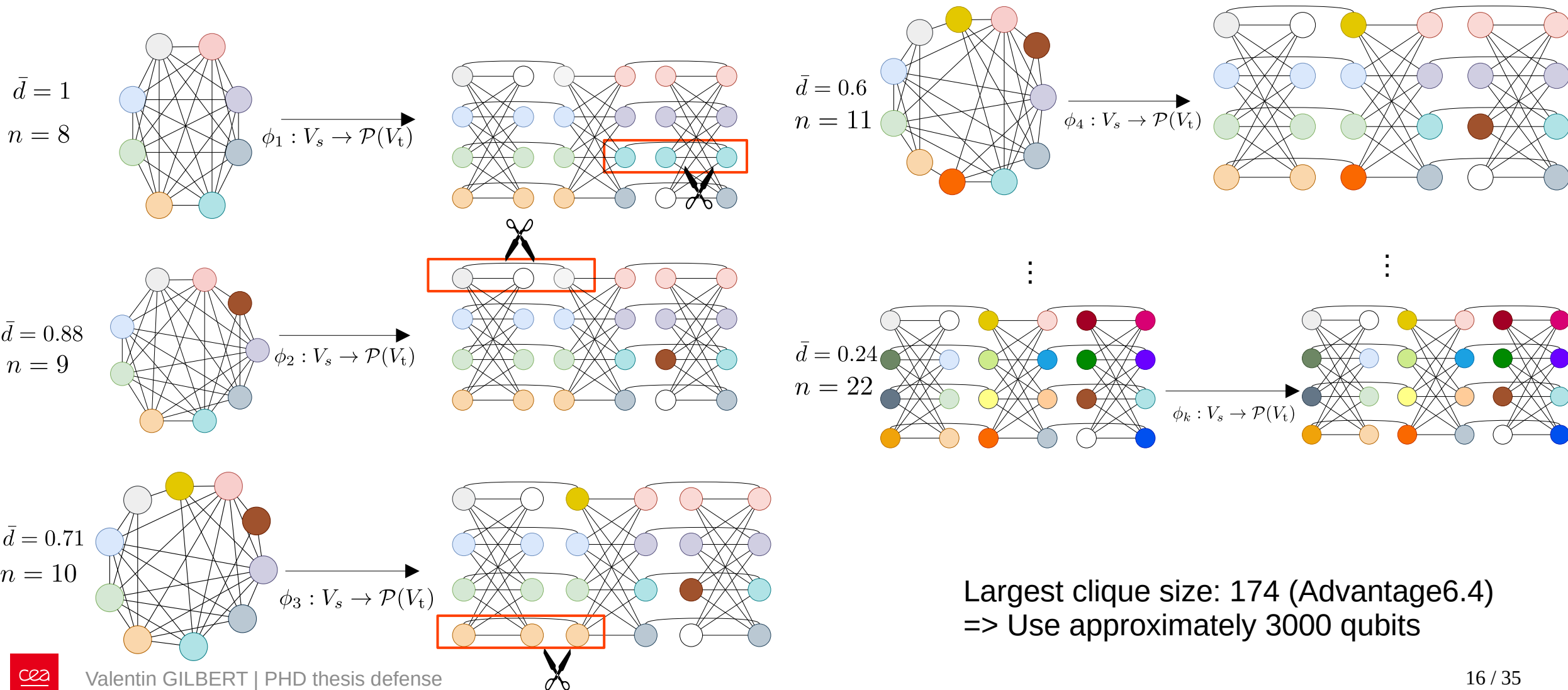
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2- Performance assessment

- Time To Solution metric (Gold standard) [RWJ⁺14]

Number of runs:

$$R = \left\lceil \frac{\log(1 - p)}{\log(1 - s)} \right\rceil$$

p : probability of getting the ground state in R runs

s : Empirical success probability

$$TTS = t_a \times R$$

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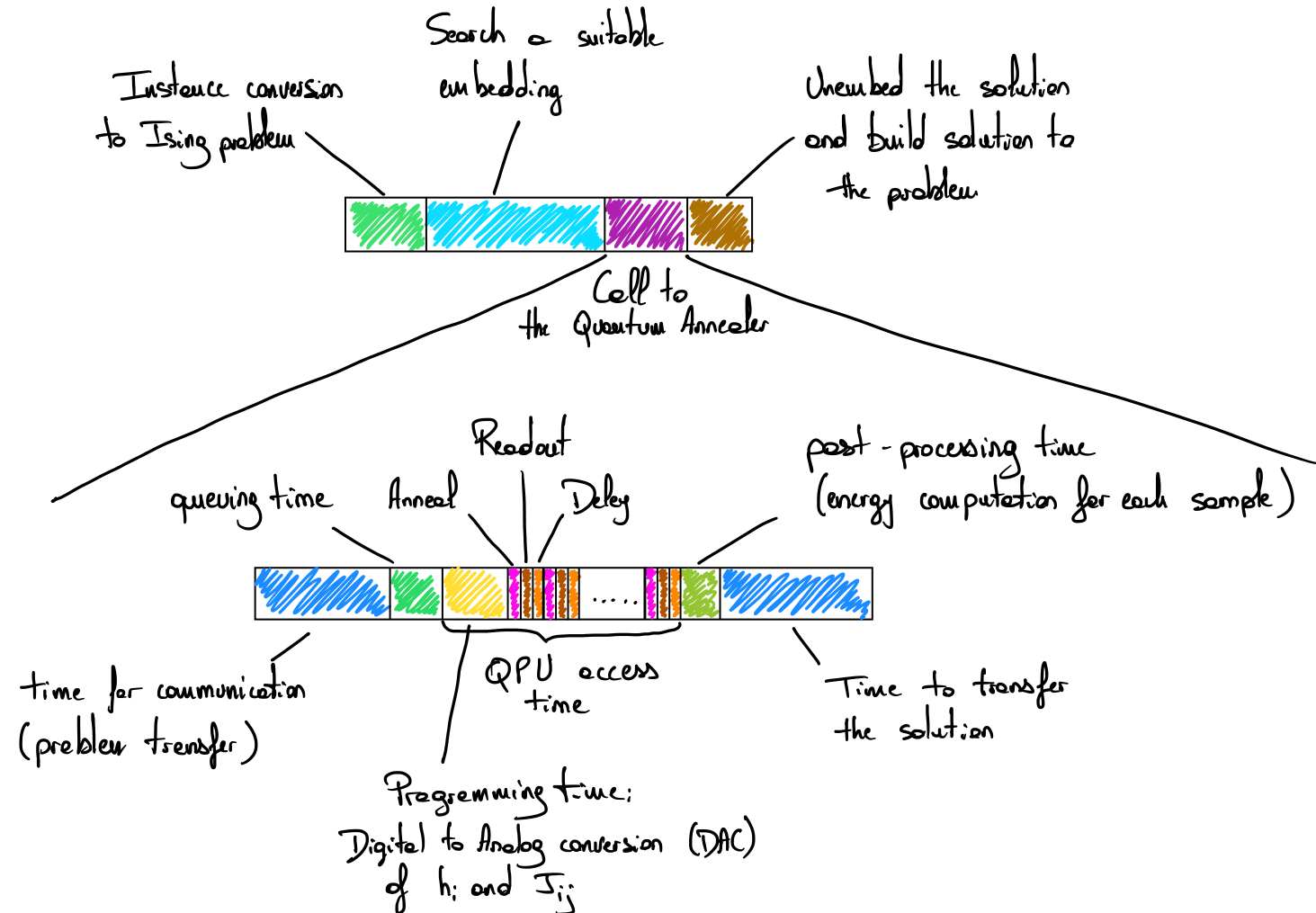
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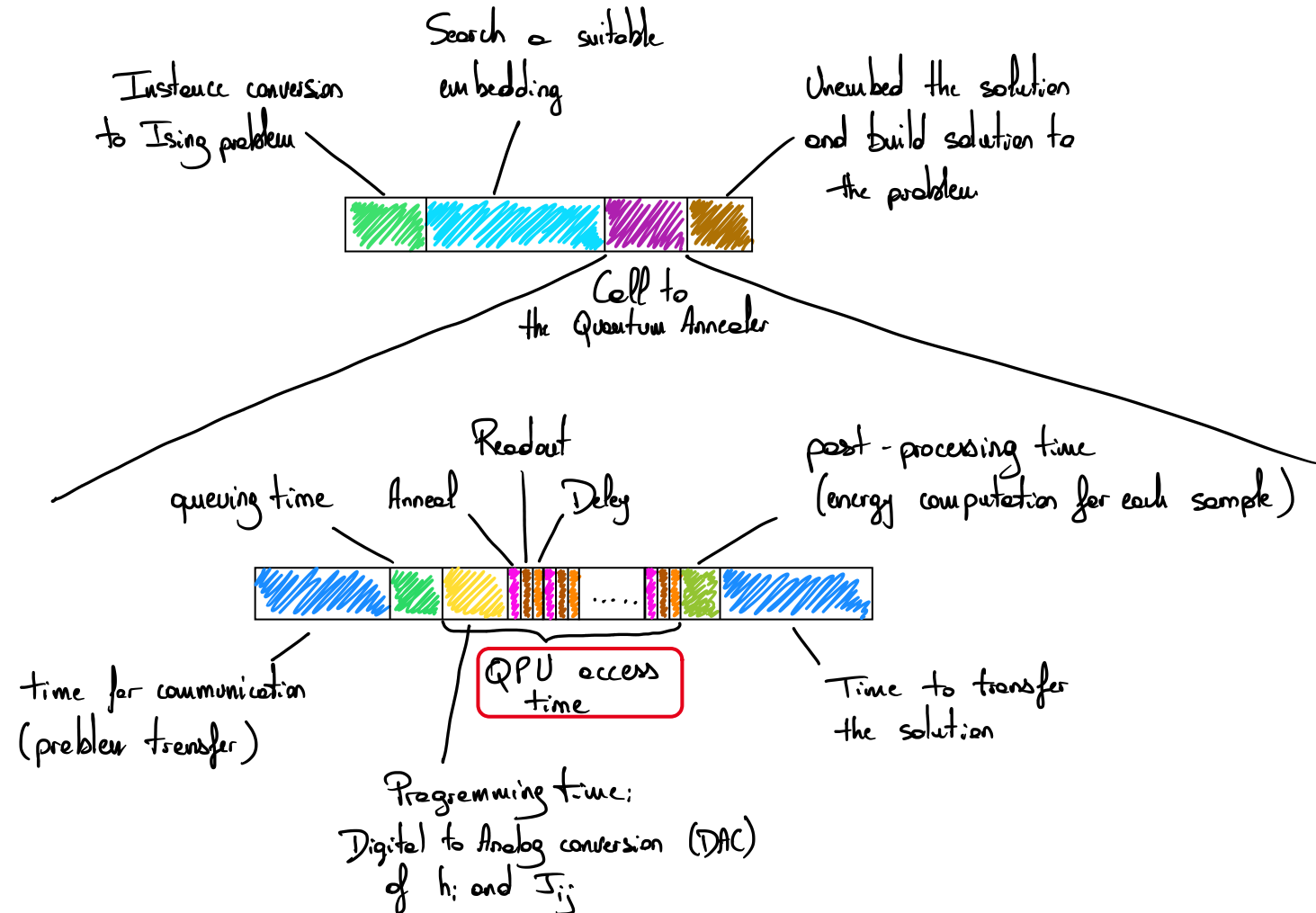
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2- Results – Max-cut problem

Density	0.02	0.03	0.04	0.05	0.06	0.07	0.08	0.09	0.1	0.2	0.3	0.4	0.5	0.6	0.7	0.8	0.9
Avg $ V_s $	1318	1062	912	810	737	680	635	597	565	395	321	277	248	226	209	195	184

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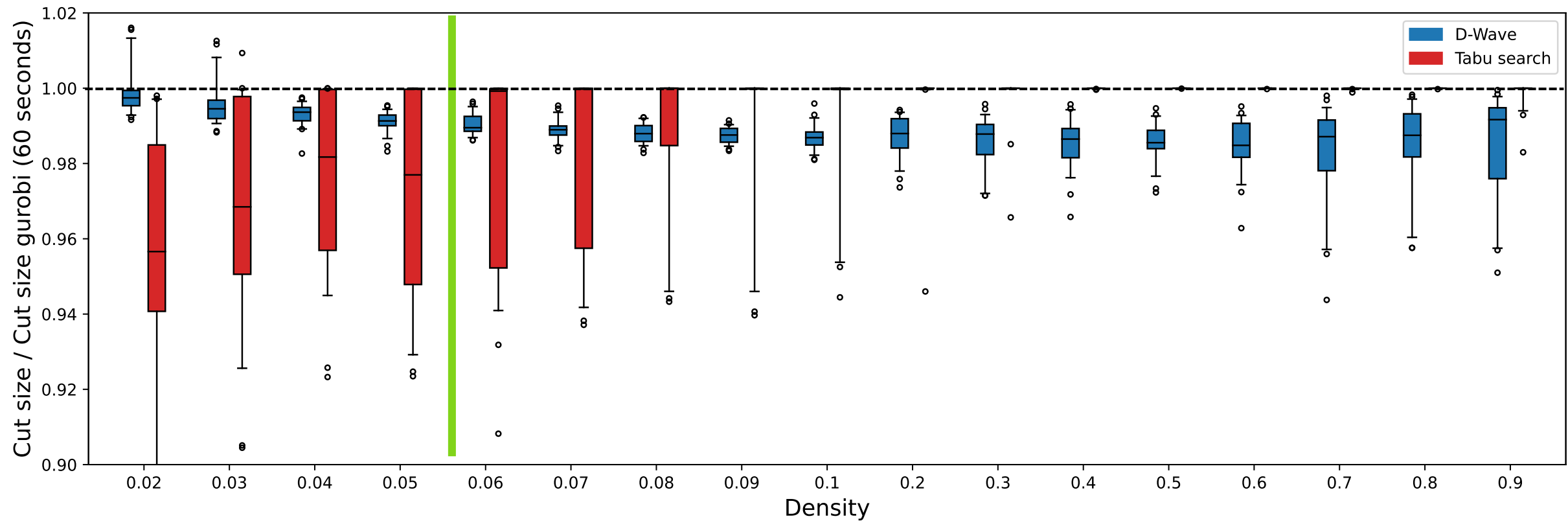
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Random Greedy Search
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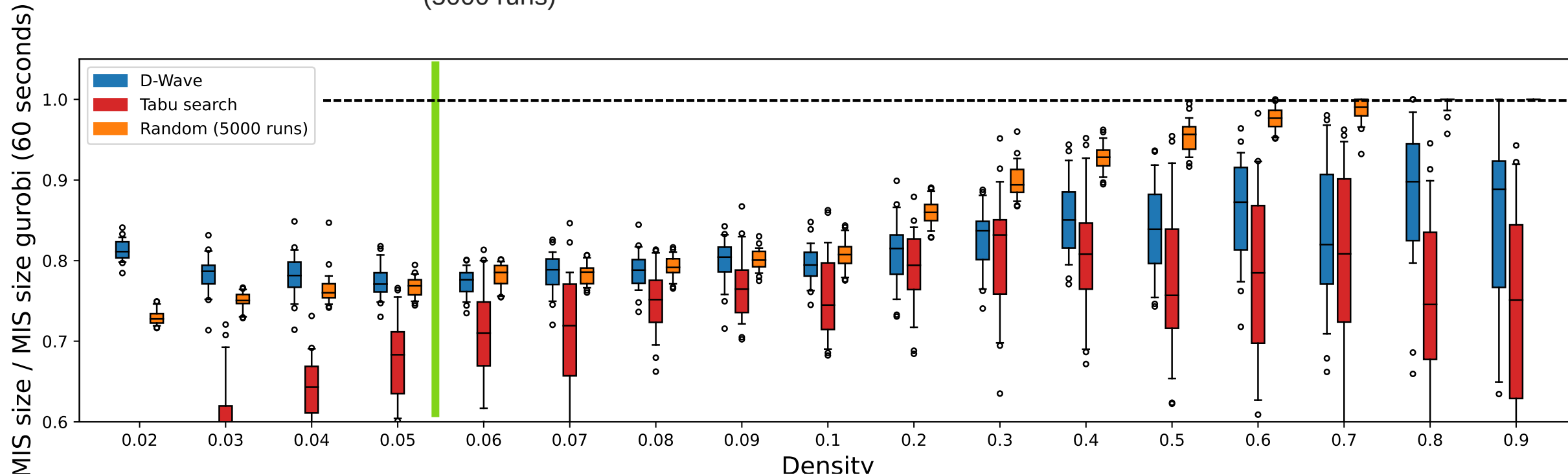
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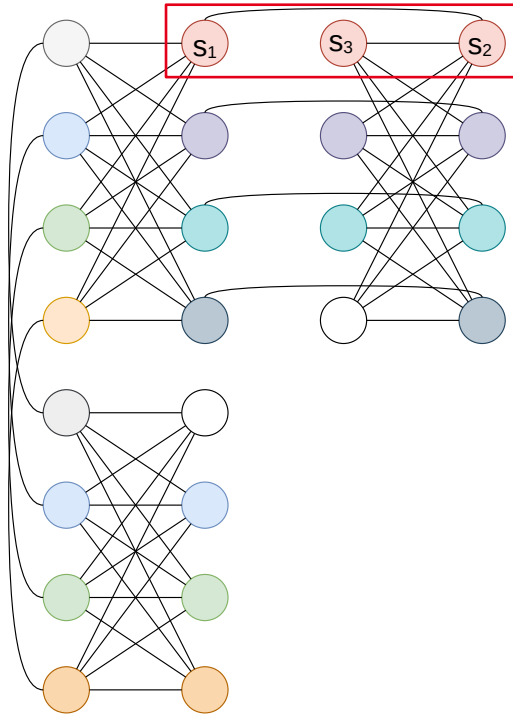
3 ■ Contribution #2 Increasing the performance of Quantum Annealers



Generated with openai

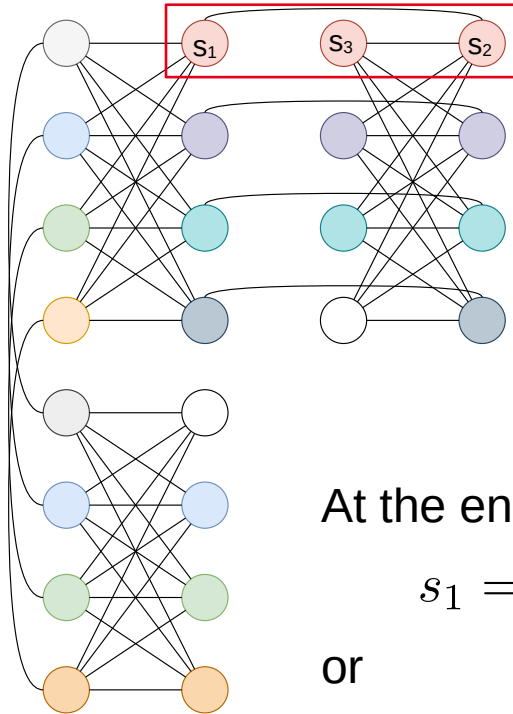
3- The embedding problem II

- Embedding step produces chains of qubits



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At the end of the annealing:

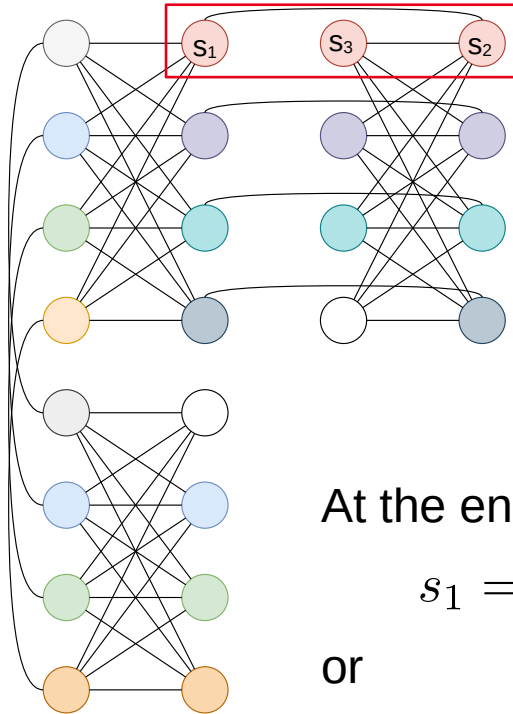
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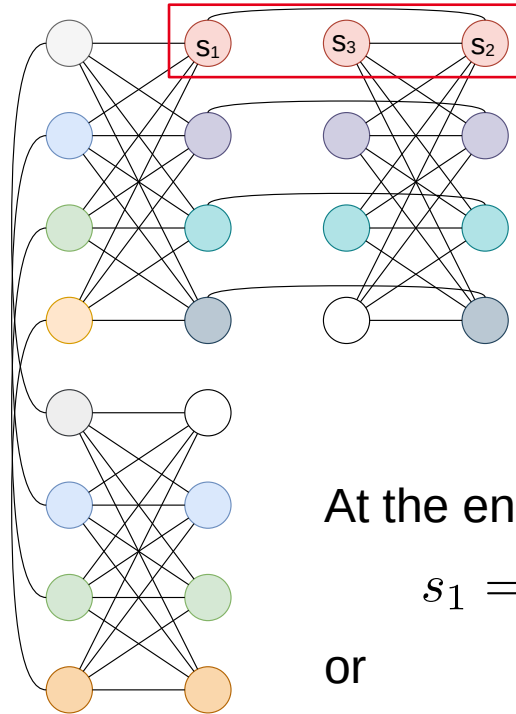
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Let's set:

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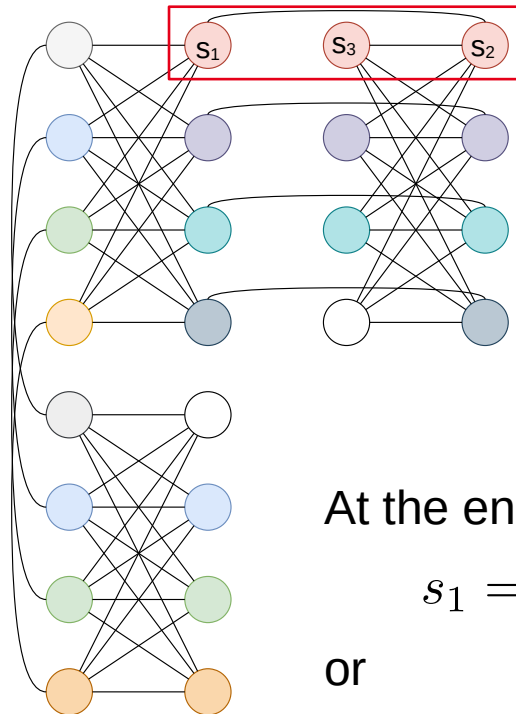
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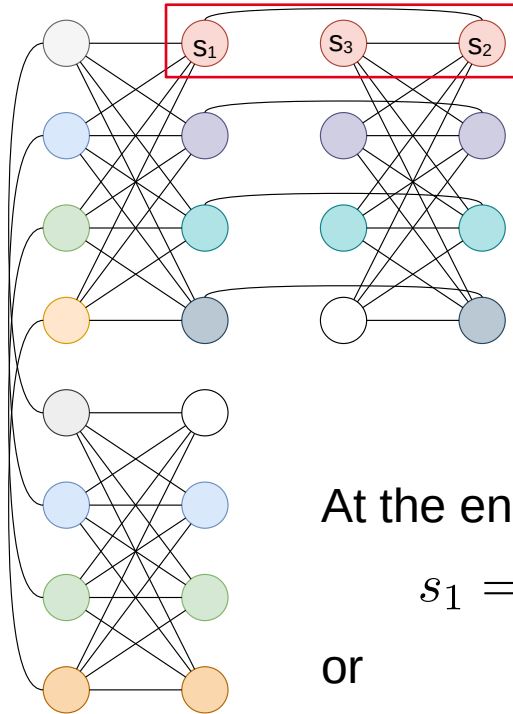
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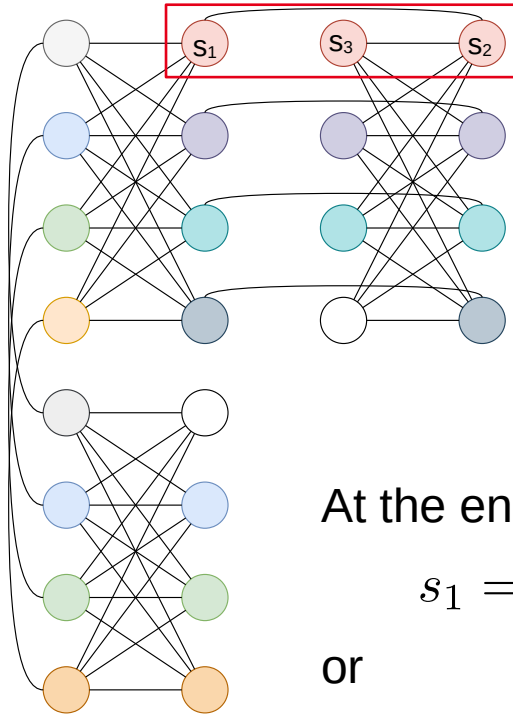
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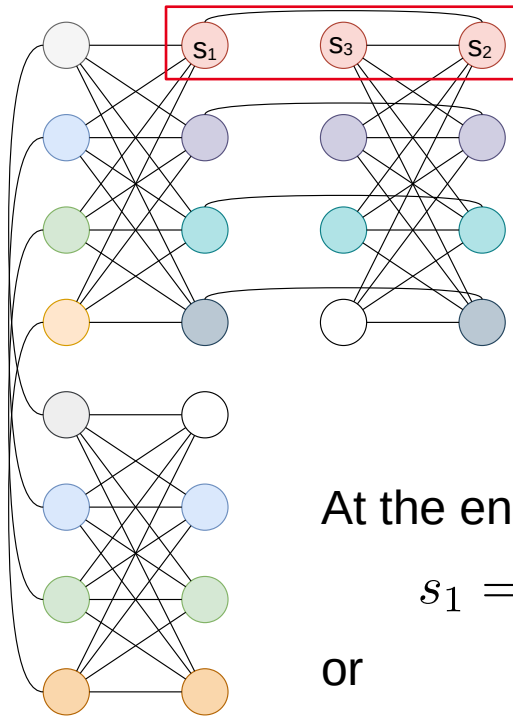
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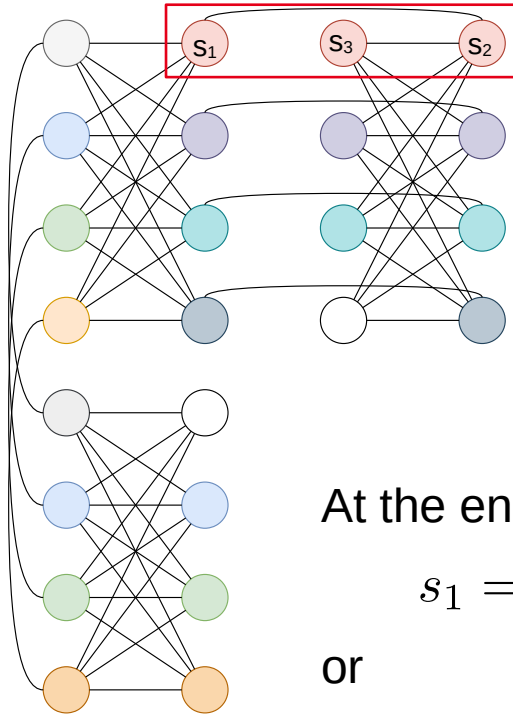
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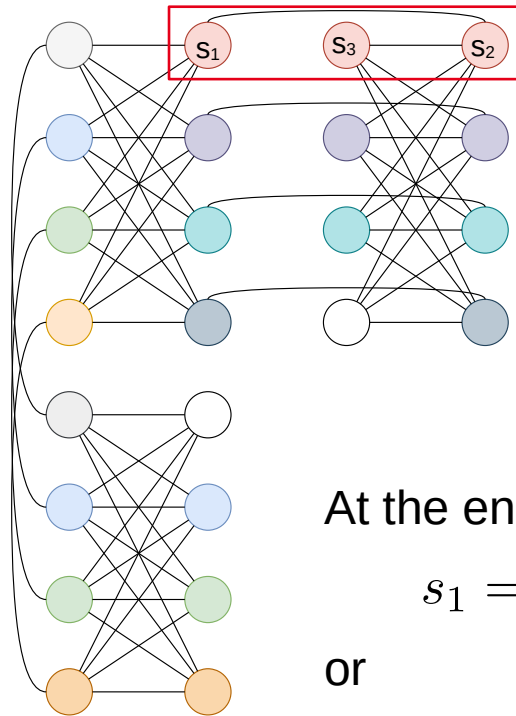
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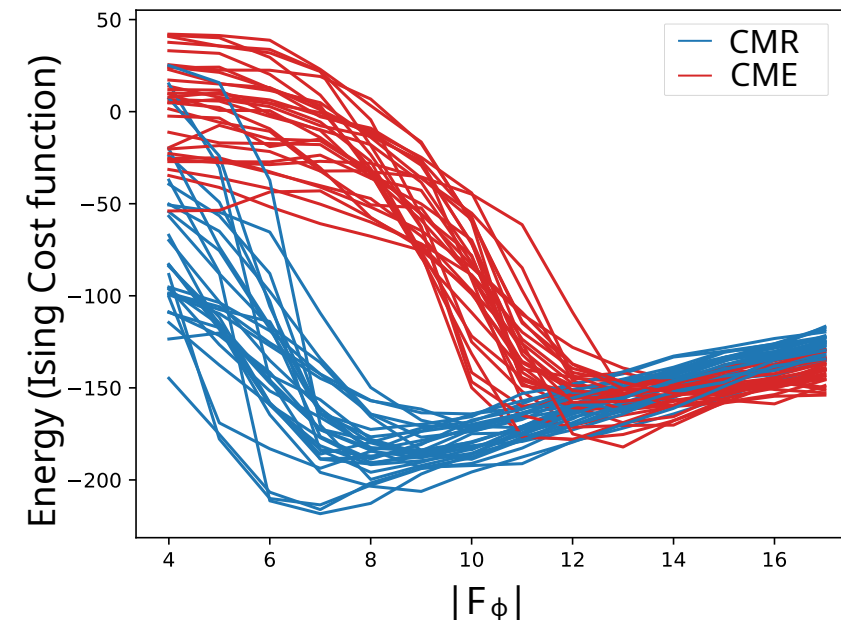


Limited programming range

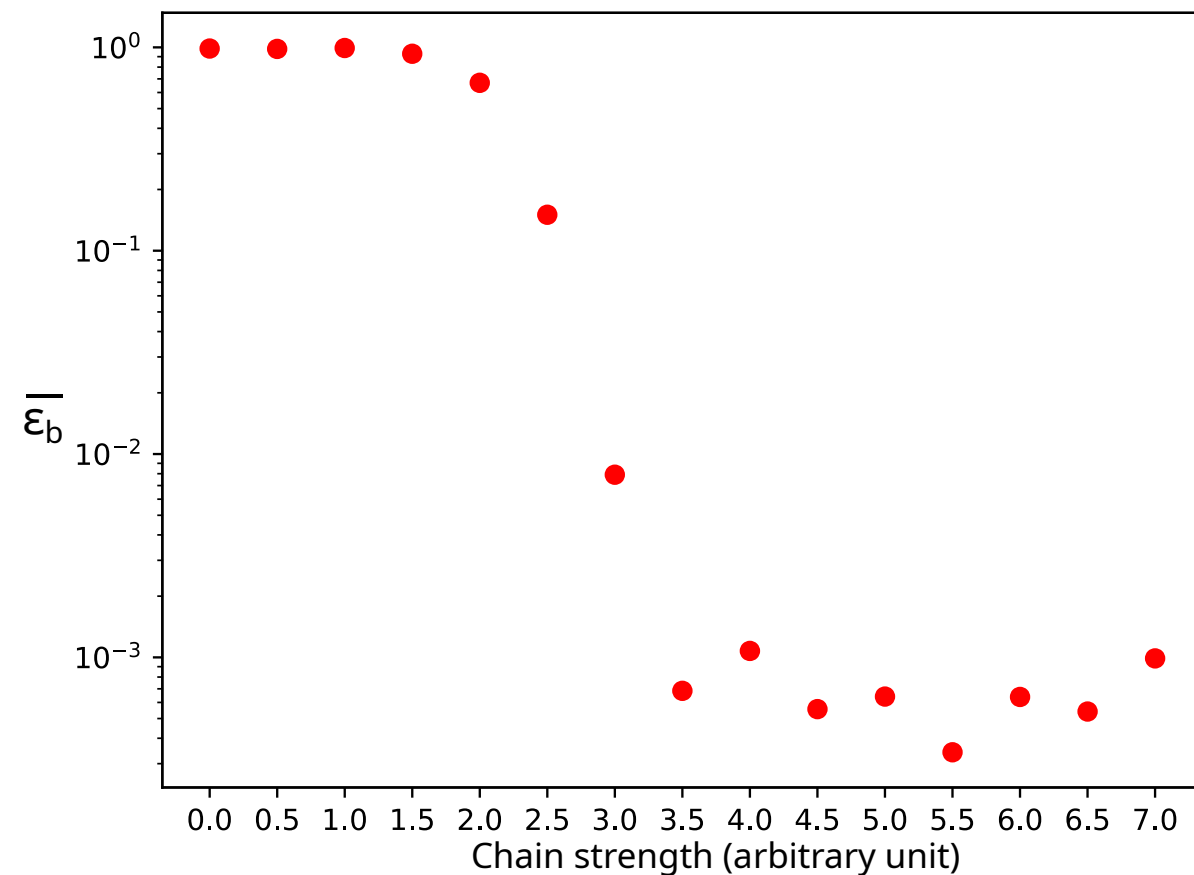
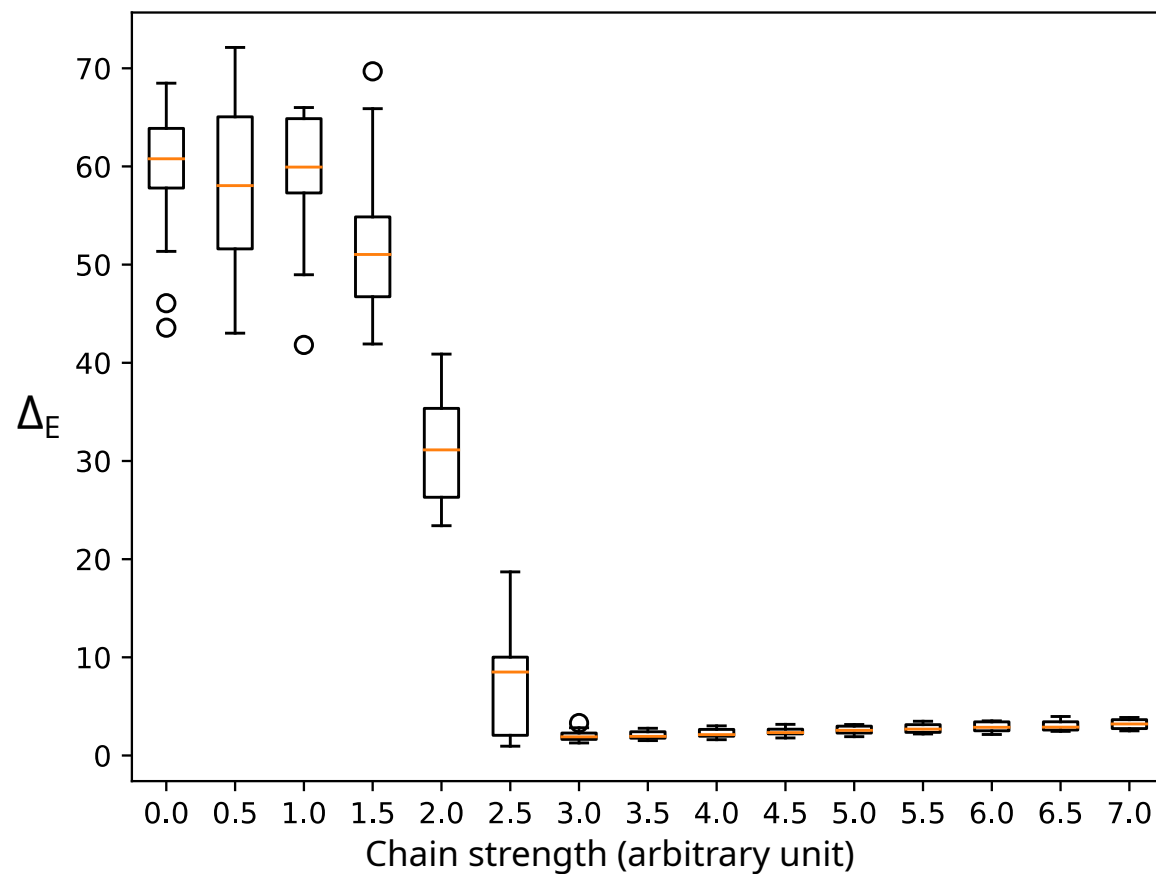
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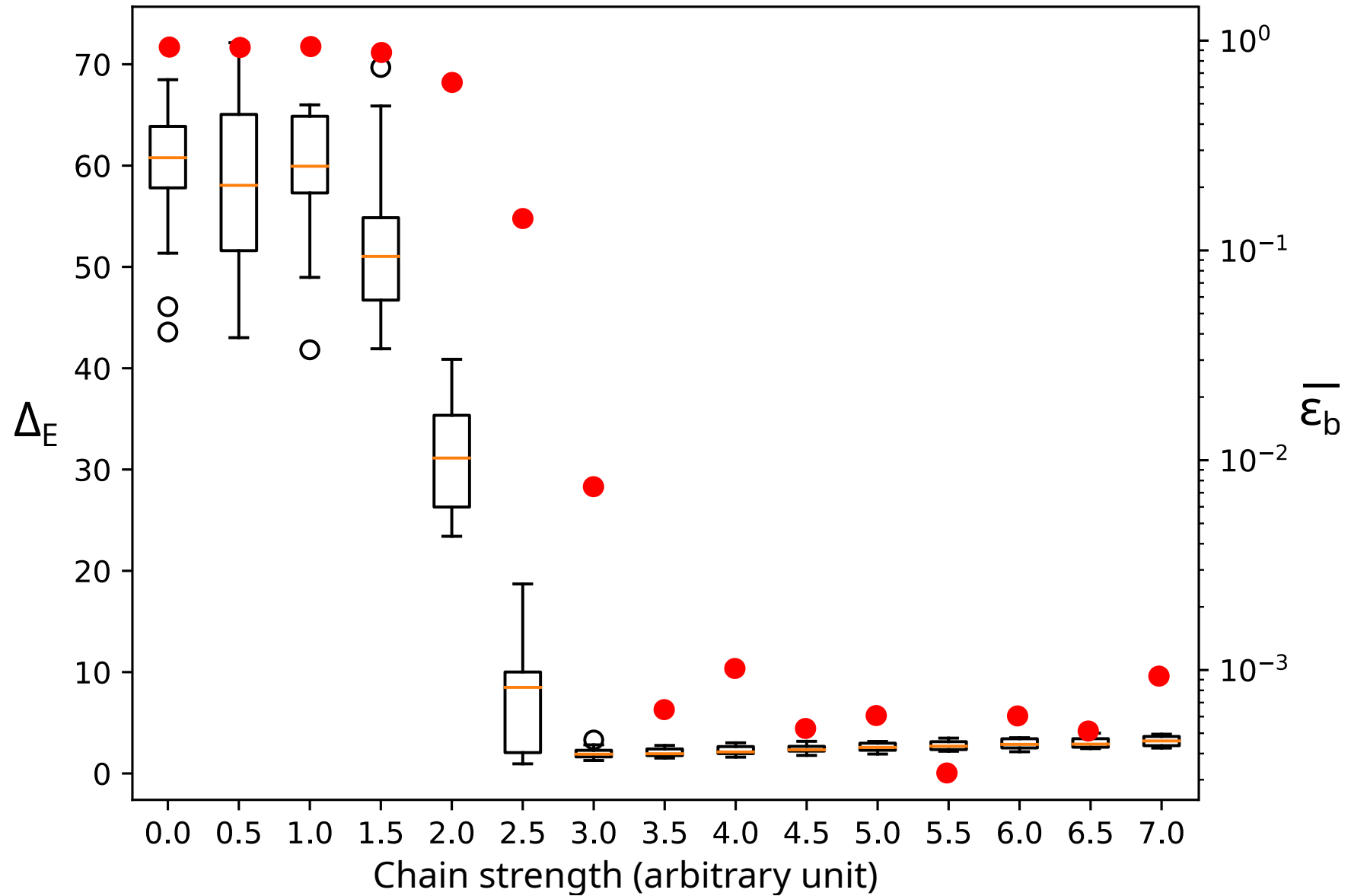
- Global chain strength: $F_\phi < 0$



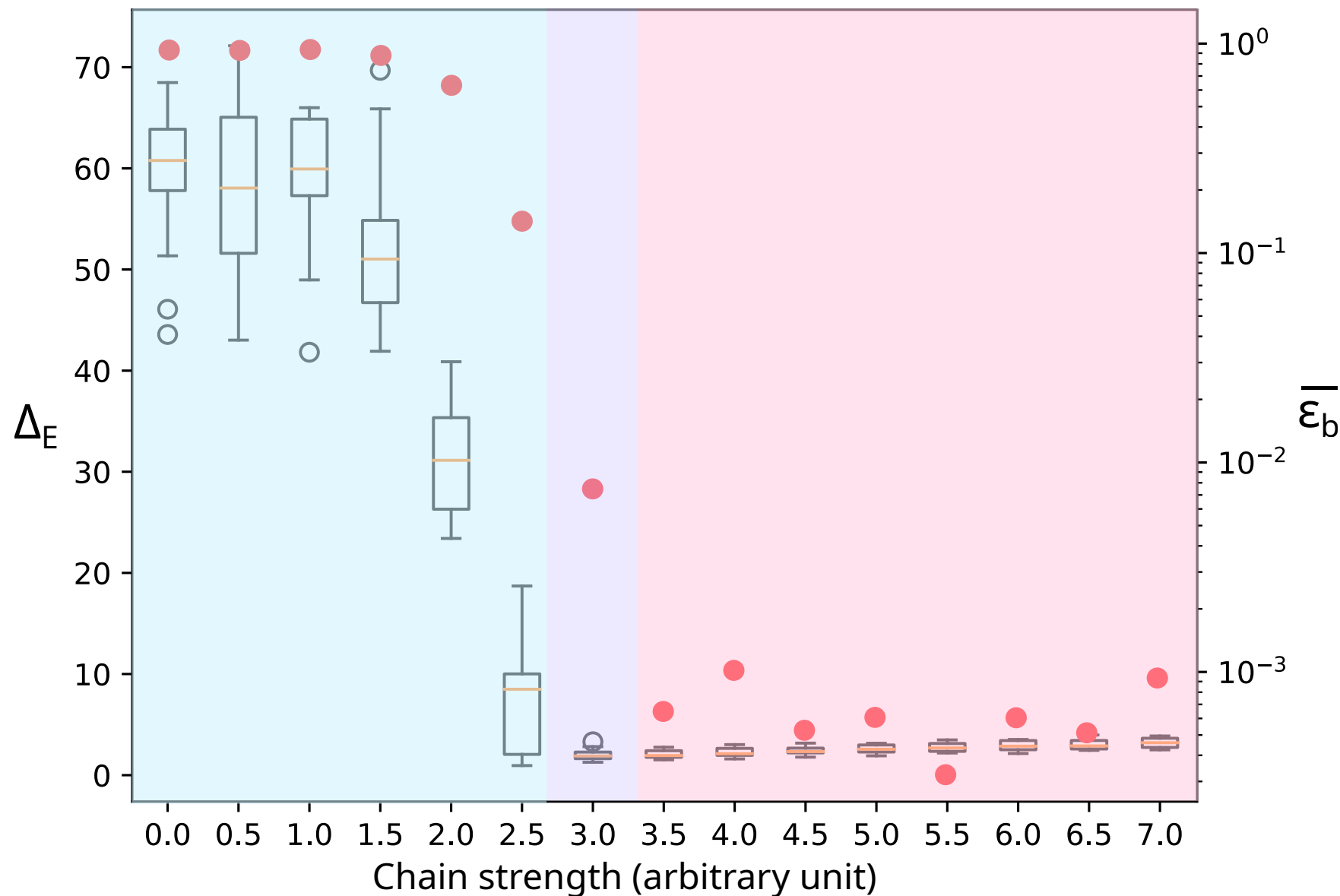
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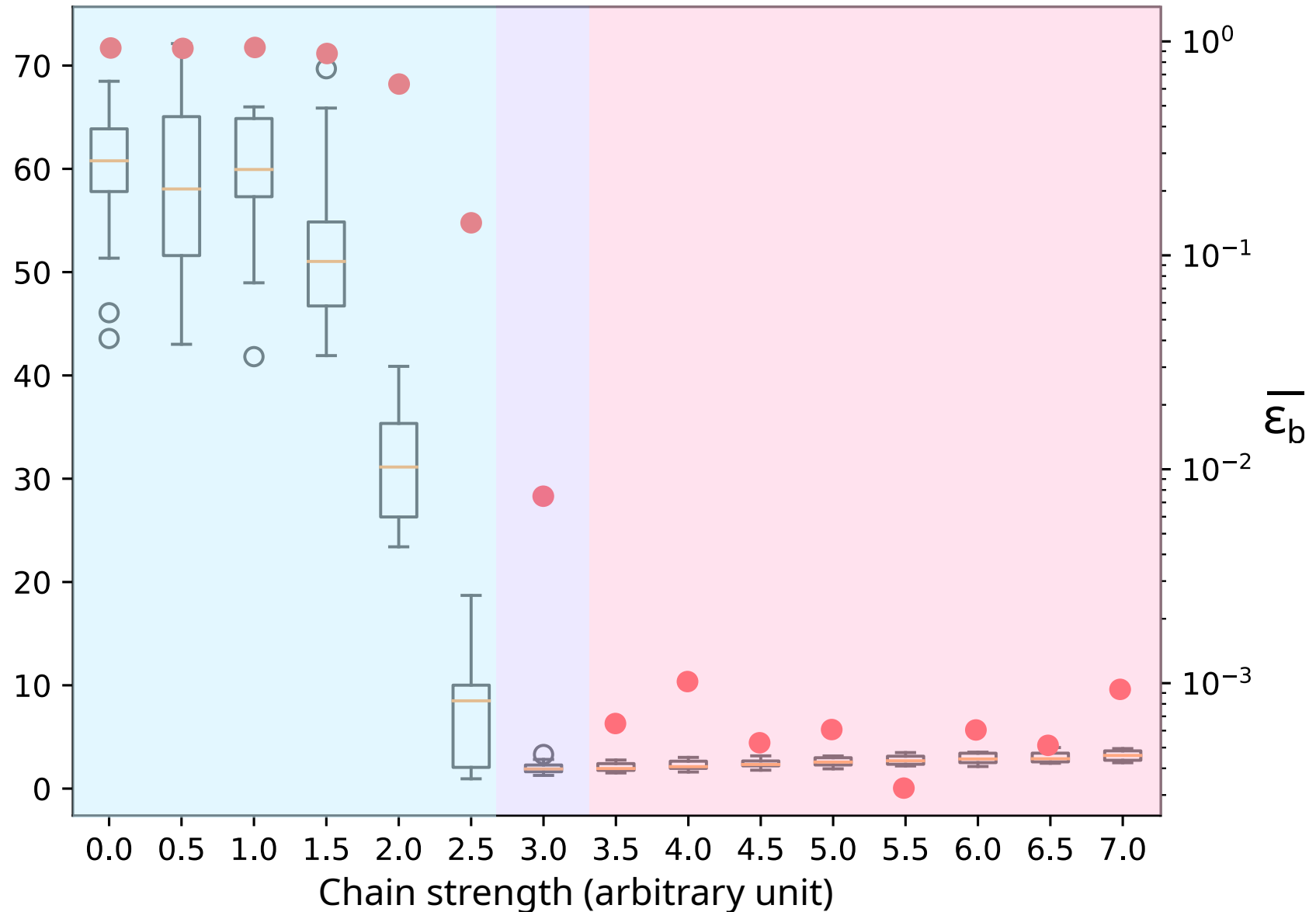


3- Chain scan & Chain breaks

Chain break intervals:
Advantage2:
 $[6 \times 10^{-3}, 2 \times 10^{-2}]$

Advantage6.4:
 $[2 \times 10^{-2}, 5 \times 10^{-2}]$

Δ_E



3- Results

- 30 instances of unweighted max-cut for each density
- Shot / pre-processing step: 128
- Final run number of shots:
Advantage2: 3072
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Advantage2_prototype2.2			Best cut size	Cut size improvement				Step
Instance size	Density	Embedding		min	max	mean	std	
$n = 40$	0.1	CMR	66.4	+0%	+0%	+0%	0%	5.4
	0.5	CMR	243	+0%	+2%	+0.2%	0%	5.8
	0.9	CMR	362.8	+2.1%	+8.2%	+5%	1.6%	4.6
$n = 80$	0.1	CMR	235.5	+0%	+0%	+0%	0%	4.7
	0.5	CME	804	+9.8%	+17.2%	+12.5%	0.2%	4.2
	0.9	CME	1435	+2%	+4.7%	+3.2%	0.6%	4.2
Advantage6.4			Best Cut size	Cut size improvement				Step
Instance size	Density	Embedding		min	max	mean	std	
$n = 100$	0.1	CMR	355.9	+0%	+0.3%	+0%	0%	4.5
	0.5	CME	1271.4	+5.6%	+14.5%	+8.8%	1.8%	2.7
	0.9	CME	2243	+1.4%	+3.7%	+2.5%	0.5%	3.7
$n = 170$	0.1	CMR	950.8	-2.1%	+0.6%	-0.5%	0.5%	2.1
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“NISQ will not change the world by itself, at least not right away; instead we should regard it as a step toward more powerful quantum technologies we hope to develop in the future.”

J. Preskill [\[Pre21\]](#)



Thank you !

A big thank to ACTIF (Association CEA des thésard.e.s d'Île-de-France)



New president: Lise.jolicoeur@cea.fr

Publications

G. Bettonte, V. Gilbert, D. Vert, S. Louise, and R. Sirdey, “Quantum approaches for wcet-related optimization problems,” in Lecture Notes in Computer Science - ICCS 2022, p. 202-217, Springer International Publishing, 2022

V. Gilbert, J. Rodriguez, S. Louise, and R. Sirdey, “Solving higher order binary optimization problems on nisd devices: experiments and limitations,” in Lecture Notes in Computer Science - ICCS 2023, p.224-232, Springer Nature Switzerland, 2023

V. Gilbert, S. Louise, and R. Sirdey, “Taqos: A benchmark protocol for quantum optimization systems,” in Lecture Notes in Computer Science – ICCS 2023, p.168-176, Springer Nature Switzerland, 2023

V. Gilbert, and S. Louise, “Quantum annealers chain strengths: A simple heuristic to set them all,” in Lecture Notes in Computer Science - ICCS 2024, p.292-306, Springer Nature Switzerland, 2024

V. Gilbert, J. Rodriguez, and S. Louise “Benchmarking quantum annealers with near-optimal minor-embedded instances,” in 2024 IEEE International Conference on Quantum Computing and Engineering (QCE), IEEE, 2024

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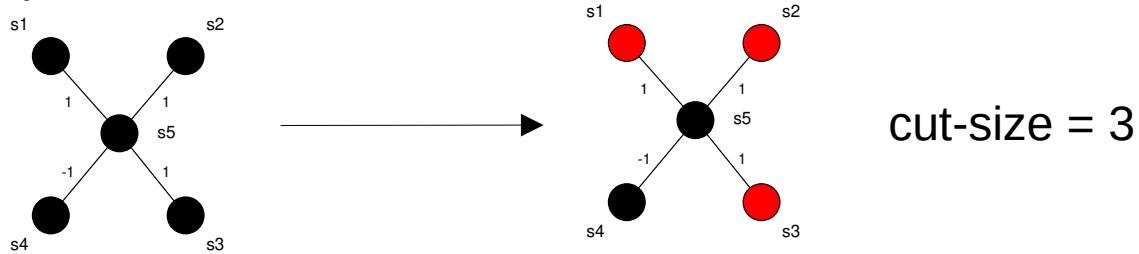
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4 ■ Appendix

Selection of best Annealing time

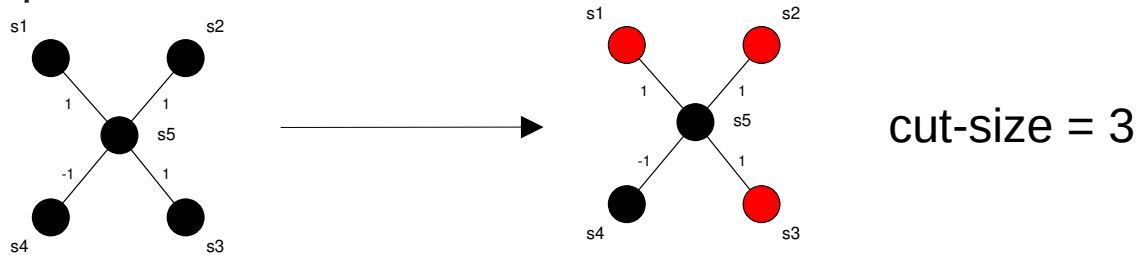
- Quality metrics are problem-dependent
 - Max-cut problem: The cut size



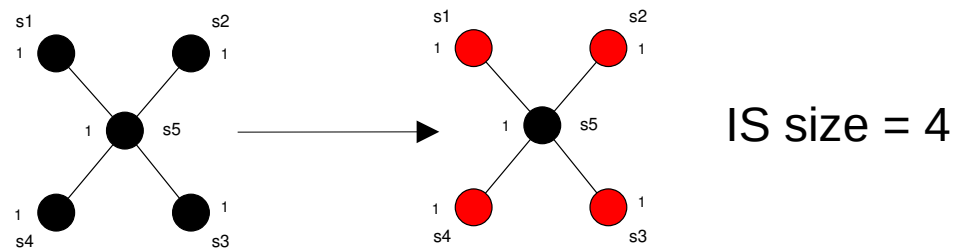
Selection of best Annealing time

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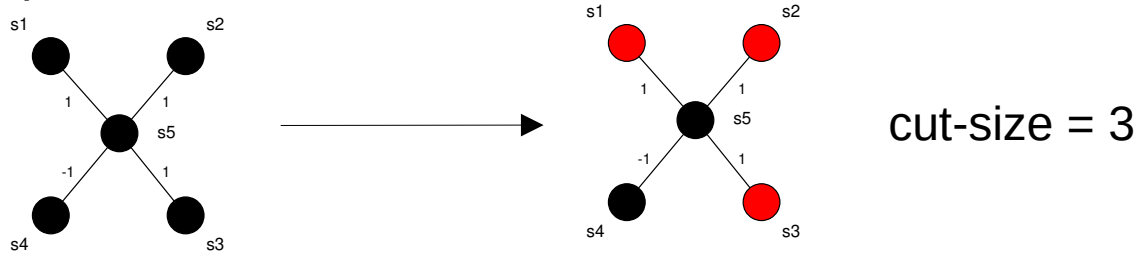
- Maximum Independent set problem: The size of the independent set.



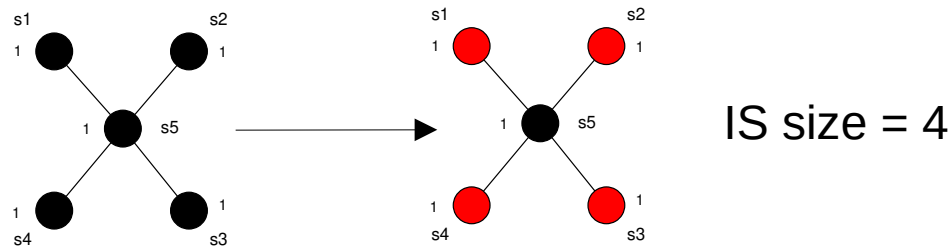
Selection of best Annealing time

- Quality metrics are problem-dependent

- Max-cut problem: The cut size



- Maximum Independent set problem: The size of the independent set.



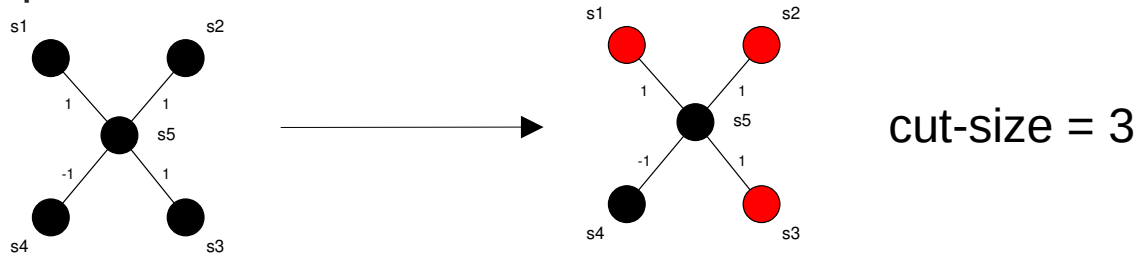
- Time measurement

- 60s time window for branch & bound algorithm
- 1s time window for Tabu Search (C implementation)
- 1s time window for D-Wave Q

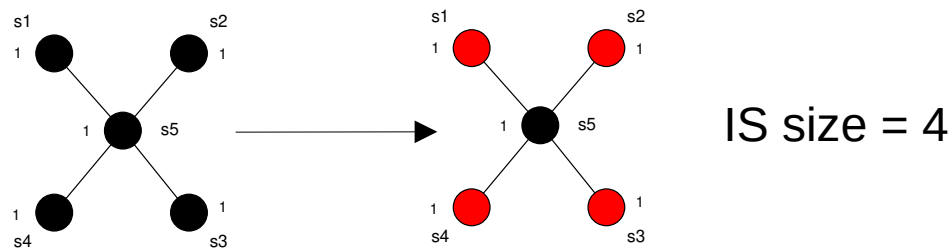
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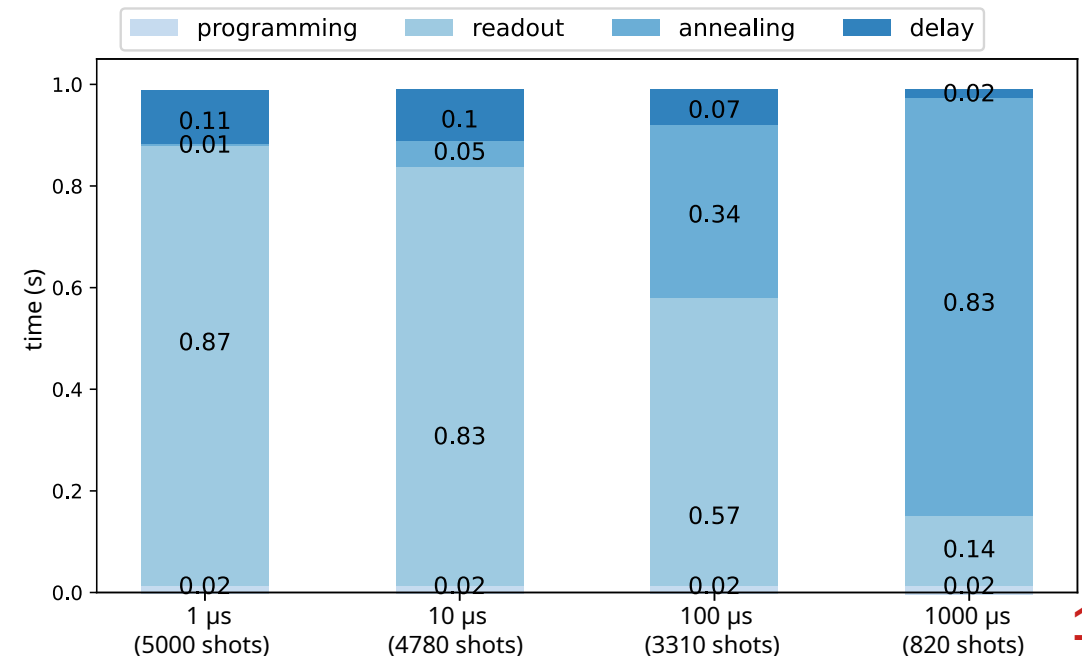


- Maximum Independent set problem: The size of the independent set.



- Time measurement

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Optimal mapping

- How to assess the quality of a mapping:
 - D-Wave QA has a regular topology
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$$n_{\phi(v)^*} = \left\{ \begin{array}{l} 1 \text{ if } \deg(v) \leq c_{\text{phys}} \\ 2 \text{ if } c_{\text{phys}} < \deg(v) \leq (2c_{\text{phys}} - 2) \\ \left\lceil \frac{\deg(v) - (2c_{\text{phys}} - 2)}{c_{\text{phys}} - 2} \right\rceil + 2 \text{ otherwise} \end{array} \right\}$$

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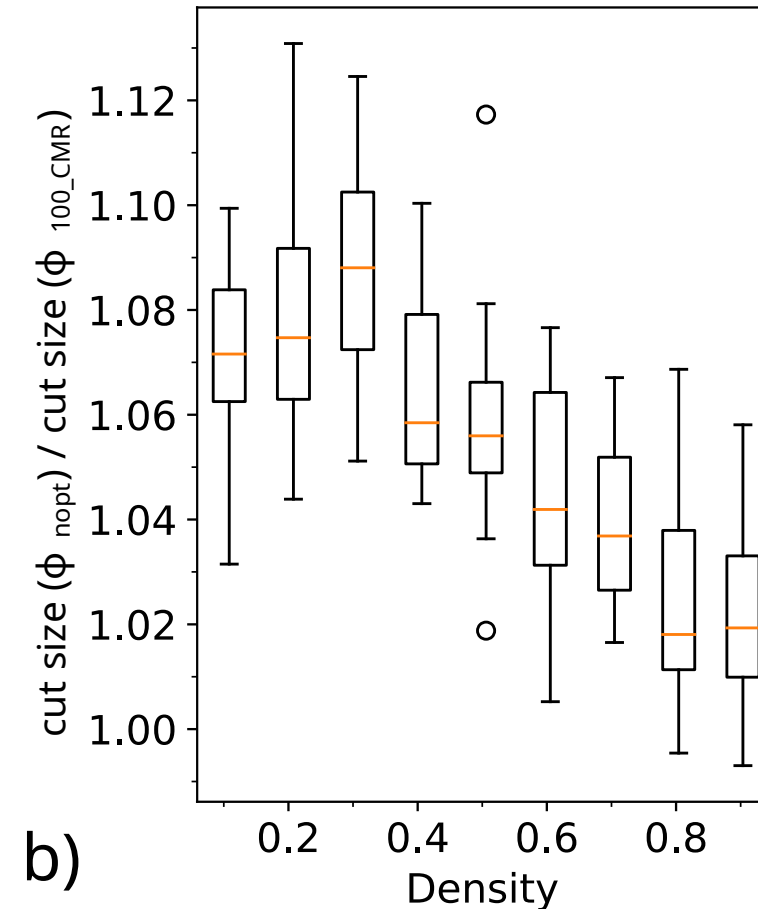
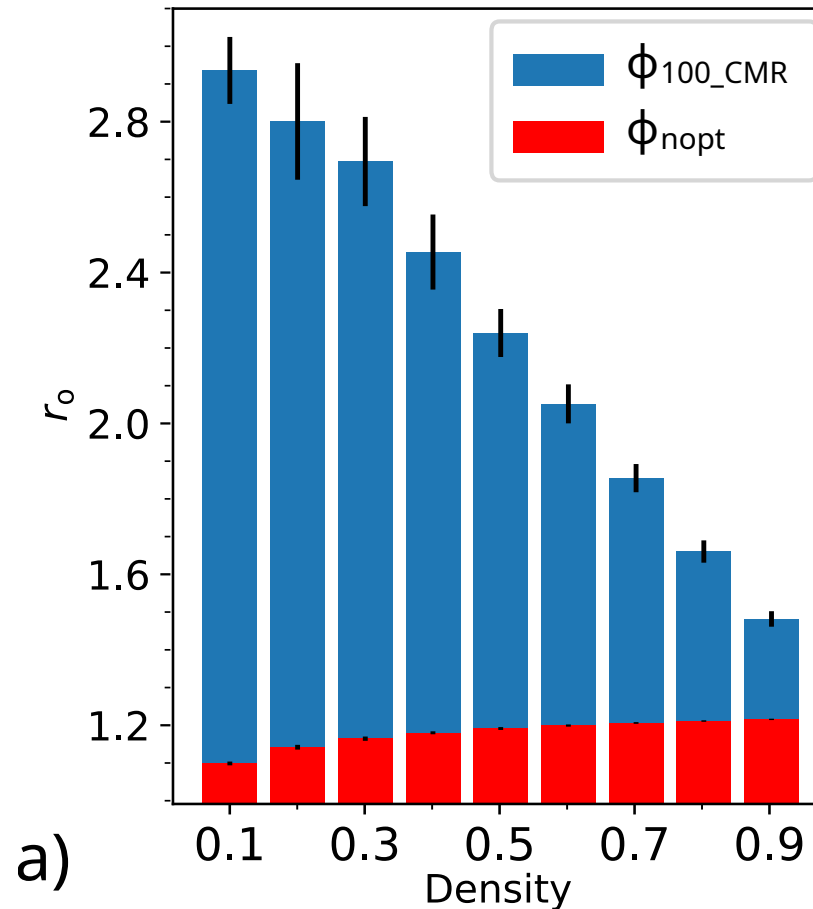
- Compute the overhead ratio considering this bound:

$$r_o = \frac{n_{\phi}}{\sum_{v \in V_s} n_{\phi(v)^*}}$$

Optimal mapping

- Comparison of the performance of our generation method against state of the art embedding method

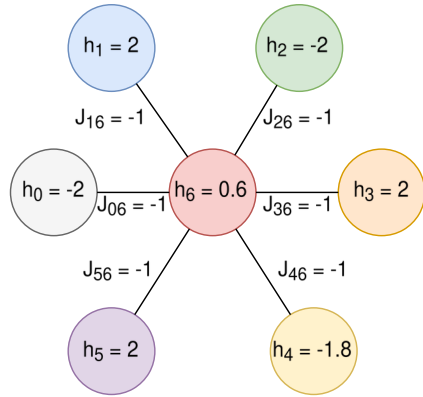
at $d = 1$
 $|V_s| = 100$
 $|V_t| = 982$



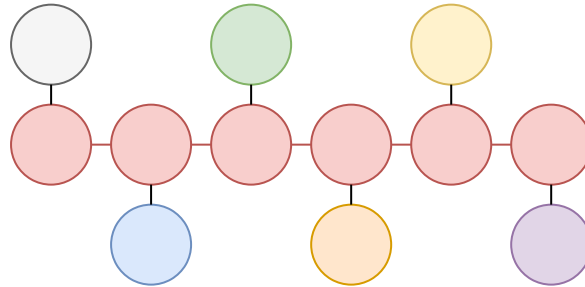
Assumption: Instances with less duplicated qubits are more easily solved by QA => Seems to be true

III- Shape of the logical qubit

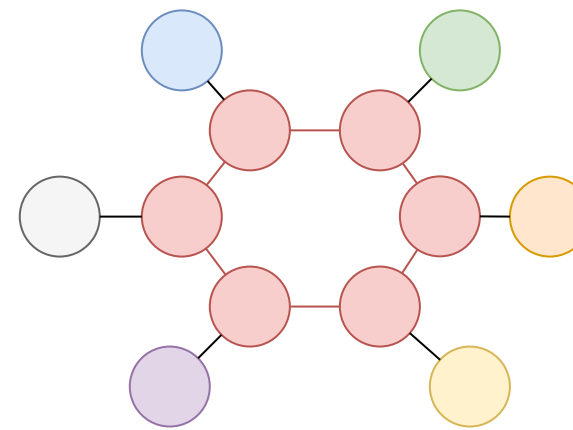
Source graph



Chain encoding



Cycle encoding



Clique encoding

